

Lexical Effect Handlers: Fast by Design, Correct by Proof

Cong Ma
University of Waterloo

Effect handler

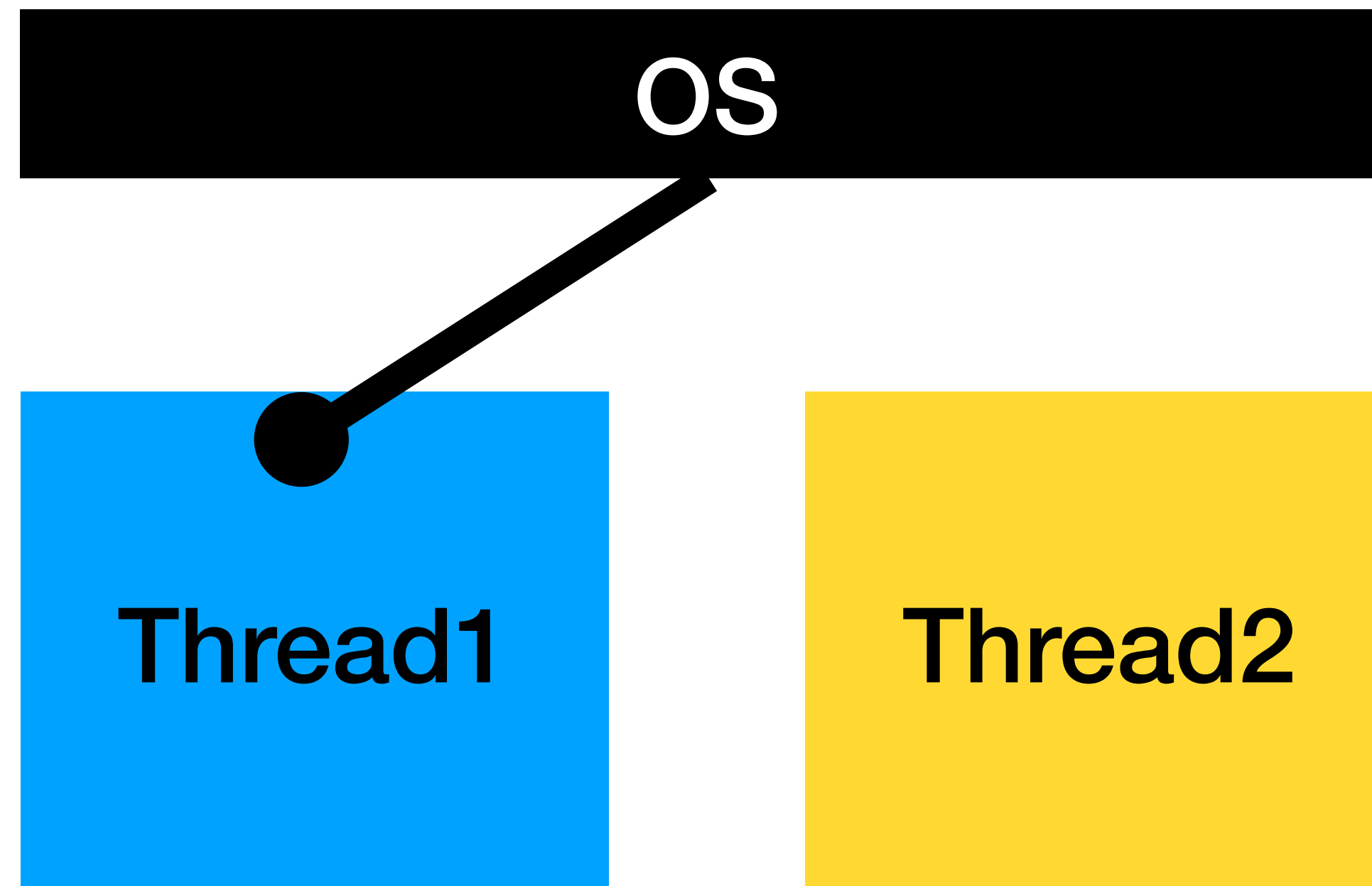
A powerful language feature

Effect handlers subsume many control flow features: `async/await`, `coroutine`, `generator`...

Effect handler

Scheduler inside OS

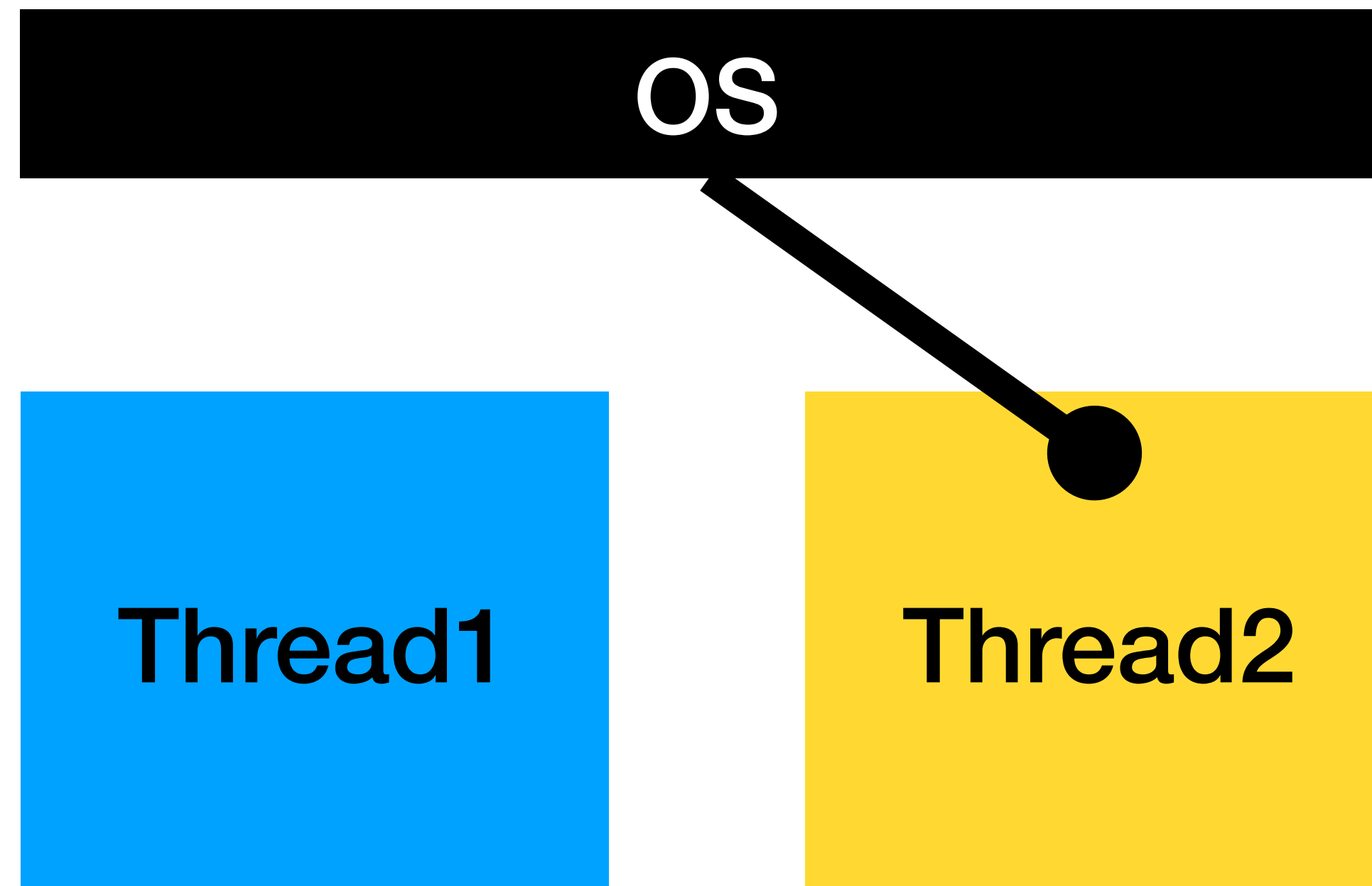
Background



Effect handler

Scheduler inside OS

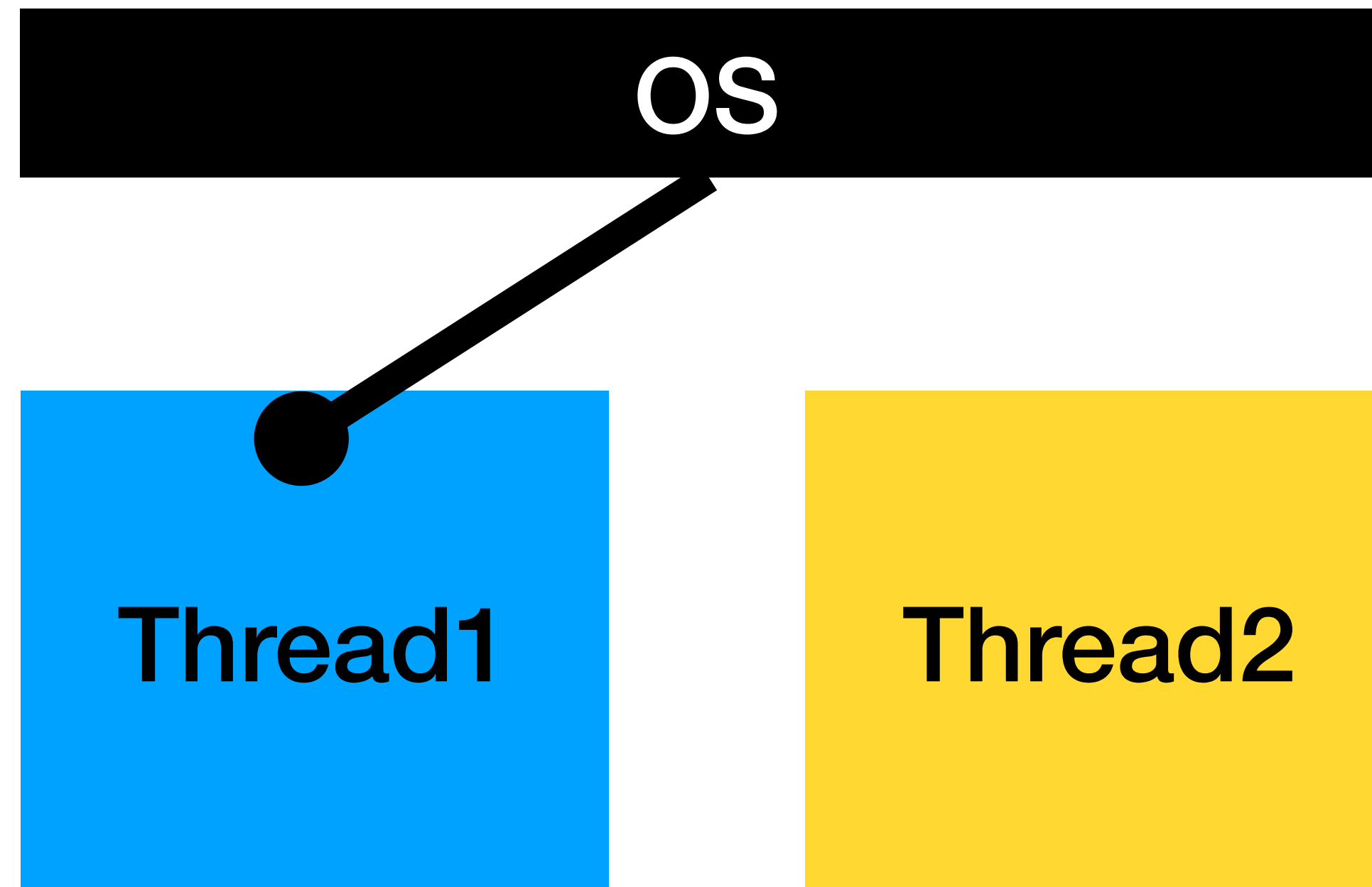
Background



Effect handler

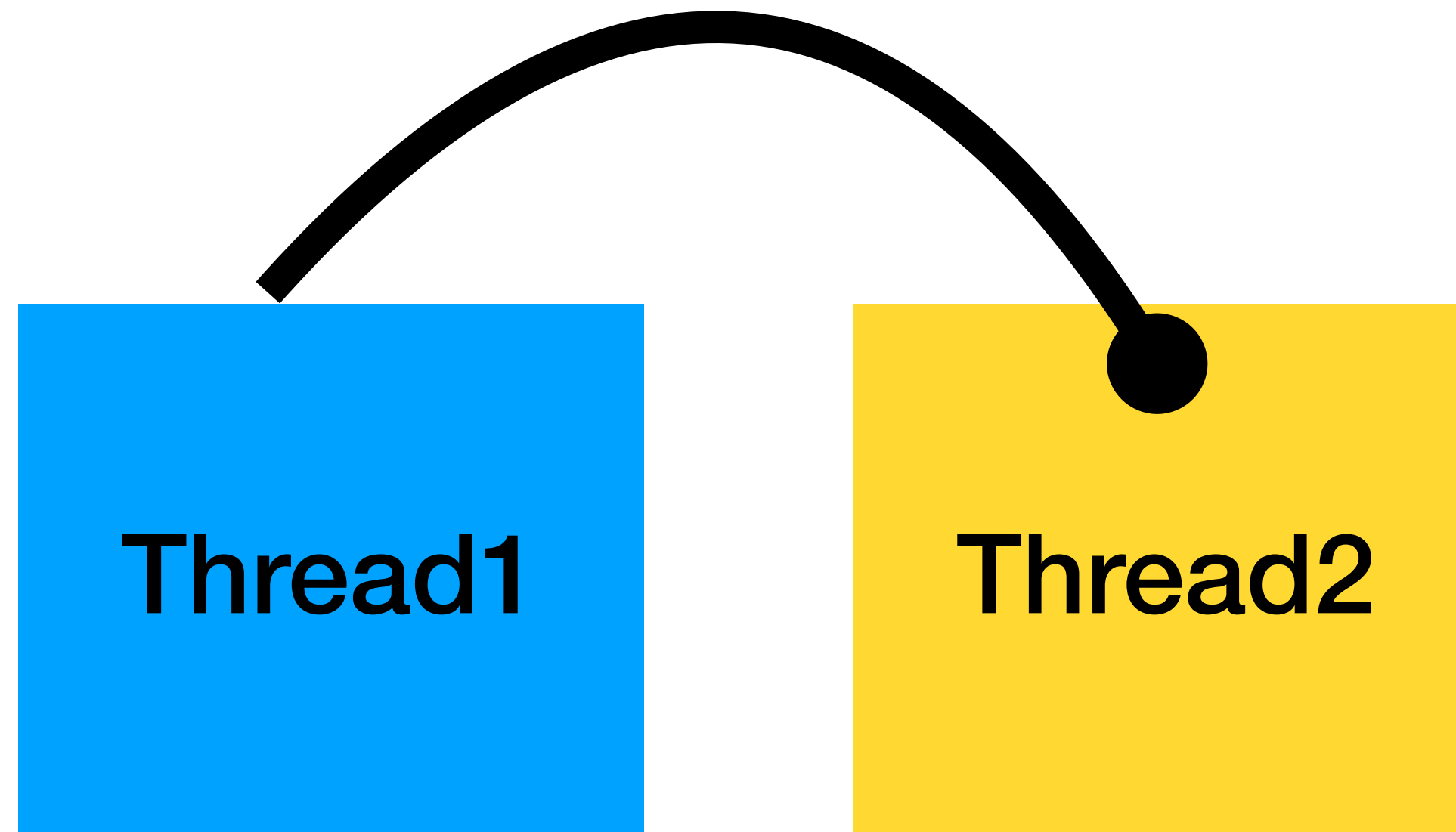
Scheduler inside OS

Background



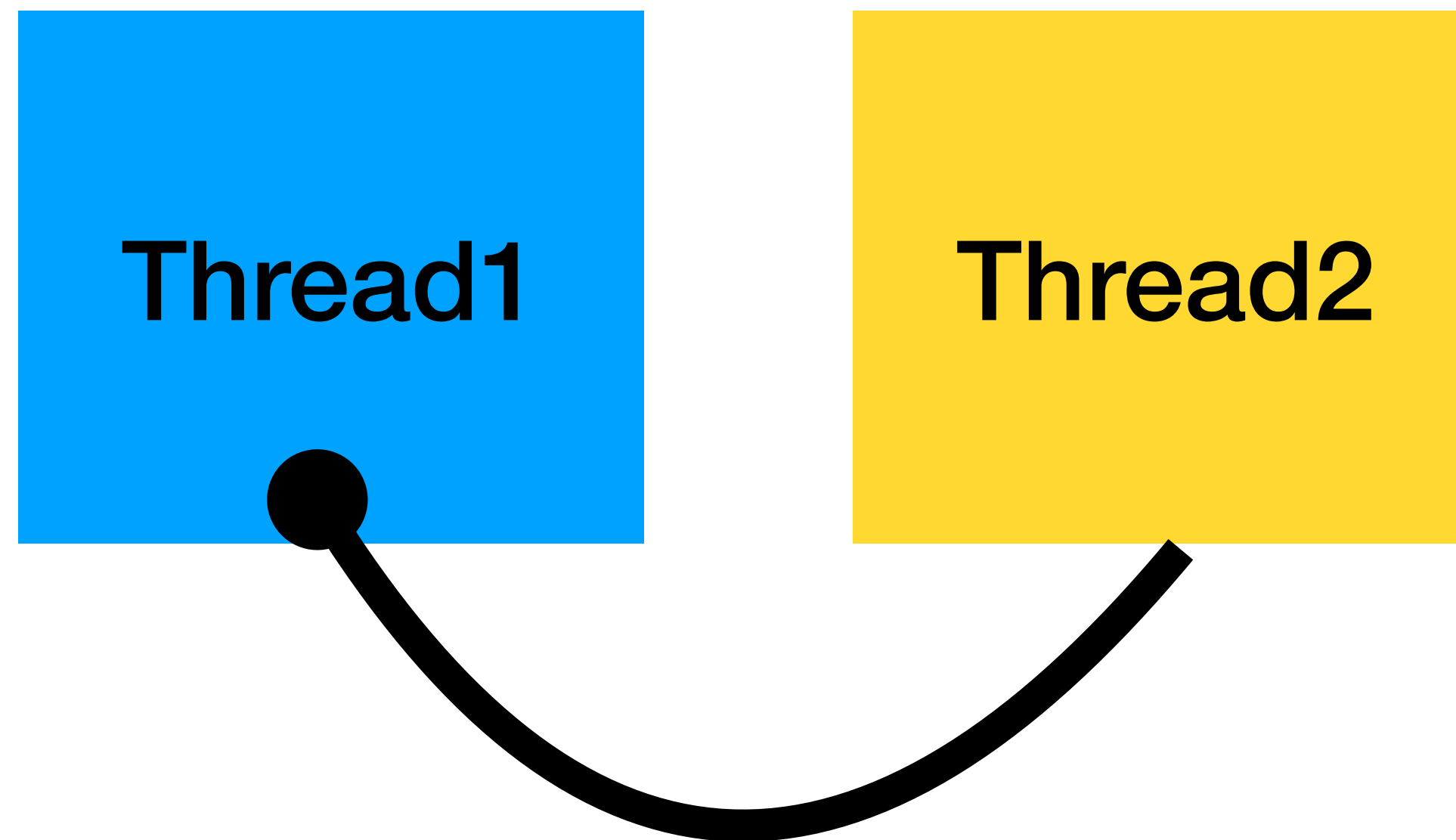
Effect handler

User-level scheduling using effect handler



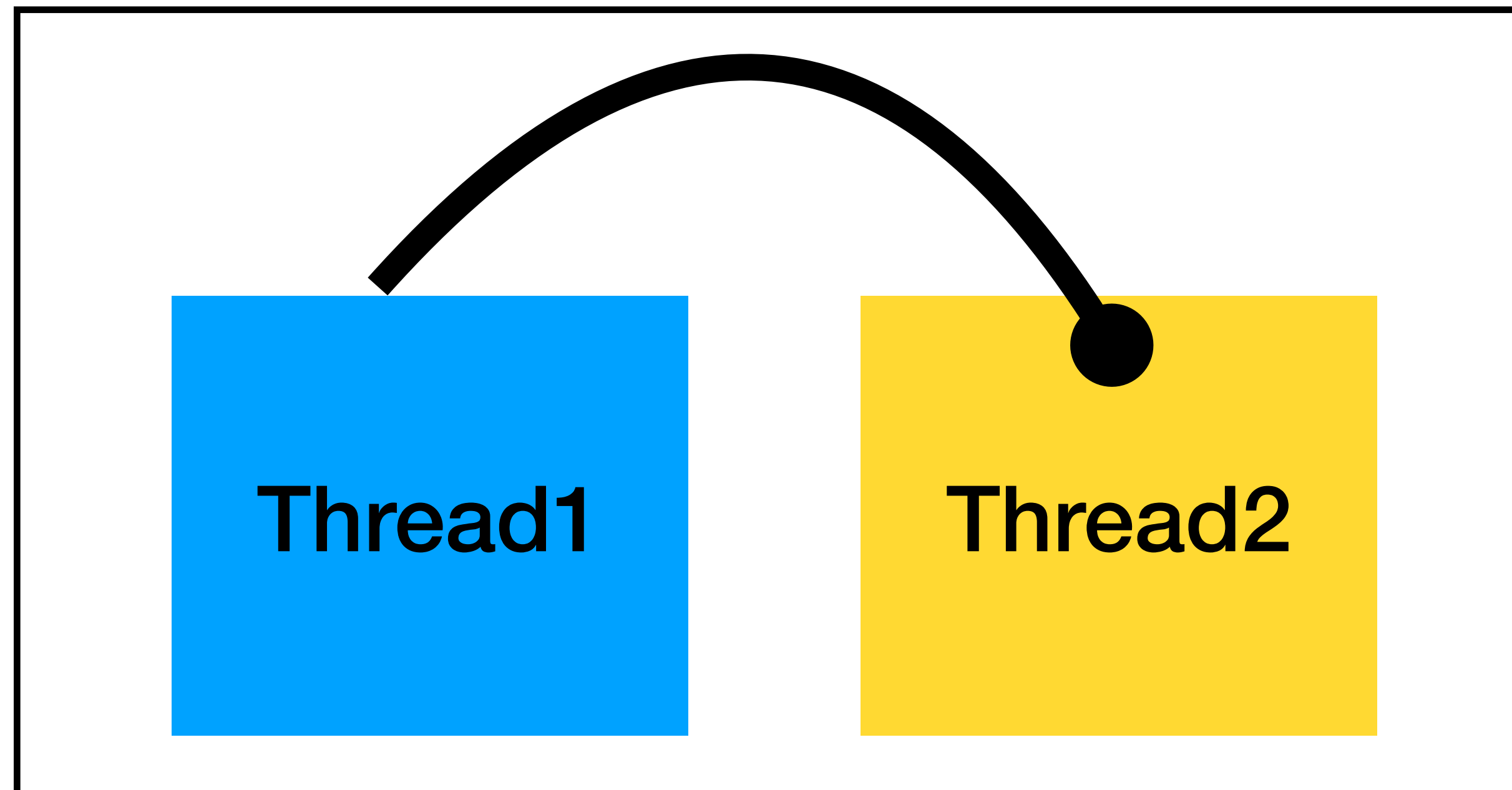
Effect handler

User-level scheduling using effect handler



Effect handler

User-level scheduling using effect handler

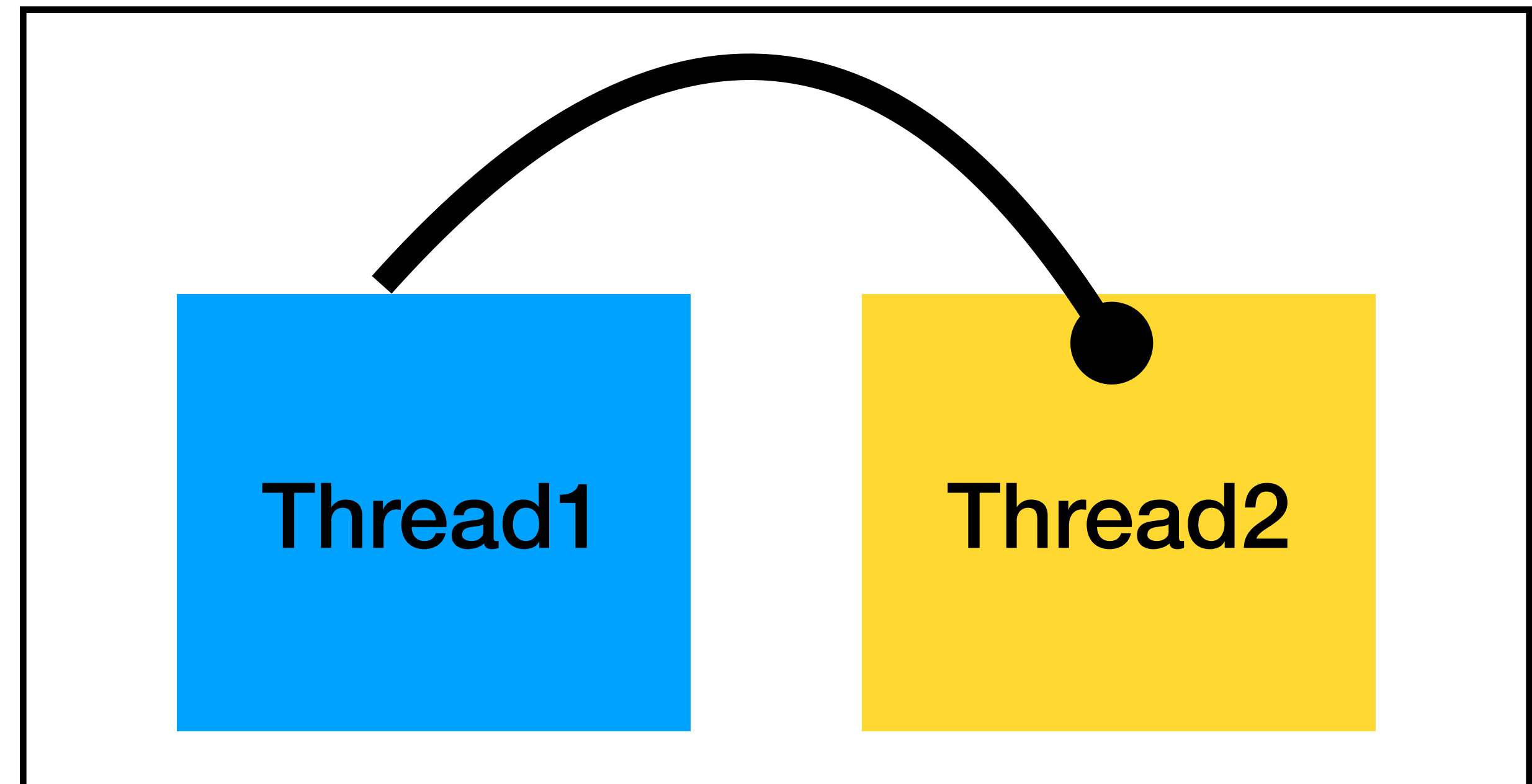


**Within the programming language,
without runtime support**

Effect handler

A powerful language feature

```
val parked = Null
def work() {
  handle
    while true
      DO_SOMETHING
      raise Yield()
  with Yield(k) =>
    val peer = parked
    parked = k
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      resume peer ()
}
work()
work()
```



Effect handler

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Common language features:

```
val
def
while
if..then..
```

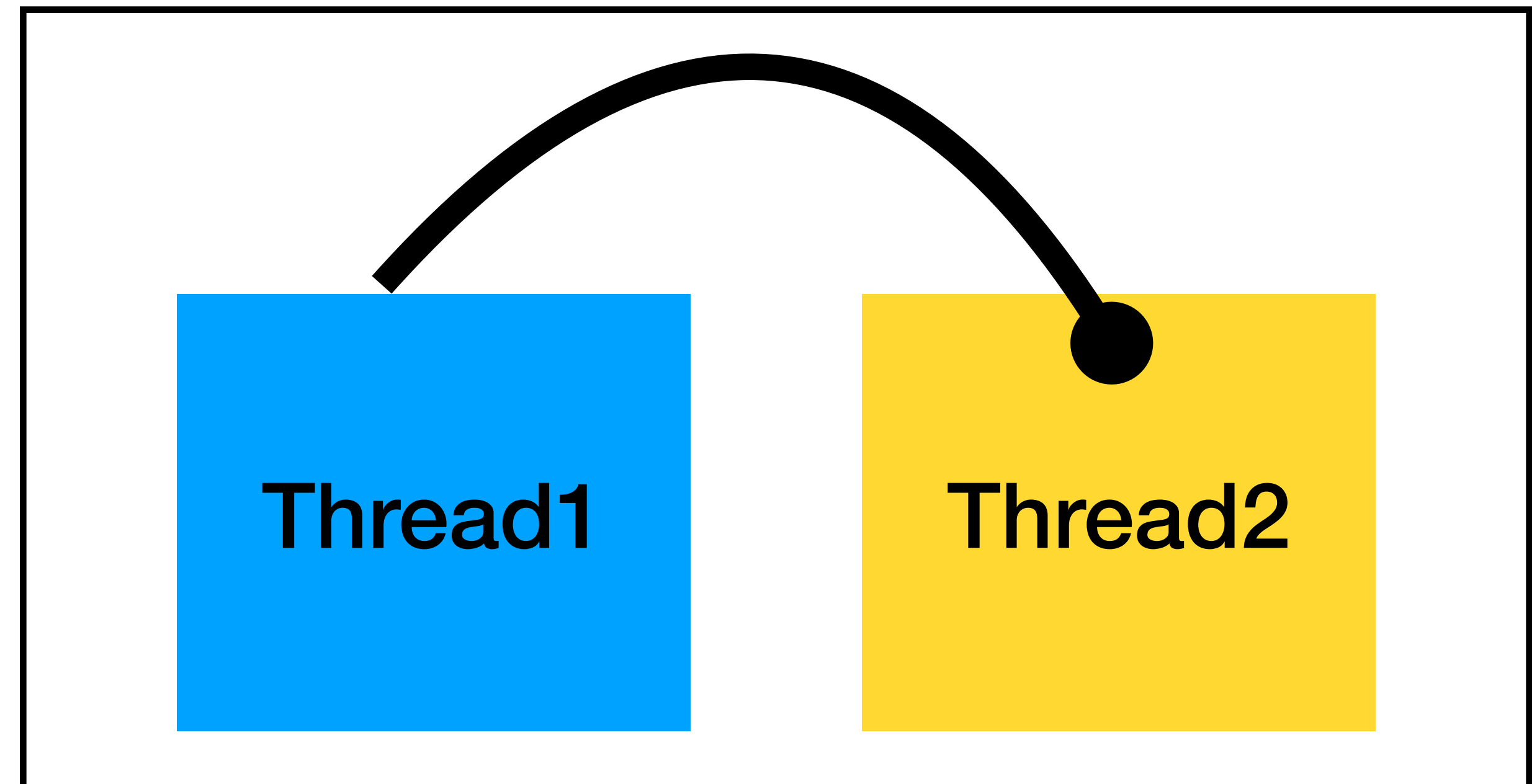
New language features:

```
handle..with..
raise
resume
```

Effect handler

A powerful language feature

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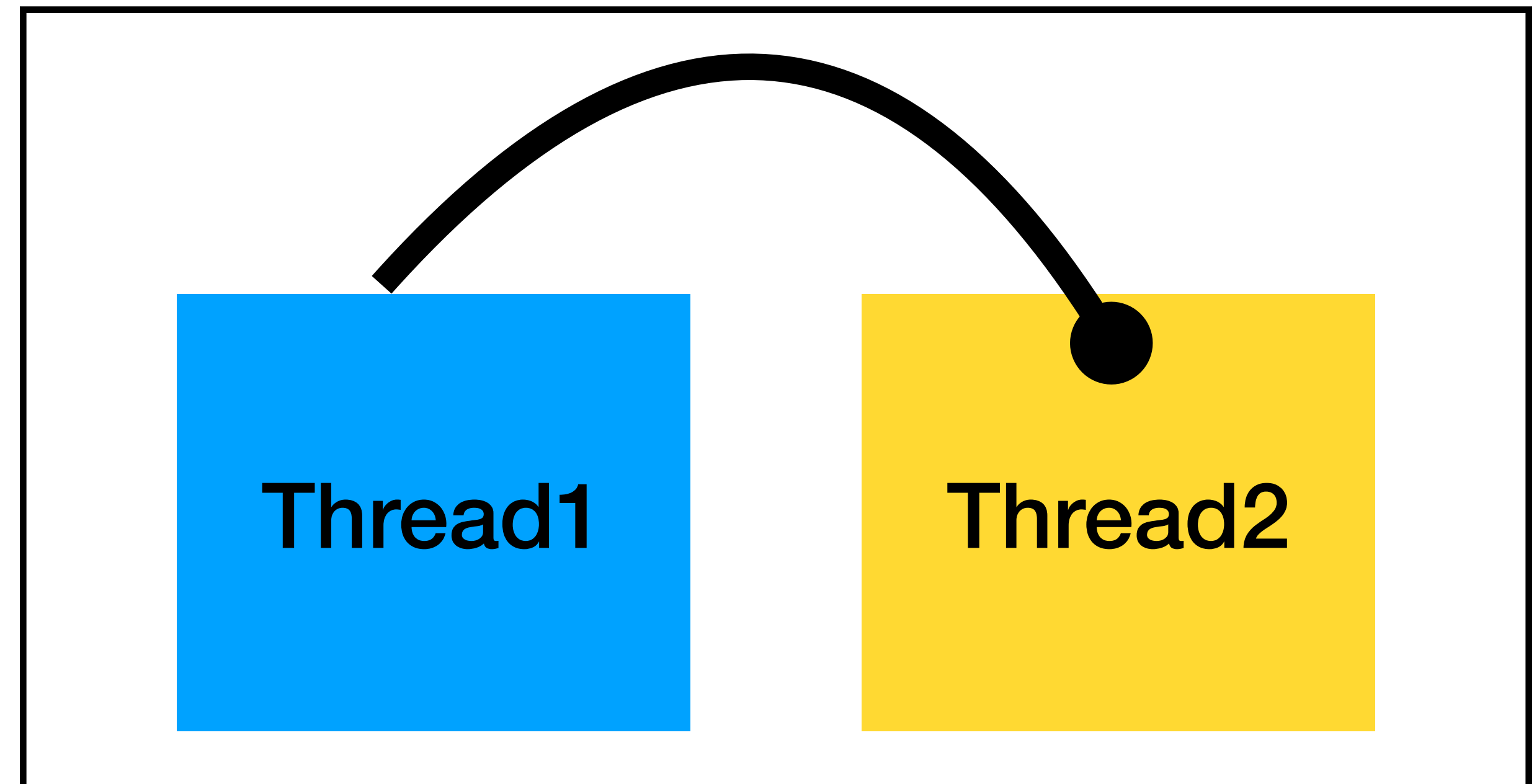
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```

```
work()
```

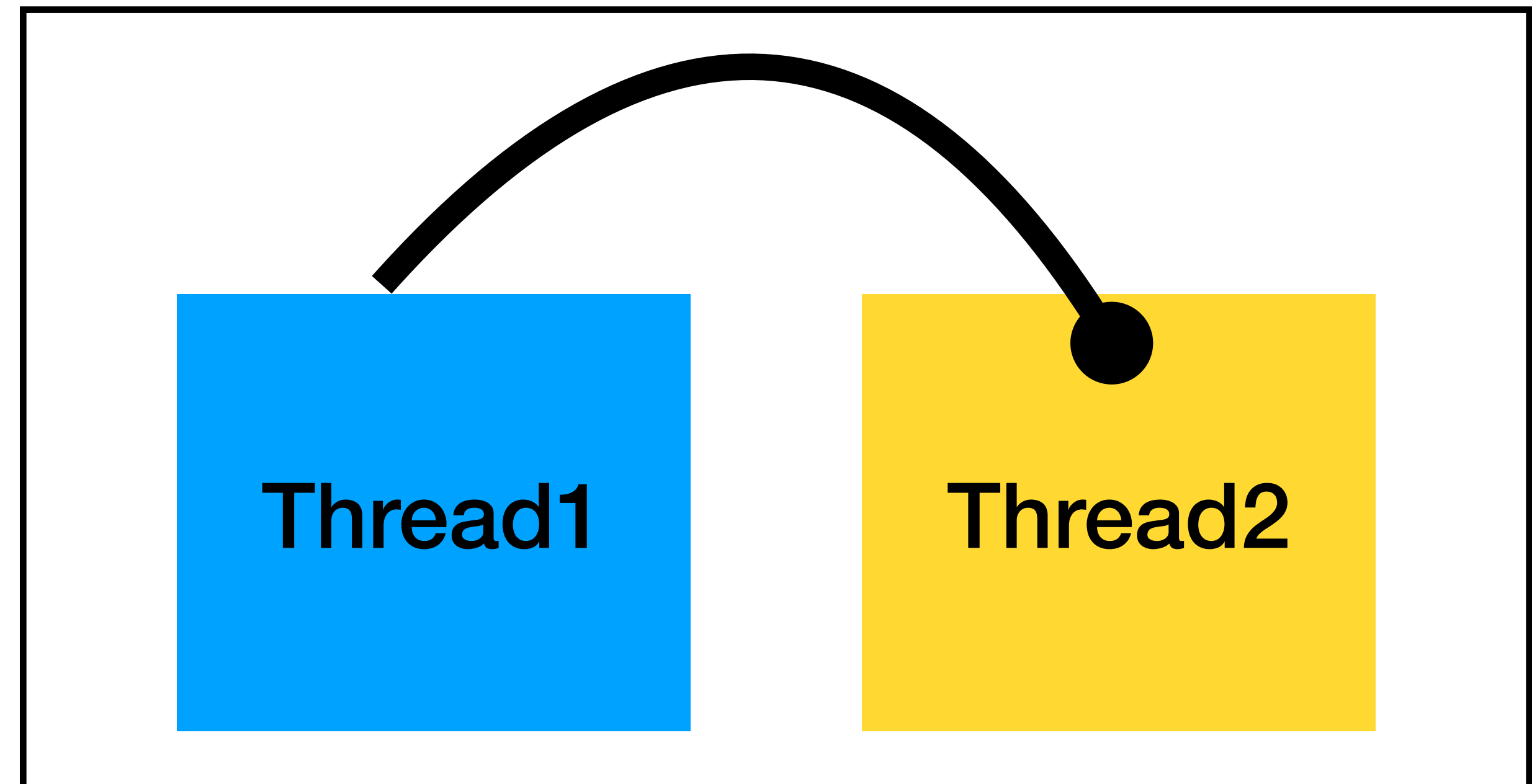
```
work()
```



Effect handler

A powerful language feature

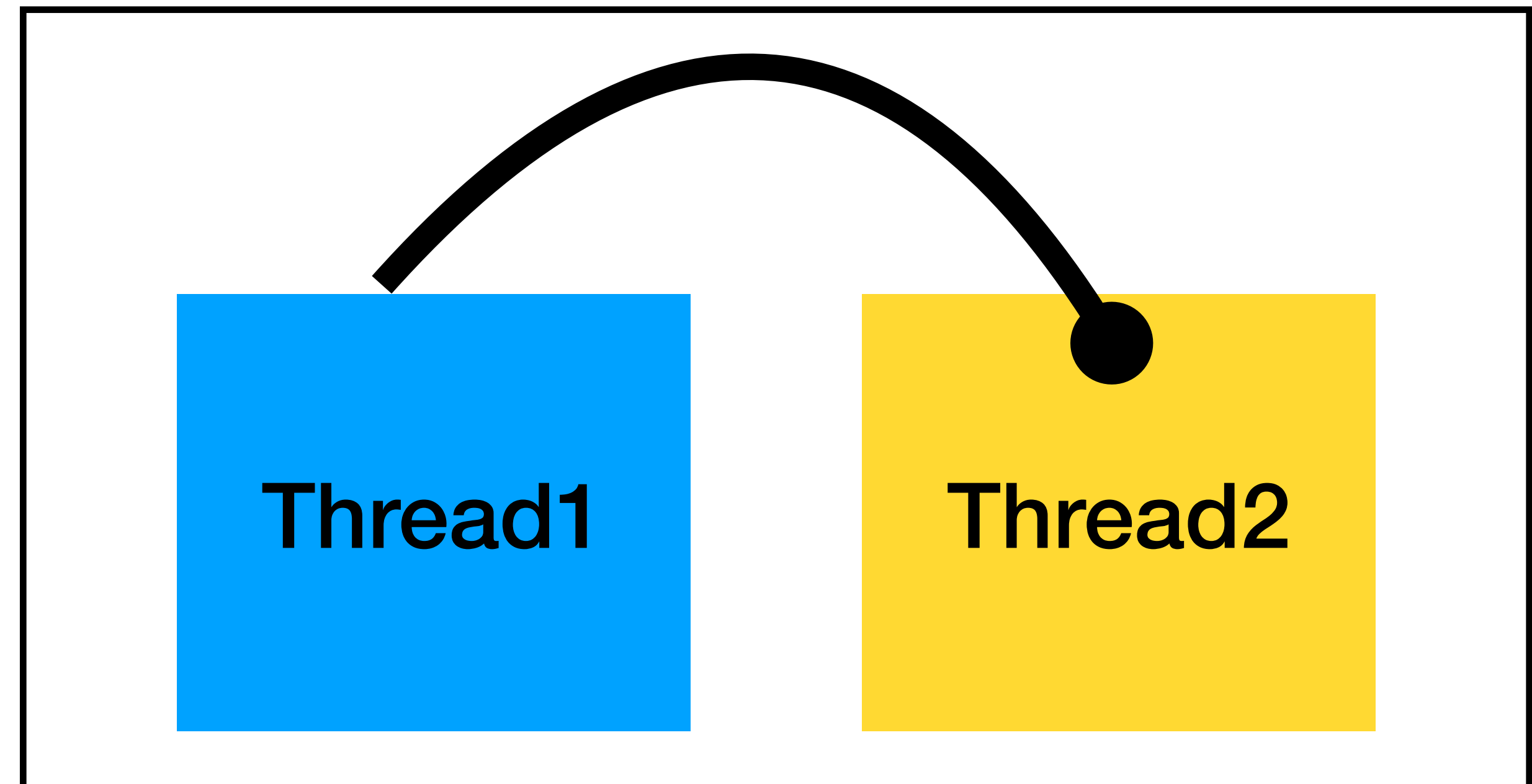
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Effect handler

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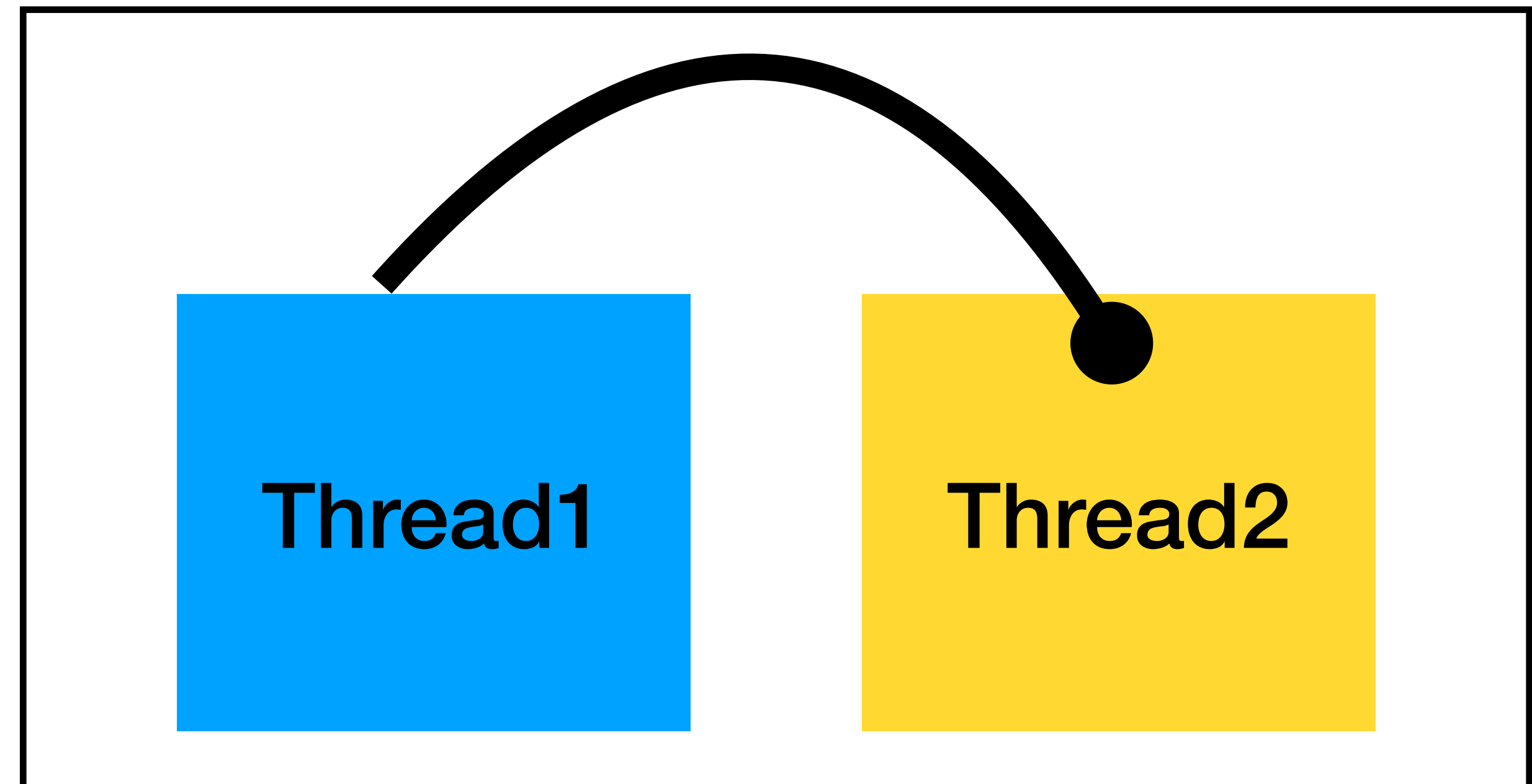
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Effect handler

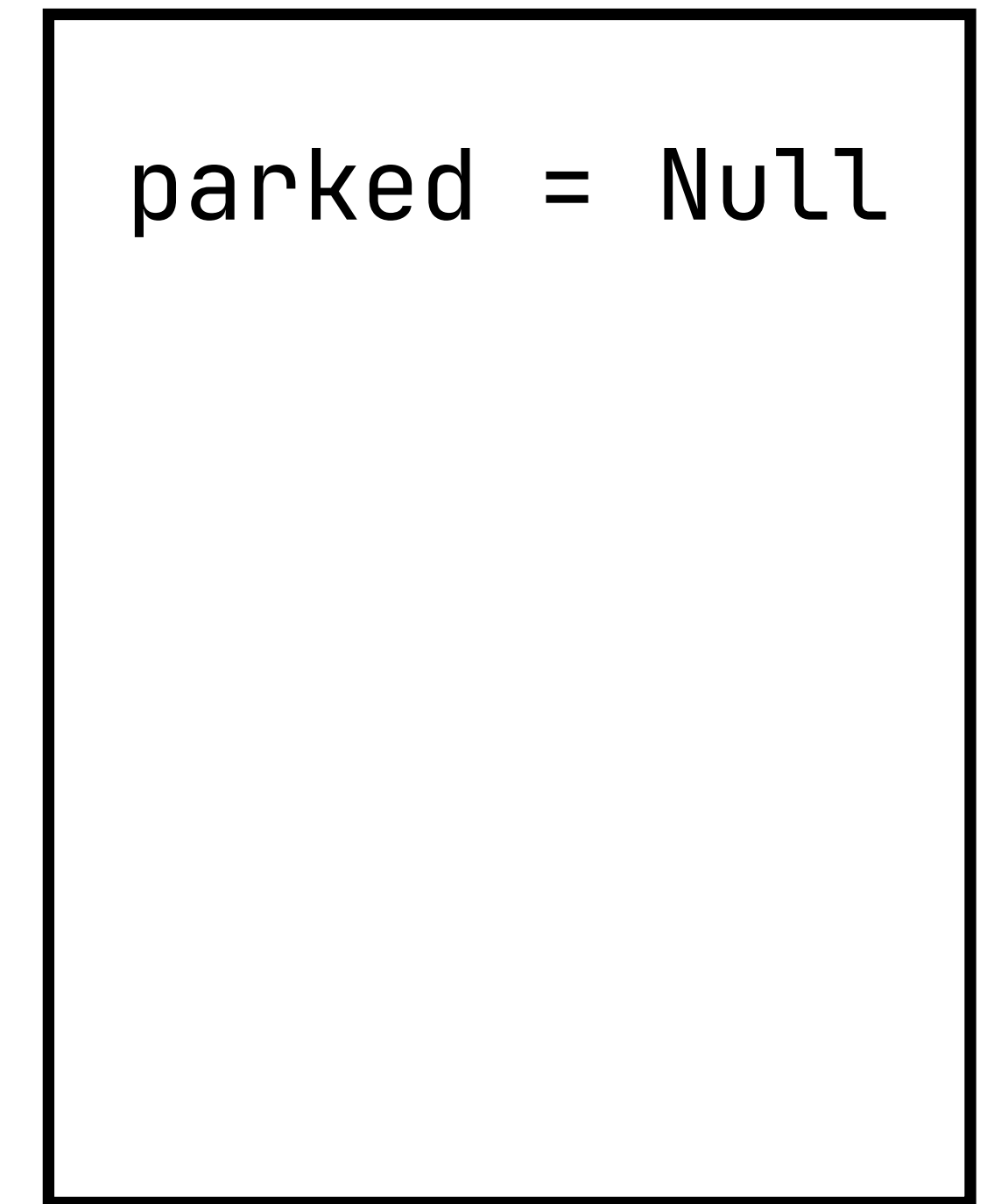
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work() ←
work()
```

Processor:



Memory:



Effect handler

A powerful language feature

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work()
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Processor:

```
var1 = 1
var2 = 2
...
```

Memory:

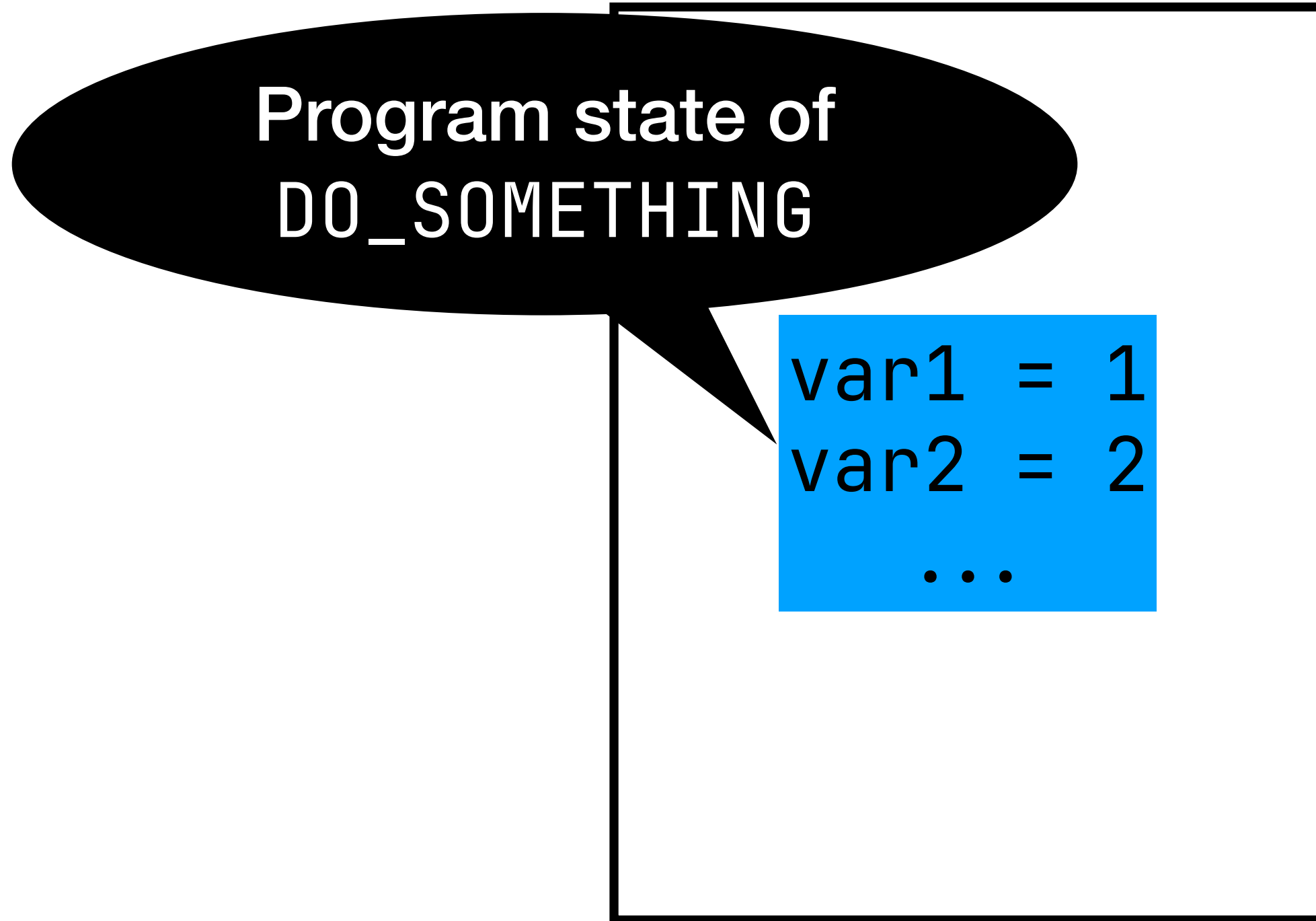
```
parked = Null
```

Effect handler

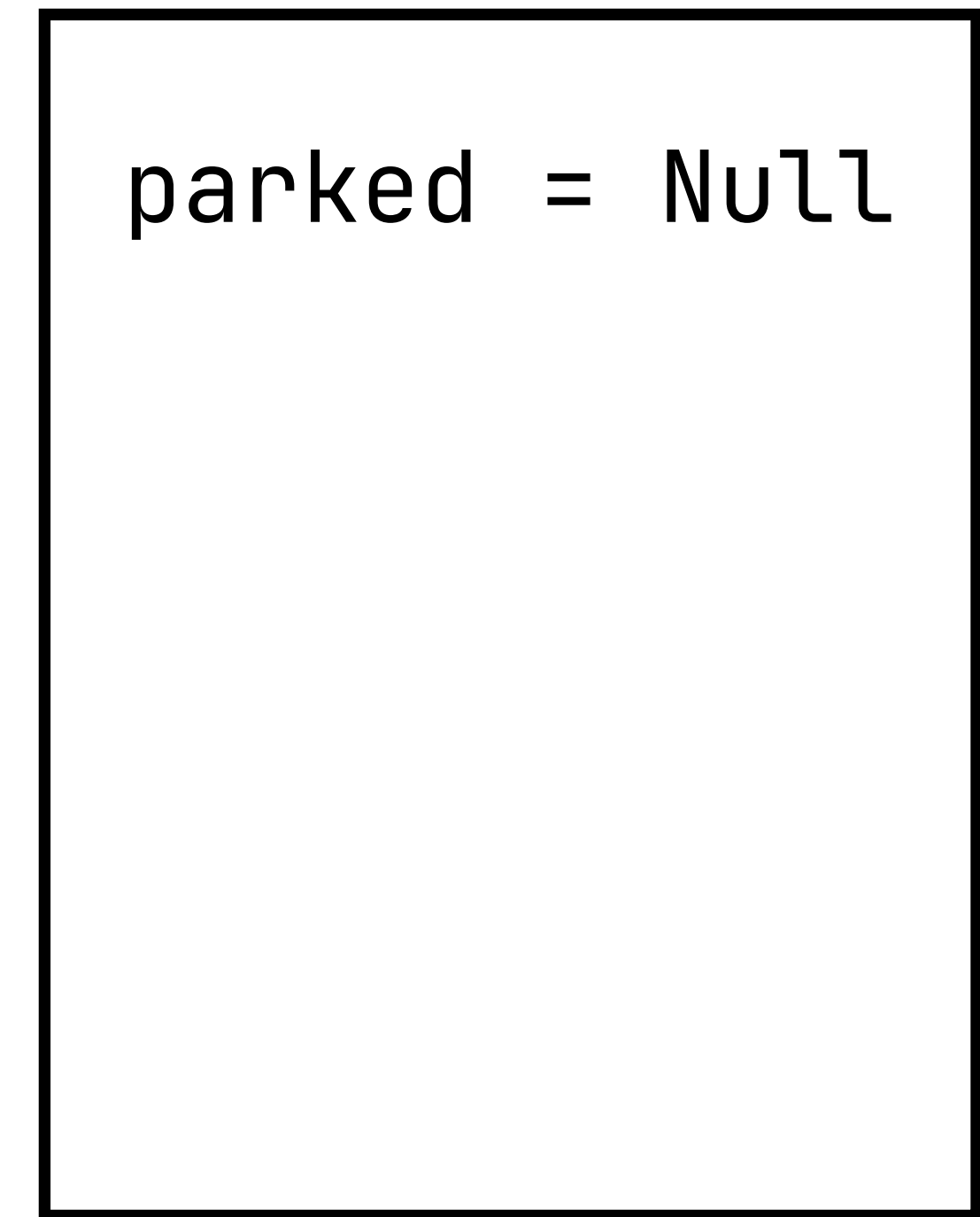
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Processor:

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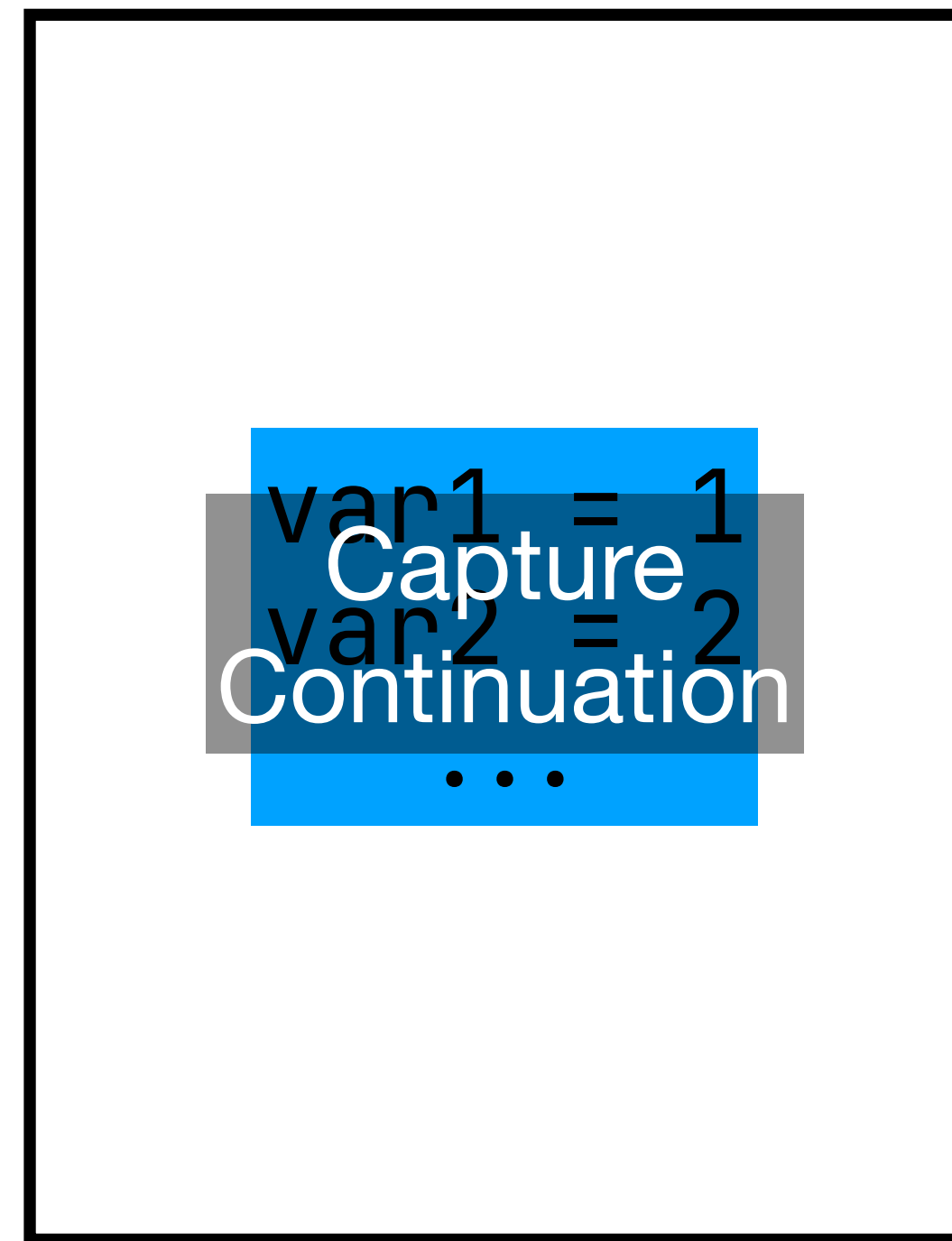
```
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Effect handler

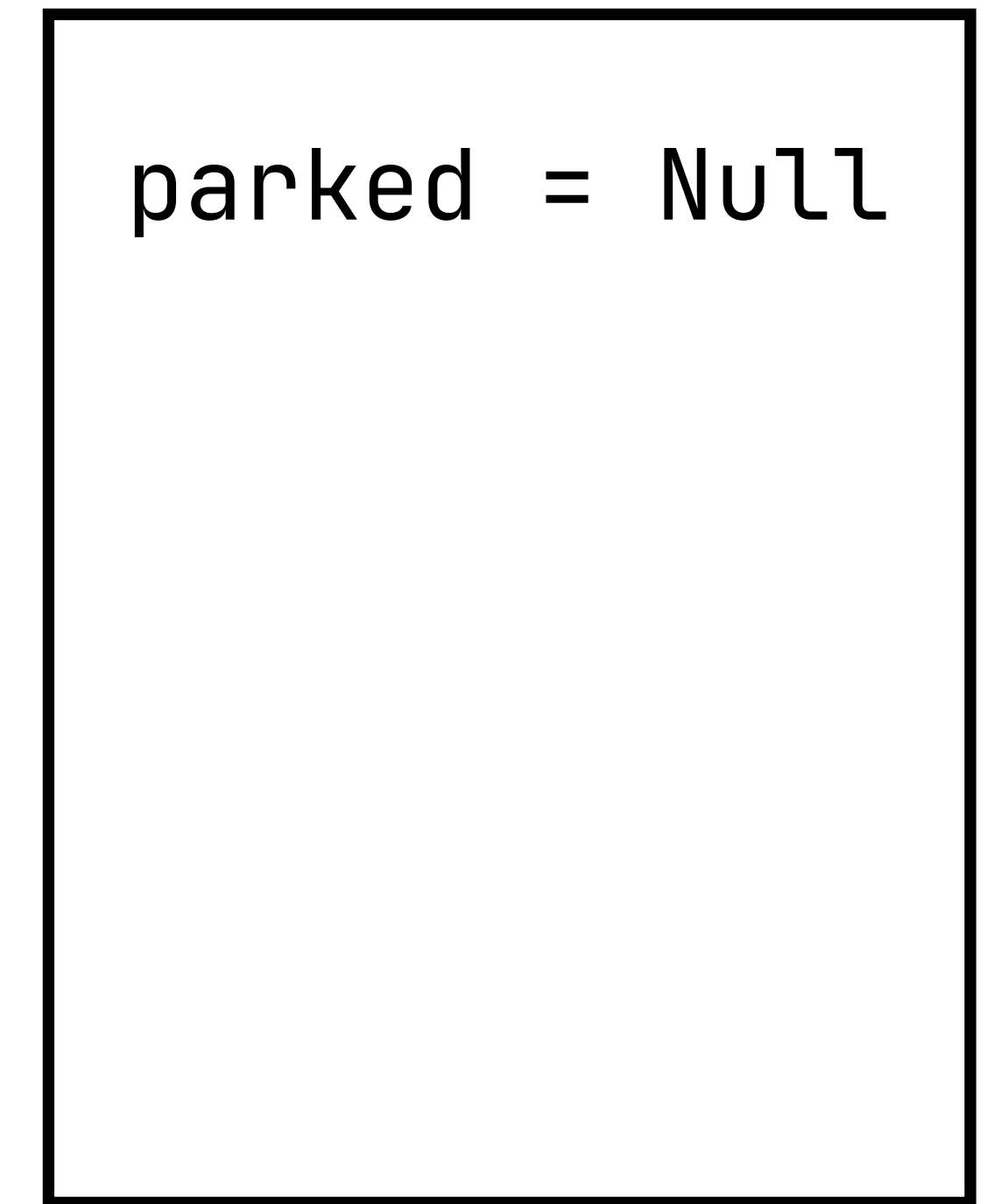
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Memory:

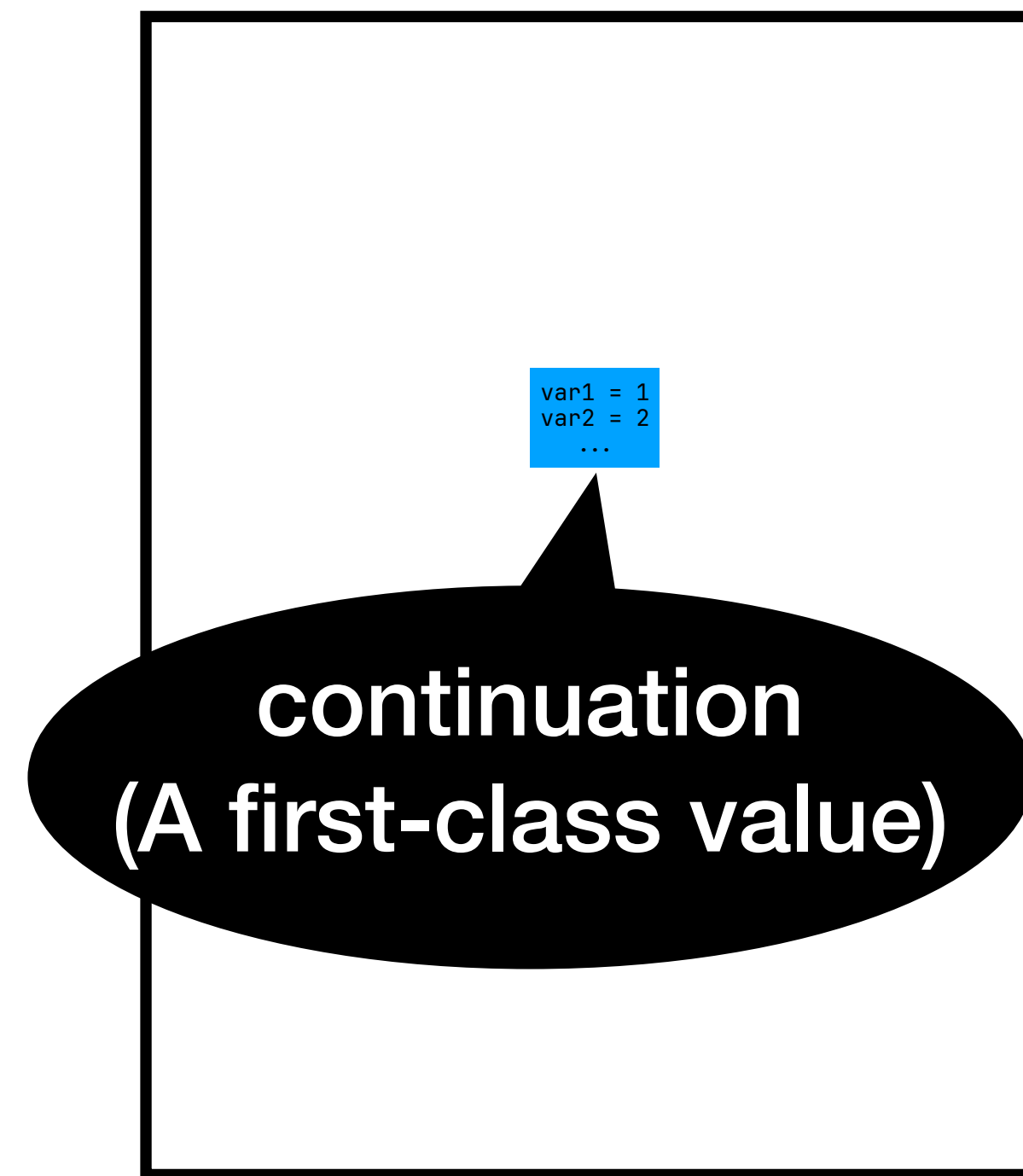


Effect handler

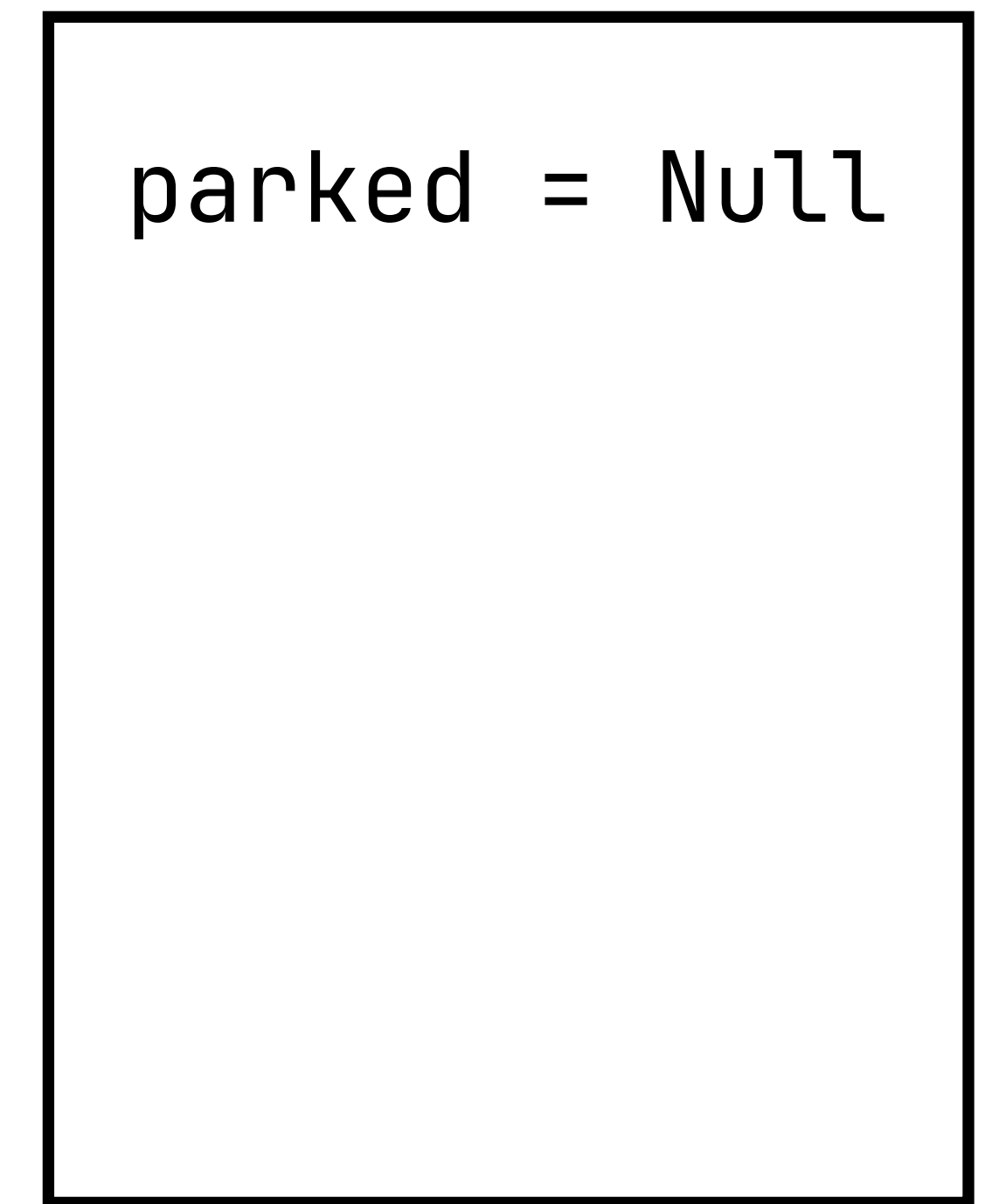
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Processor:



Memory:

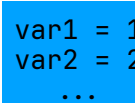


Effect handler

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```

Processor:

k = 

Memory:

parked = Null

Effect handler

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```

Processor:

```
k = var1 = 1  
var2 = 2  
...  
peer = Null
```

Memory:

```
parked = Null
```

Effect handler

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...  
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```

Memory:

```
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var2 = 2  
...
```

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Memory:

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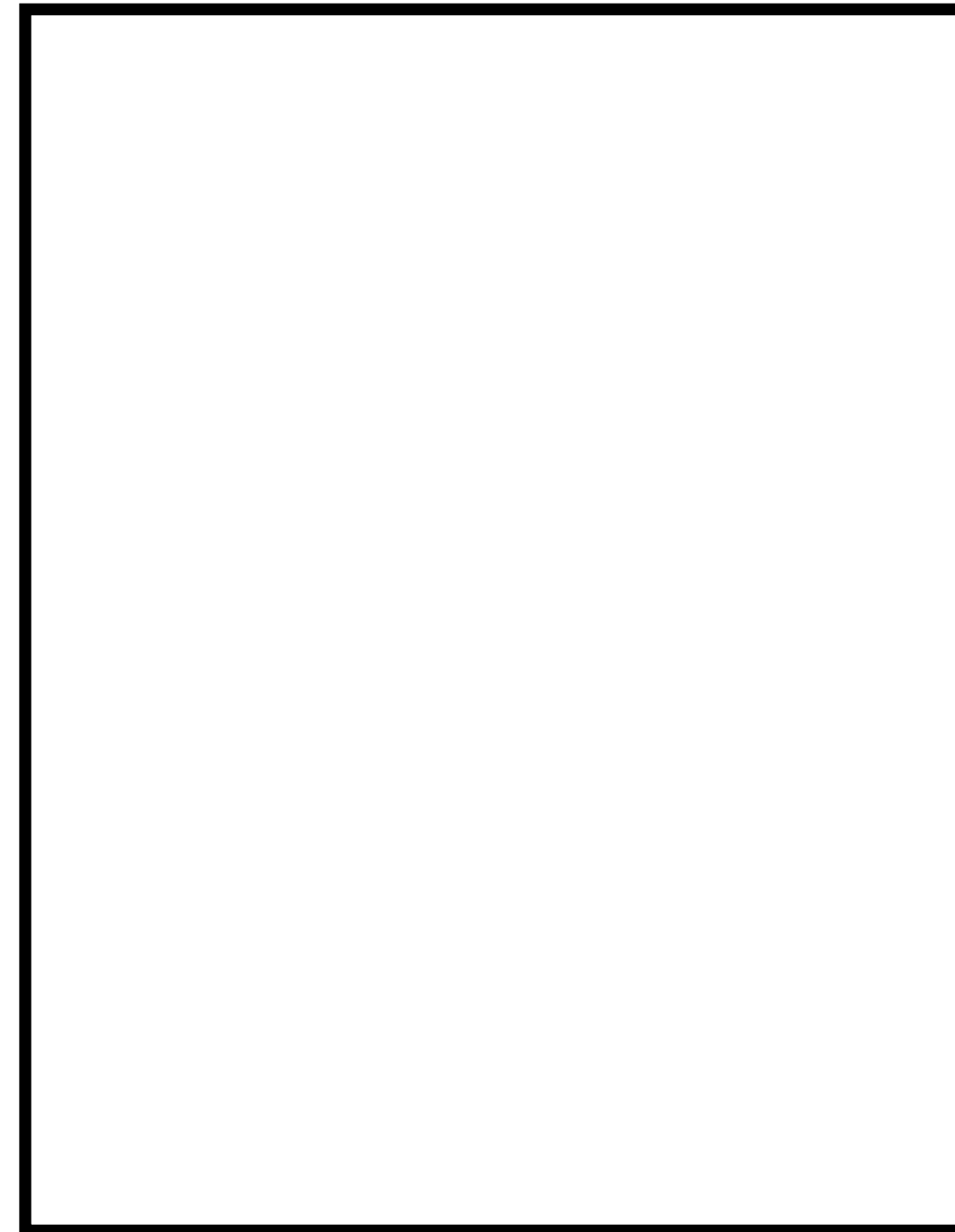
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Processor:



Memory:

parked = 

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Memory:

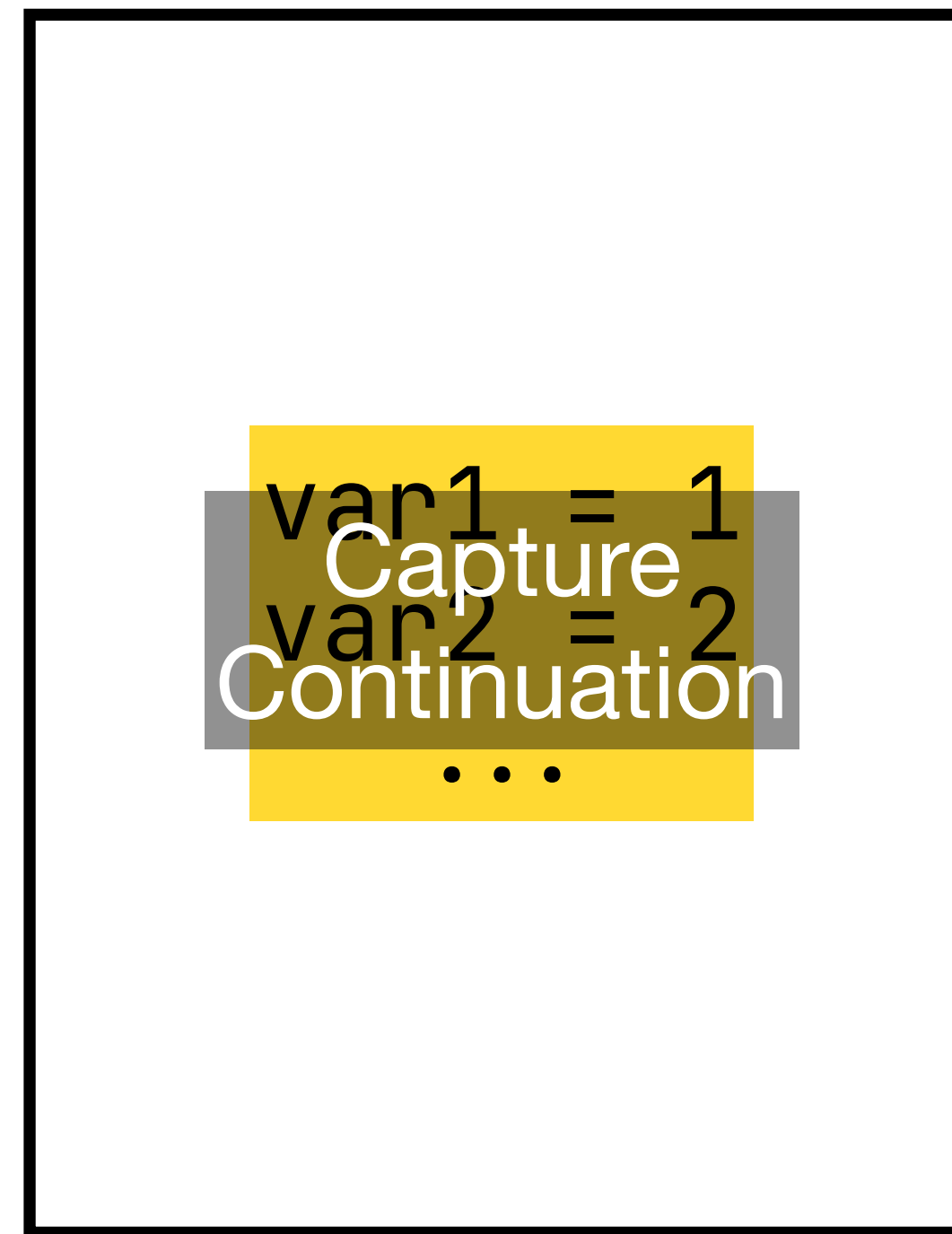
```
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```

Effect handler

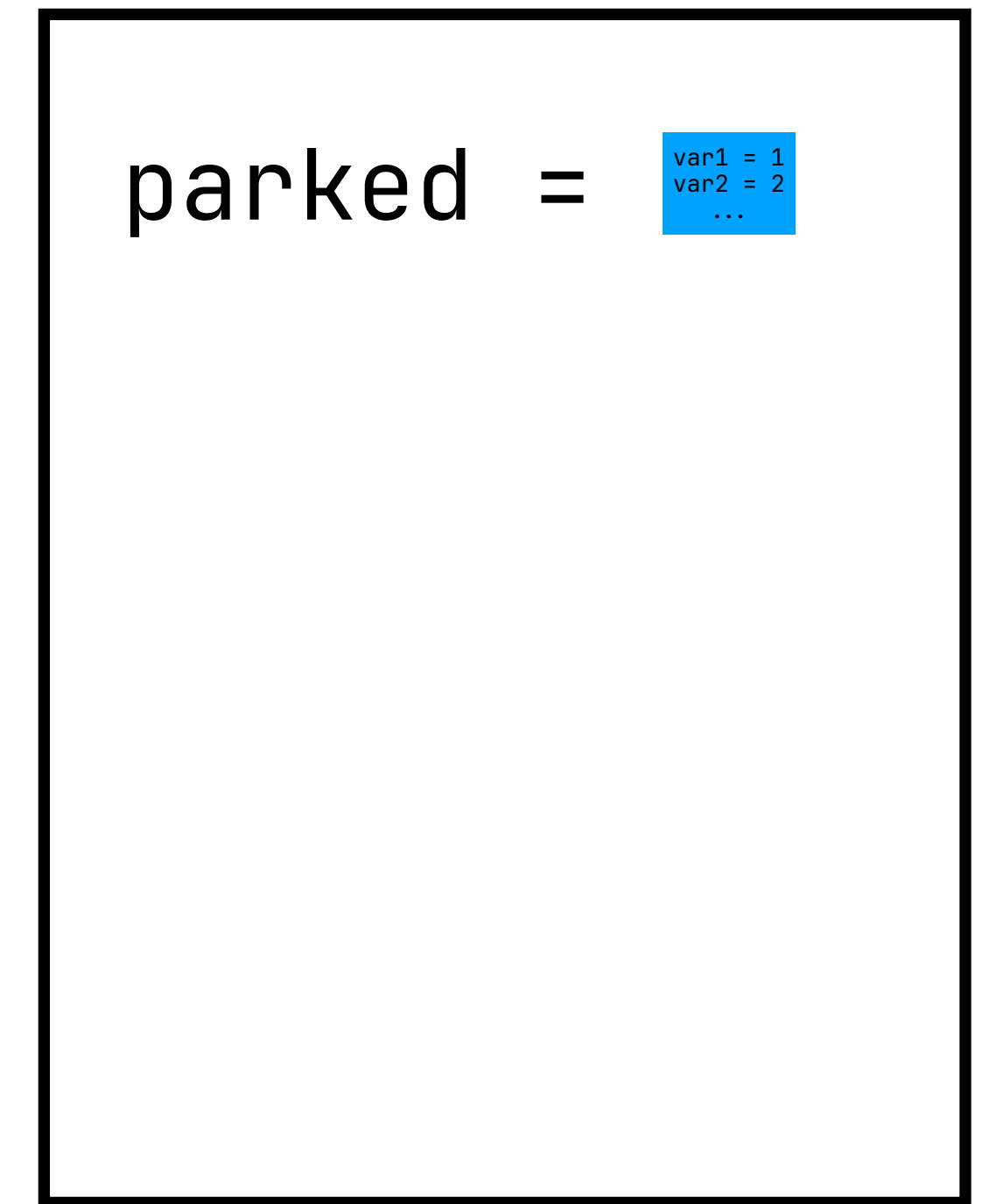
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Memory:



Effect handler

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Processor:

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Memory:

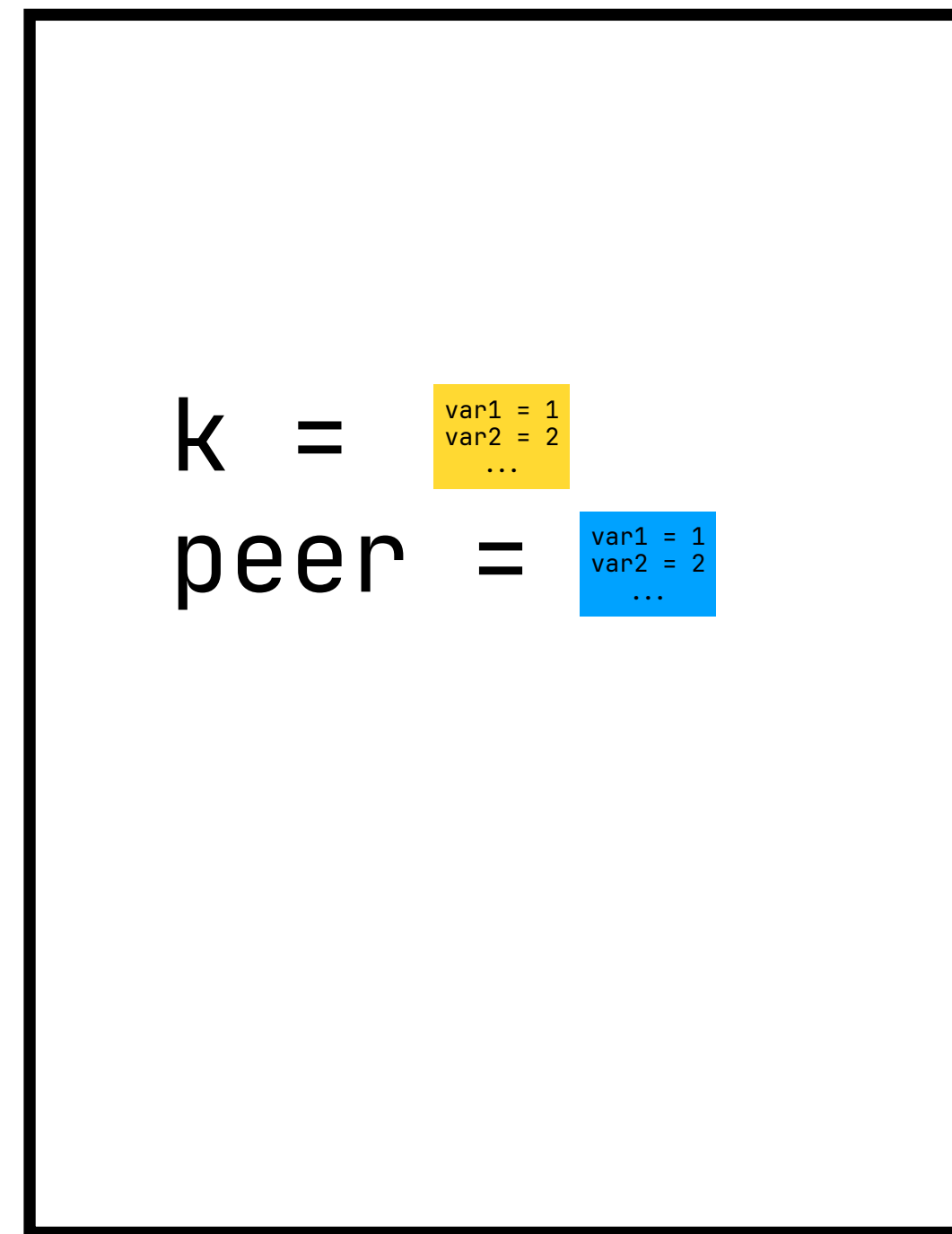
parked =

Effect handler

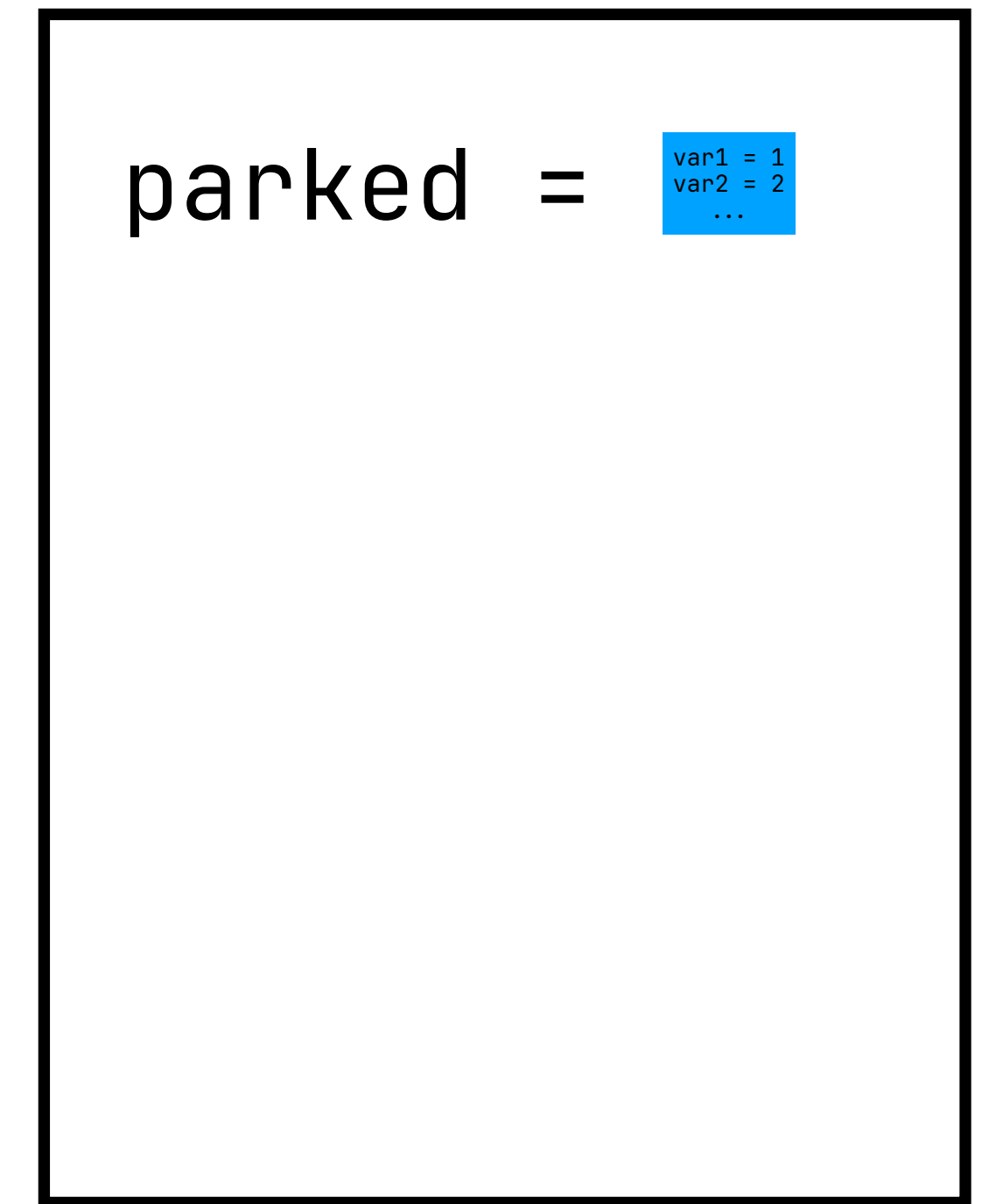
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Processor:



Memory:



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Processor:

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...  
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...
```

Memory:

```
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...
```

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...
```

Memory:

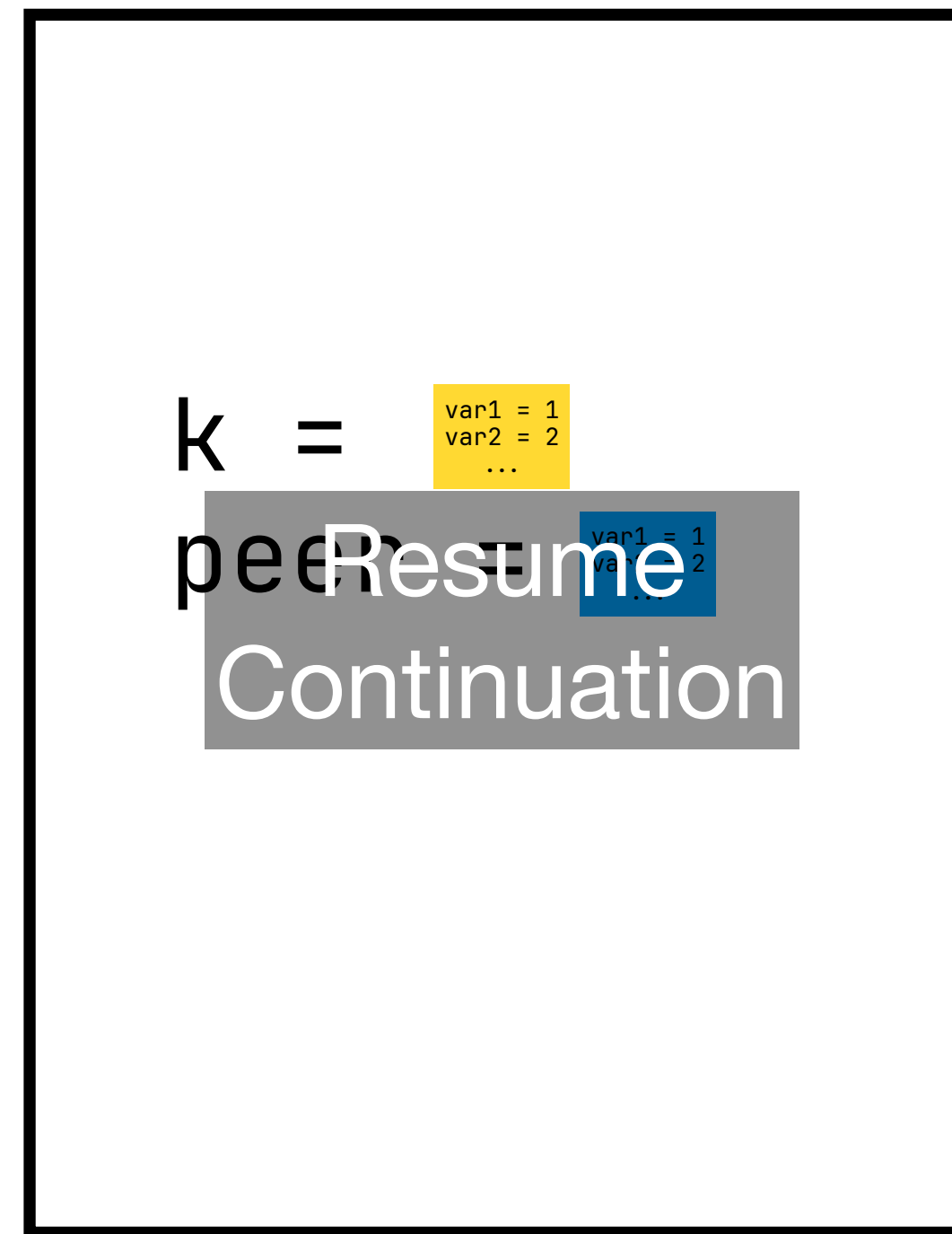
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Effect handler

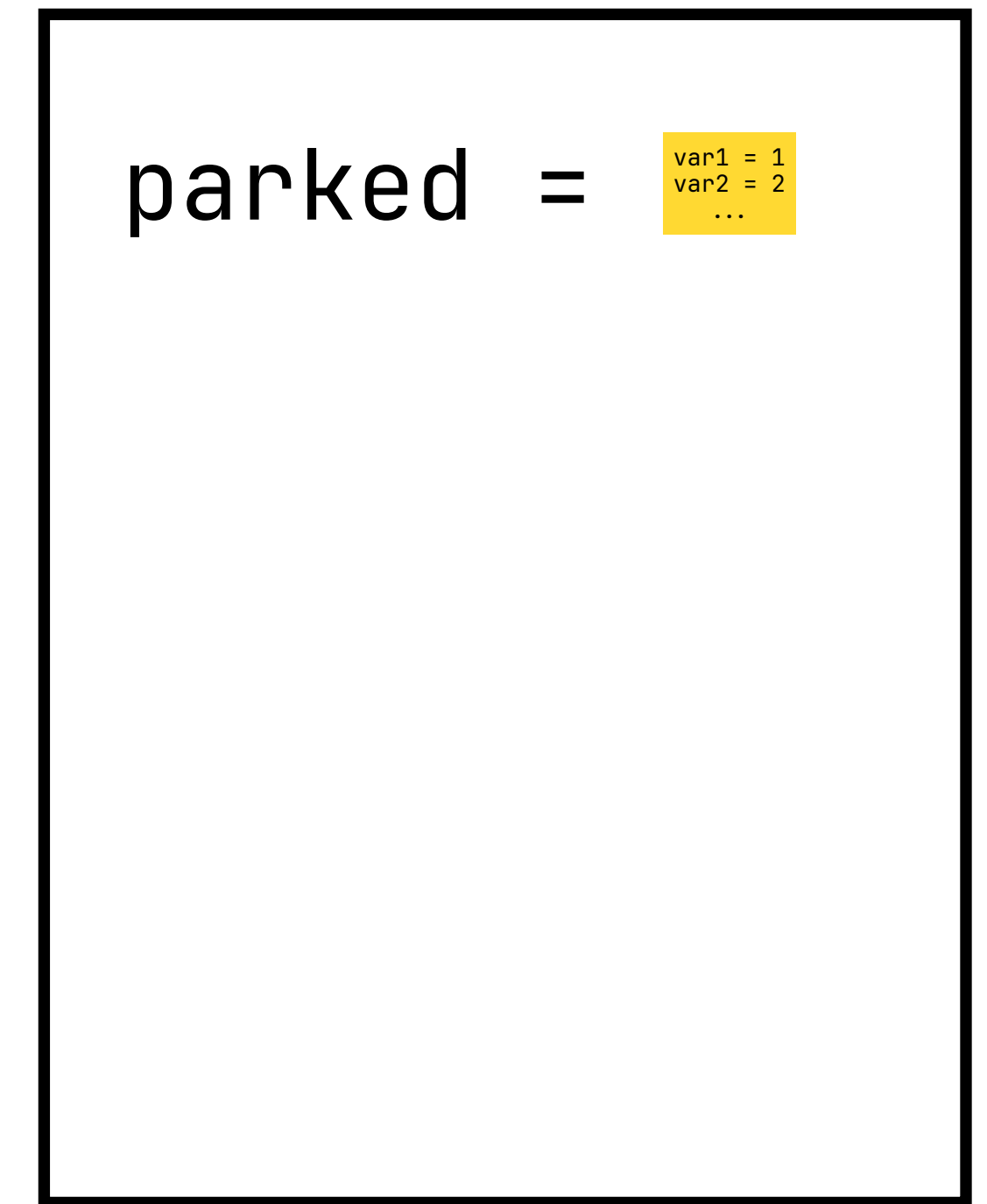
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Processor:

```
var1 = 1
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```

Memory:

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parked = var1 = 1
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Effect handler

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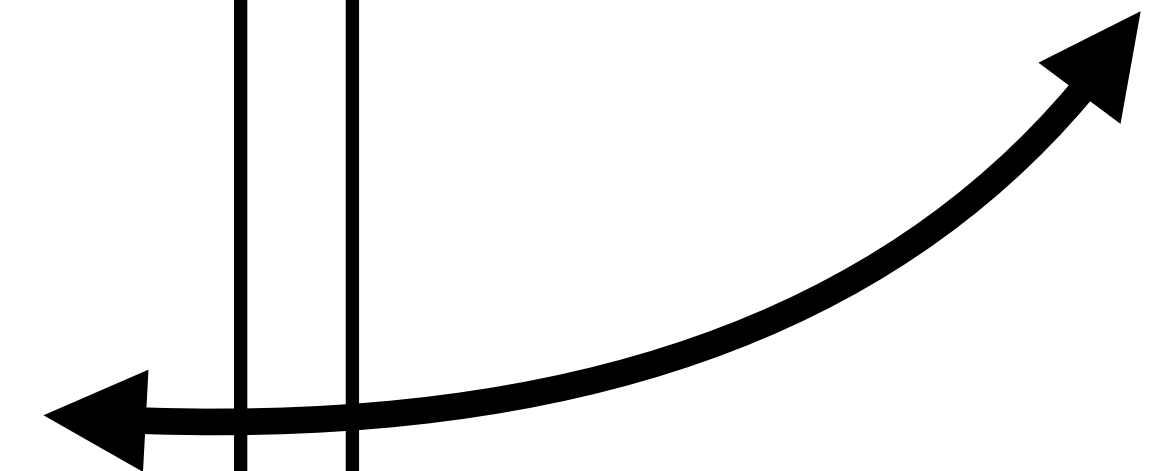
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Effect handler

A powerful language feature

Language Definition

Syntax

+

Semantics

Effect handler

A powerful language feature

Language Definition

Syntax

+

Semantics

`handle E with H`

`raise ...`

`resume ...`

Capture continuation

Resume continuation

Effect handler

A powerful language feature

Language Definition

Syntax

+

Semantics

handle E with H

active research

Install a handler

raise ...

Find handler; Capture continuation

resume ...

Resume continuation

Effect handler

A powerful language feature

Language Definition

Syntax

+

Semantics

handle E with H

active research

Install a handler

raise ...

Find handler; Capture continuation

resume ...

Resume continuation

Introducing lexical effect handler

Lexical effect handler

Last bit of semantics: install a handler

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work()
work()
```

Lexical effect handler

Last bit of semantics: install a handler

```
handle
  while true
    DO_SOMETHING
    raise yield()
with yield(k) ⇒
  SCHEDULER
```

Lexical effect handler

Last bit of semantics: install a handler

```
handle
  while true
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variable

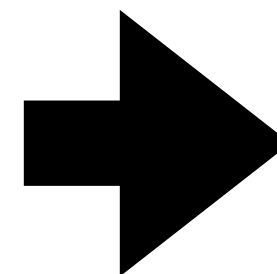
Lexical effect handler

Last bit of semantics: install a handler

```
handle
  while true
    DO_SOMETHING
    raise yield()
with yield(k) ⇒
  SCHEDULER
```

variable

reduce to



```
handle
  while true
    DO_SOMETHING
    raise #314()
with #314(k) ⇒
  SCHEDULER
```

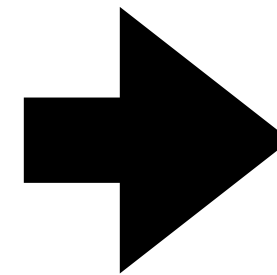
fresh name

Lexical effect handler

Last bit of semantics: install a handler

```
handle
  while true
    DO_SOMETHING
    raise yield()
with yield(k) ⇒
  SCHEDULER
```

reduce to



```
handle
  while true
    DO_SOMETHING
    raise #314()
with #314(k) ⇒
  SCHEDULER
```

handler
frame

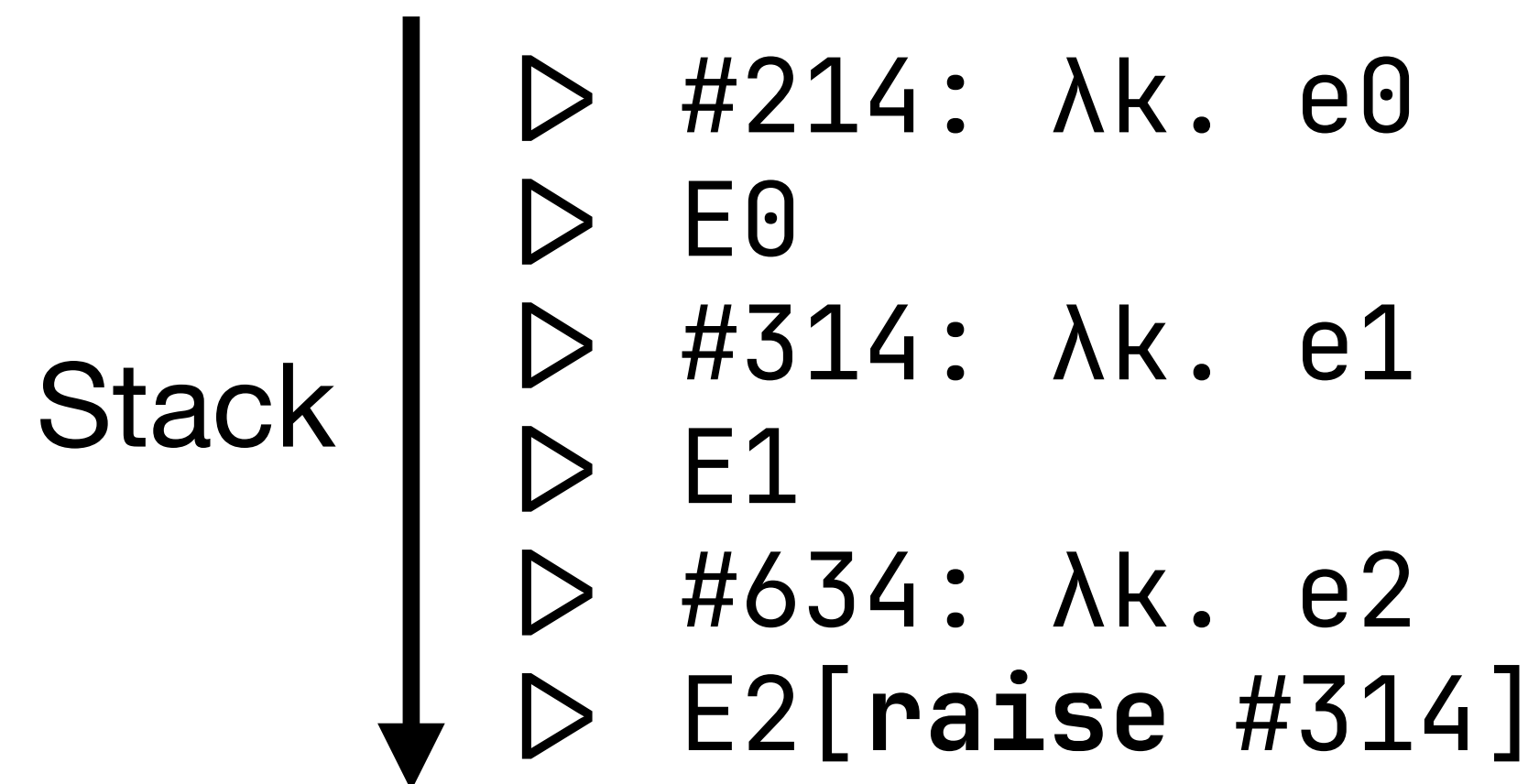
function
frame

Another presentation:

```
▷ #314: SCHEDULER
▷ while true {...raise #314()}
```

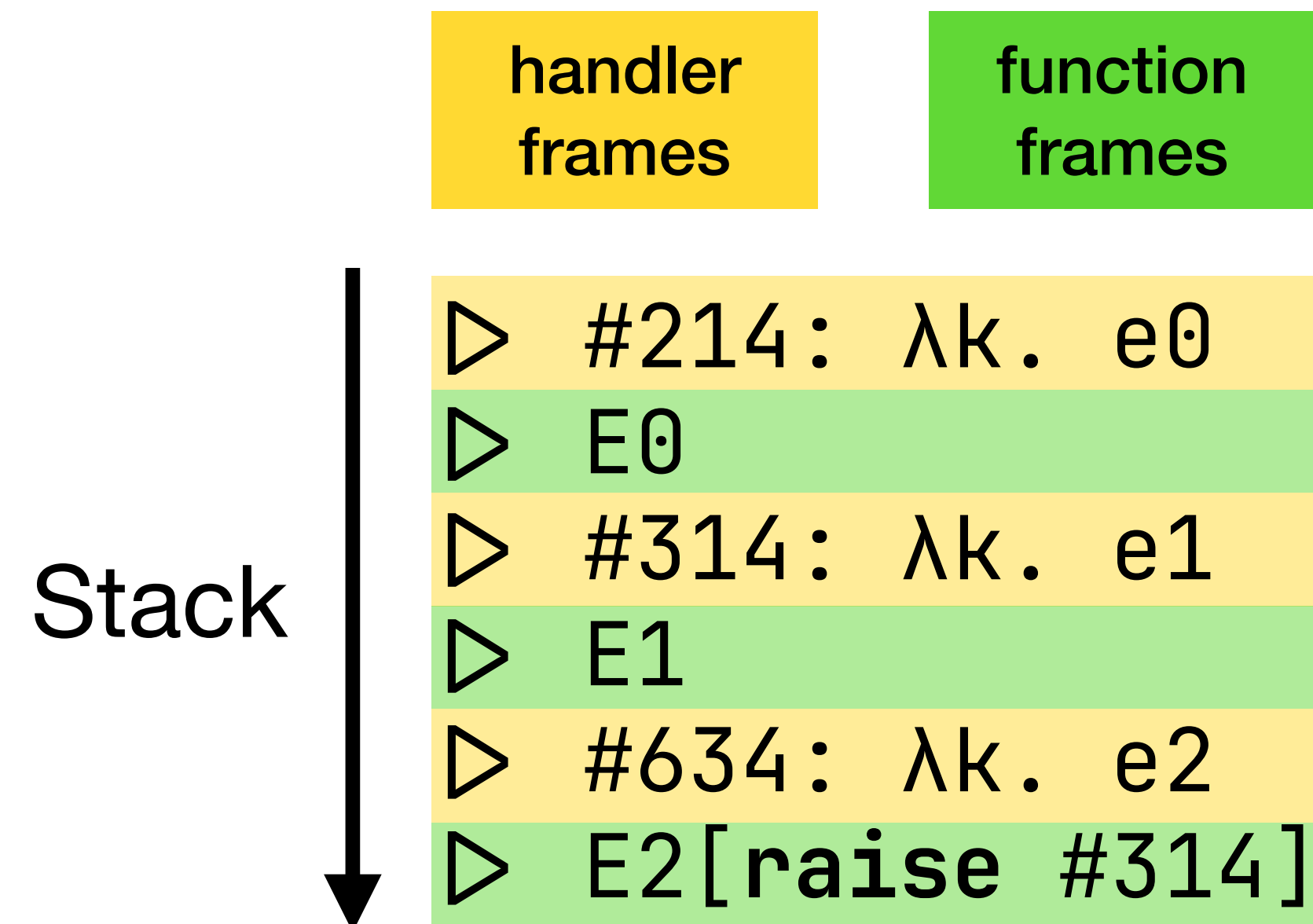
Lexical effect handler

Last bit of semantics: install a handler



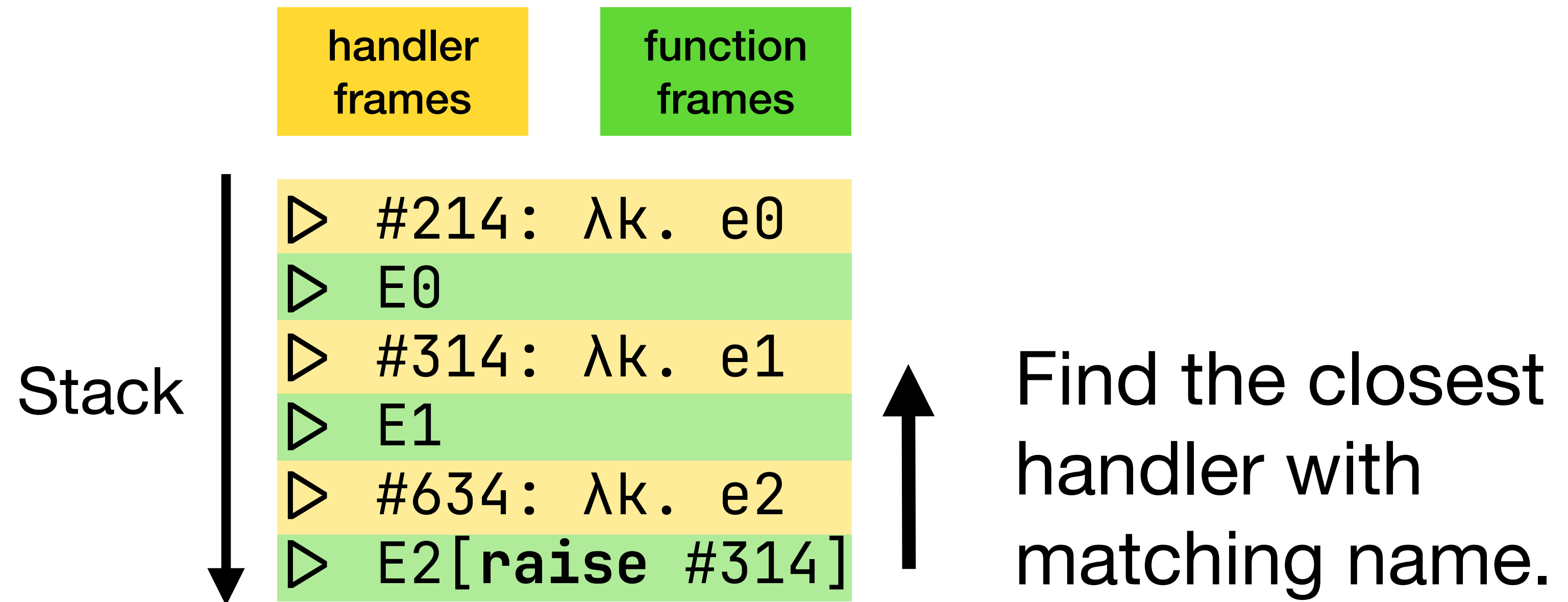
Lexical effect handler

Last bit of semantics: install a handler



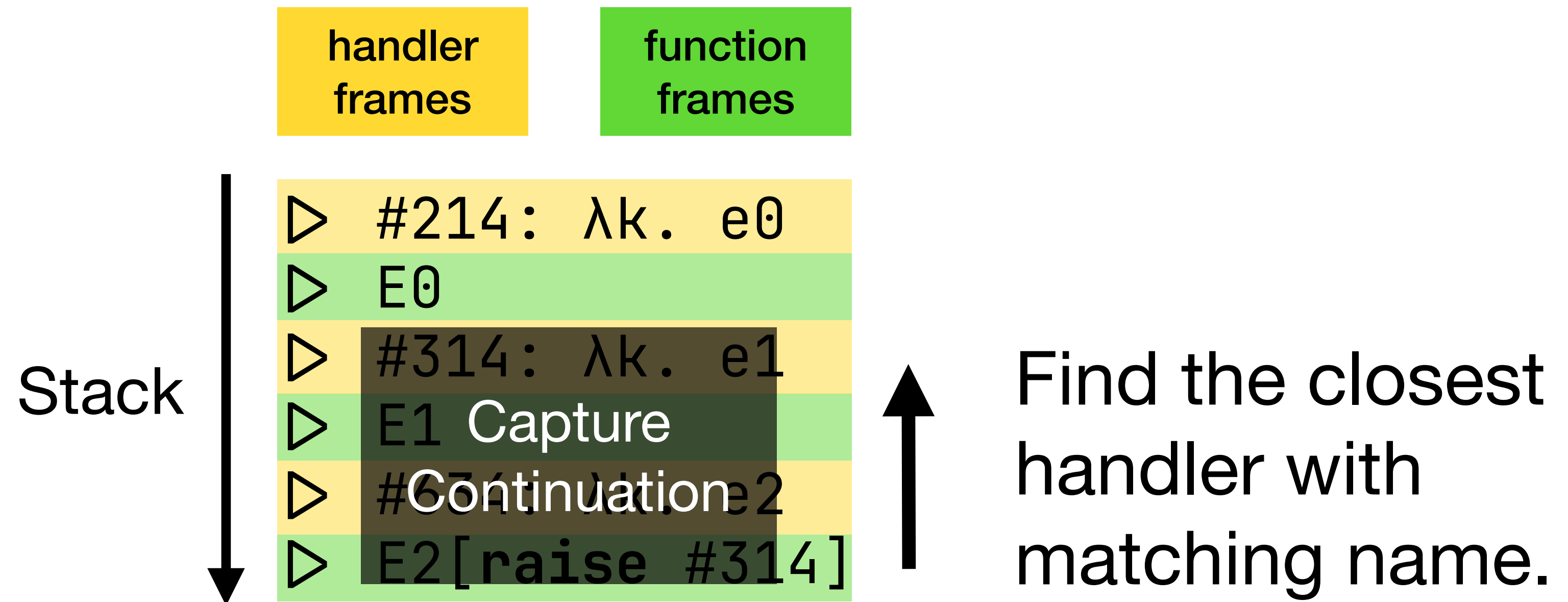
Lexical effect handler

Last bit of semantics: install a handler



Lexical effect handler

Last bit of semantics: install a handler



Lexical effect handler

A powerful language feature

Syntax

handle E with H

raise ...

resume ...

+

Semantics

Install a handler

Find handler; Capture continuation

Resume continuation

Lexical effect handler

A powerful language feature

Syntax

handle E with H

raise ...

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+

Semantics

Install a handler

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Algebraic Effect
Plotkin and Power

Effect Handler
Plotkin and Pretnar

Lexical Handler
Zhang and Myers

 OCaml

Lexa

 WA

仓颉

2003

2013

2019

2022

2024

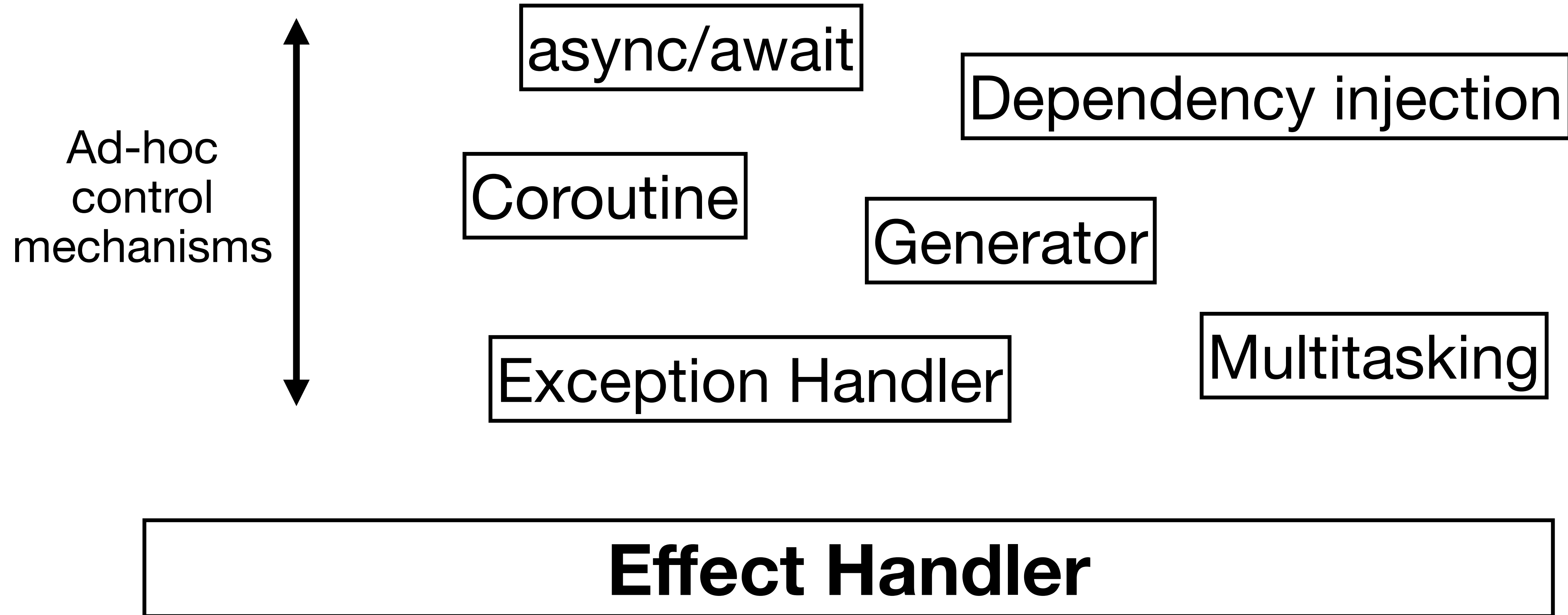
WIP

WIP

Background

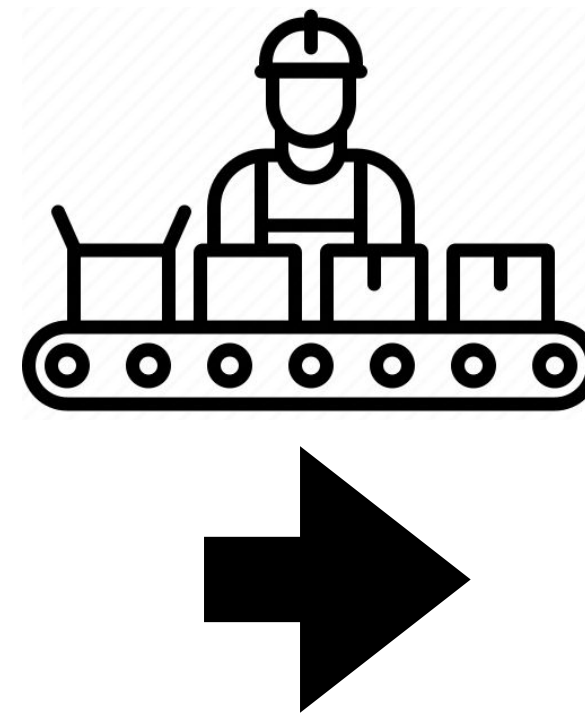
Effect handler

A powerful language feature



high-level effect handlers in Lexa

```
handle E with H  
  raise ...  
  resume ...
```



low-level stack switching in assembly

```
ENTER  
RAISE  
RESUME
```

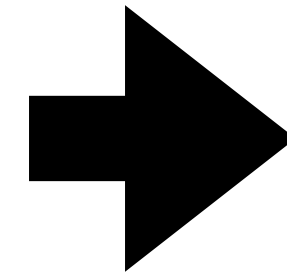
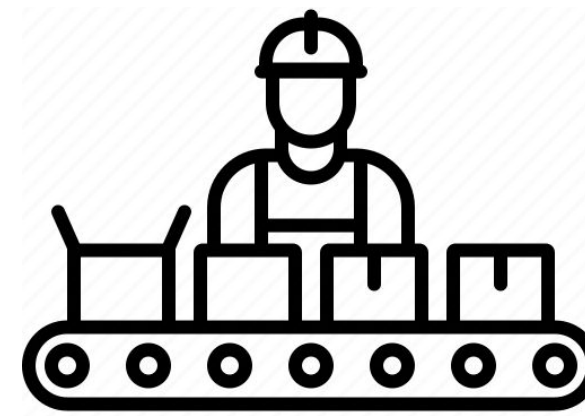
Direct Lexa
OOPSLA24

Zero Lexa
OOPSLA25

Multi Lexa
work-in-progress

high-level effect handlers in Lexa

```
handle E with H  
  raise ...  
  resume ...
```



low-level stack switching in assembly

```
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RESUME
```

Direct Lexa

A large, hollow black triangle pointing upwards. Inside the triangle, the text "tradeoffs between performance and expressivity" is centered and written in a sans-serif font.

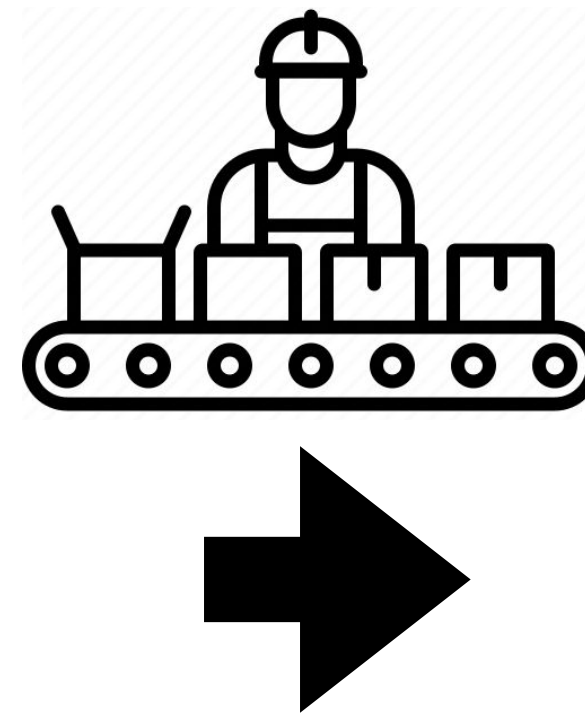
tradeoffs
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Multi Lexa

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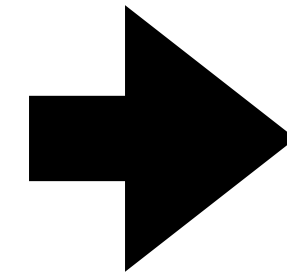
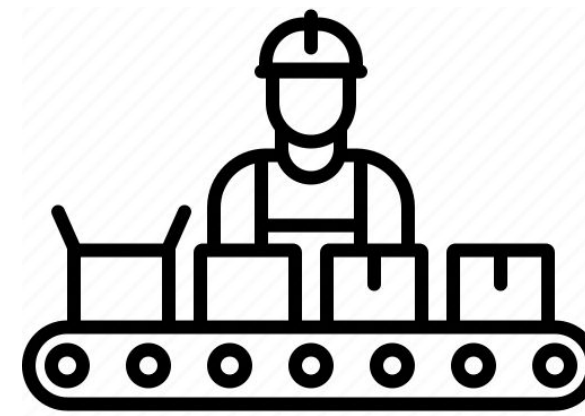
formally
proved
correct
[[*E*]] = [[*E*]]

Multi Lexa

Zero Lexa

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Multi Lexa

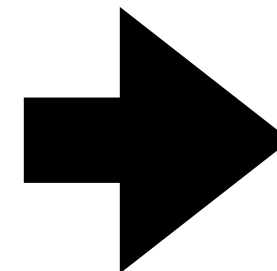
Zero Lexa

Direct Lexa

Idea 1: use stack addresses as handler names

```
handle
  while true
    DO_SOMETHING
    raise yield()
with yield(k) ⇒
  SCHEDULER
```

reduce to



```
▷ #314: SCHEDULER
▷ E[raise #314()]
```

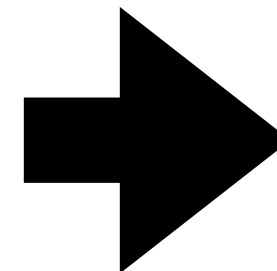
globally
unique names

Direct Lexa

Idea 1: use stack addresses as handler names

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handle
  while true
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    raise yield()
with yield(k) ⇒
  SCHEDULER
```

reduce to



stack addresses

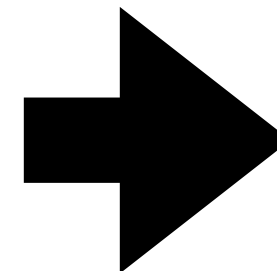
```
0x123 ▷ #314: SCHEDULER
0x124 ▷ E[raise #314()]
```

Direct Lexa

Idea 1: use stack addresses as handler names

```
handle
  while true
    DO_SOMETHING
    raise yield()
with yield(k) ⇒
  SCHEDULER
```

reduce to



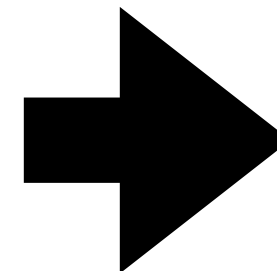
```
0x123 ▷ SCHEDULER
0x124 ▷ E[raise ?()]
```

Direct Lexa

Idea 1: use stack addresses as handler names

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reduce to



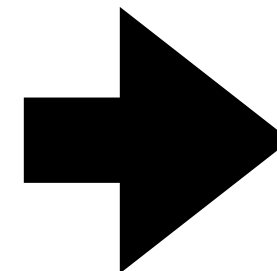
```
0x123 ▷ SCHEDULER
0x124 ▷ E[raise 0x123()]
```

Direct Lexa

Idea 1: use stack addresses as handler names

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    raise yield()
with yield(k) ⇒
  SCHEDULER
```

reduce to



```
0x123 ▷ SCHEDULER
0x124 ▷ E[raise 0x123()]
```

locate handler
in $O(1)$

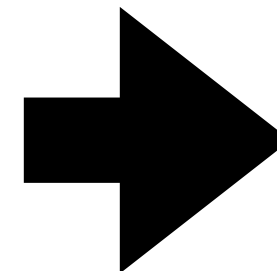
Direct Lexa

Idea 1: use stack addresses as handler names

This design will come back and bite us—more on that later.

```
handle
  while true
    DO_SOMETHING
    raise yield()
with yield(k) ⇒
  SCHEDULER
```

reduce to



```
0x123 ▷ SCHEDULER
0x124 ▷ E[raise 0x123()]
```

Direct Lexa

Idea 2: Stack switching

Direct Lexa

Idea 2: Stack switching

INSTALL

RAISE

RESUME

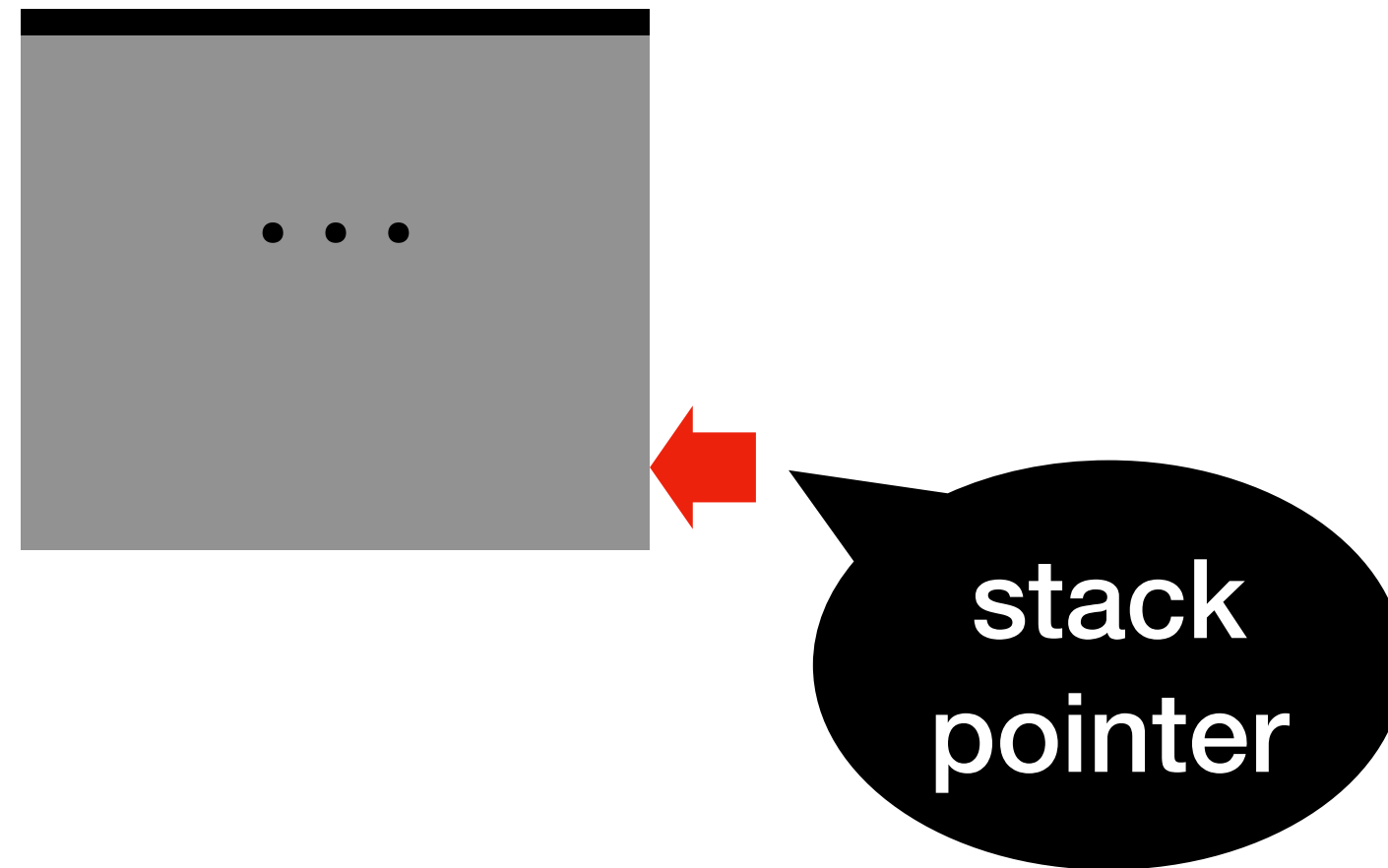
Direct Lexa Stack switching

```
handle
  while true
    DO_SOMETHING
  with yield(k) =>
    SCHEDULER
```

INSTALL

RAISE

RESUME



Direct Lexa Stack switching

```
handle  
  while true  
    DO_SOMETHING  
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    SCHEDULER
```

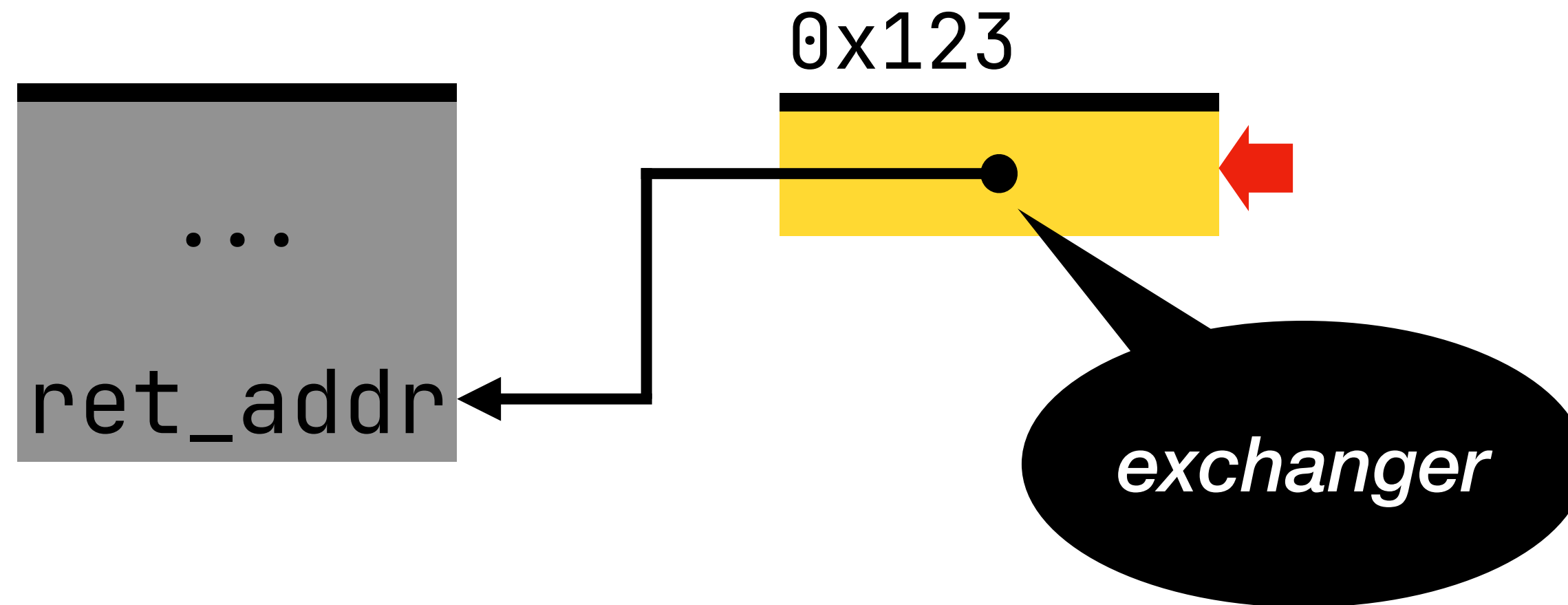
INSTALL
RAISE
RESUME



Direct Lexa Stack switching

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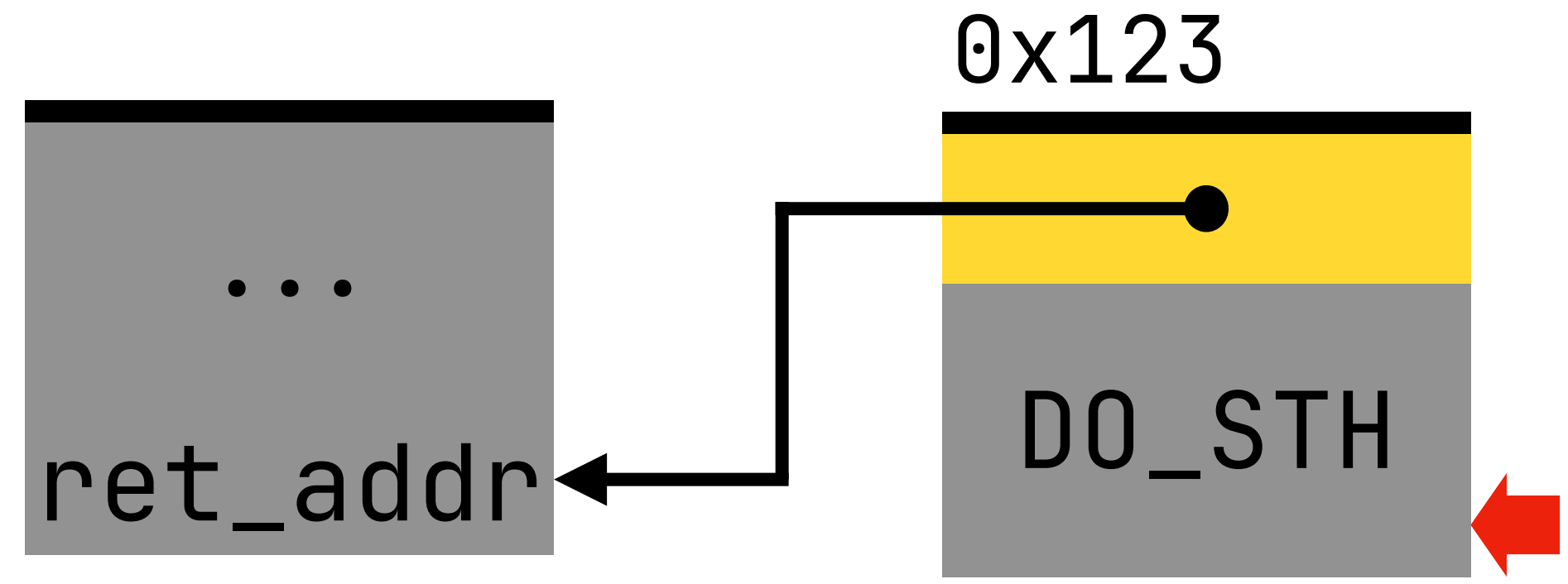
INSTALL
RAISE
RESUME



Direct Lexa Stack switching

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```

INSTALL
RAISE
RESUME



Direct Lexa Stack switching

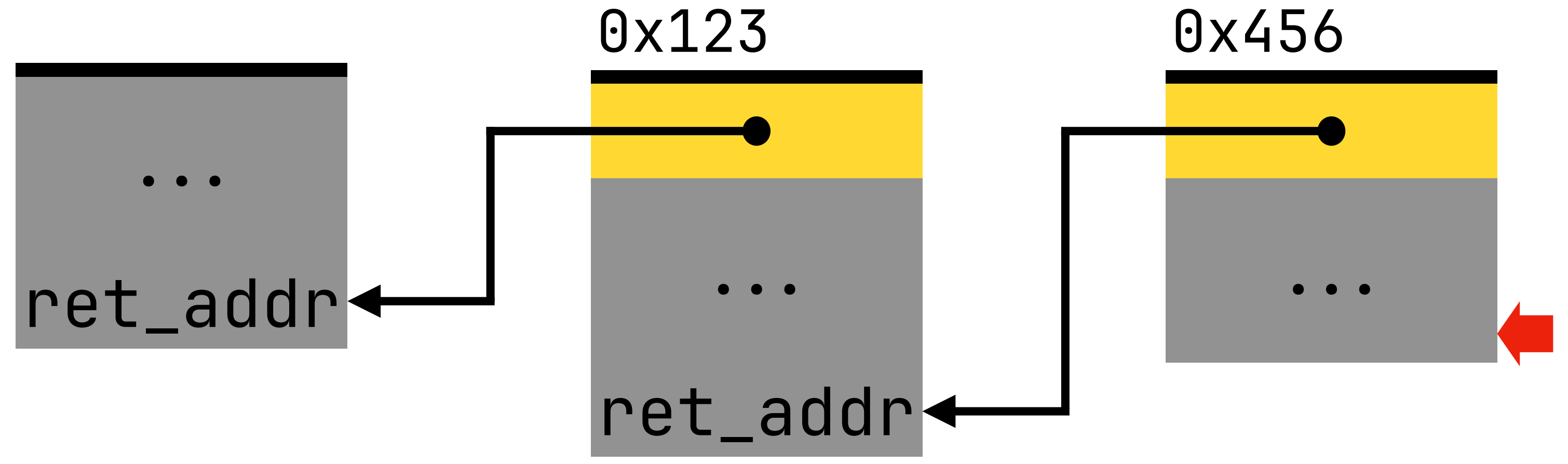
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```

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  ...  
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  with ...
```

INSTALL

RAISE

RESUME

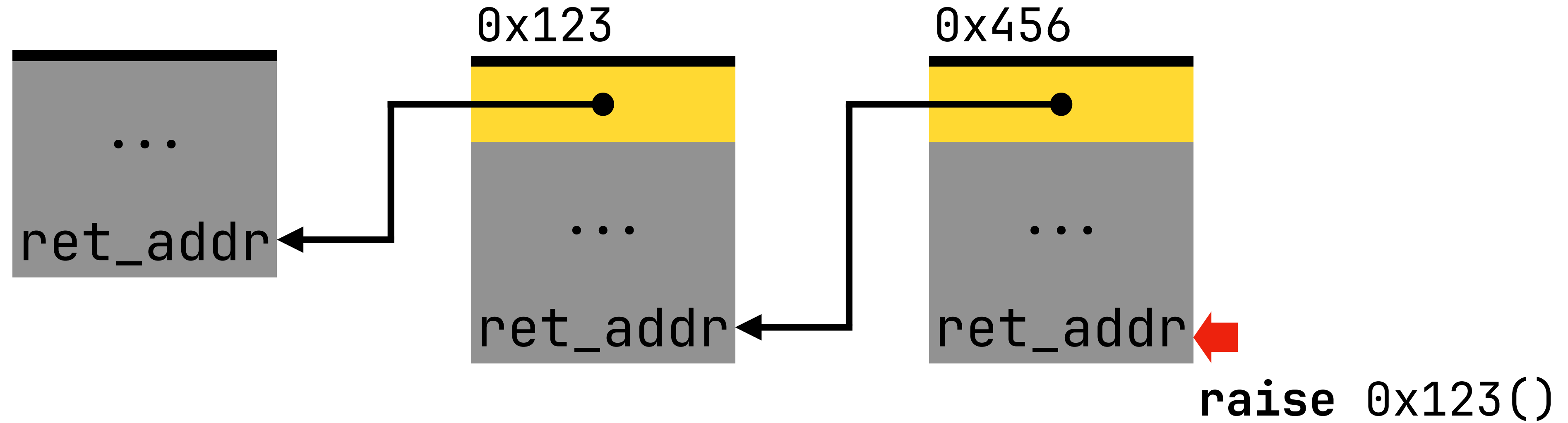


Direct Lexa Stack switching

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  with yield(k) ⇒  
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```

```
handle  
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INSTALL
RAISE
RESUME

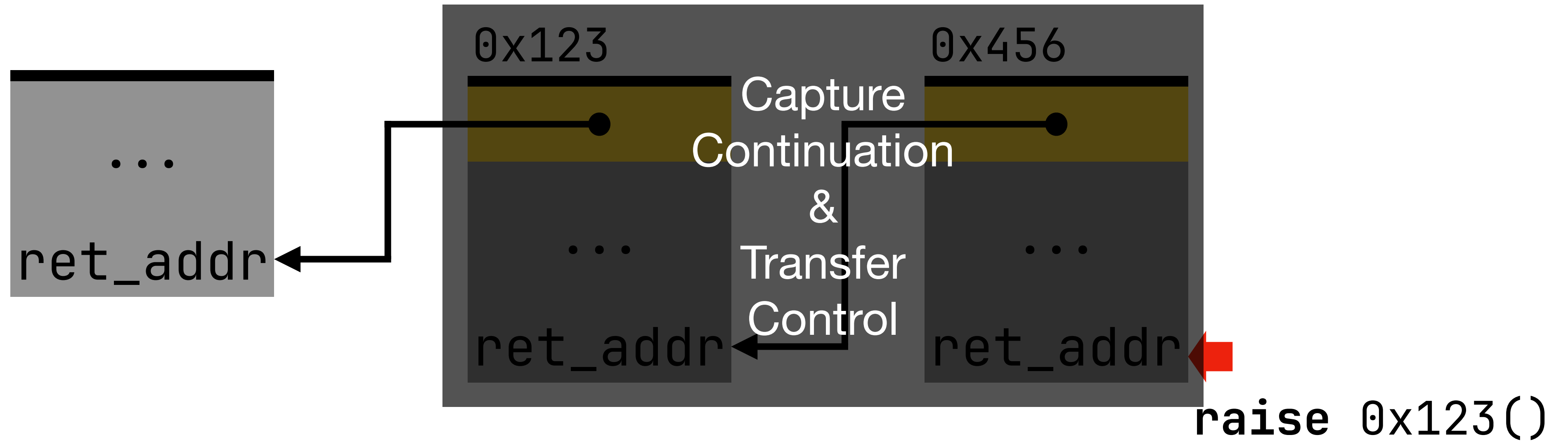


Direct Lexa Stack switching

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INSTALL
RAISE
RESUME

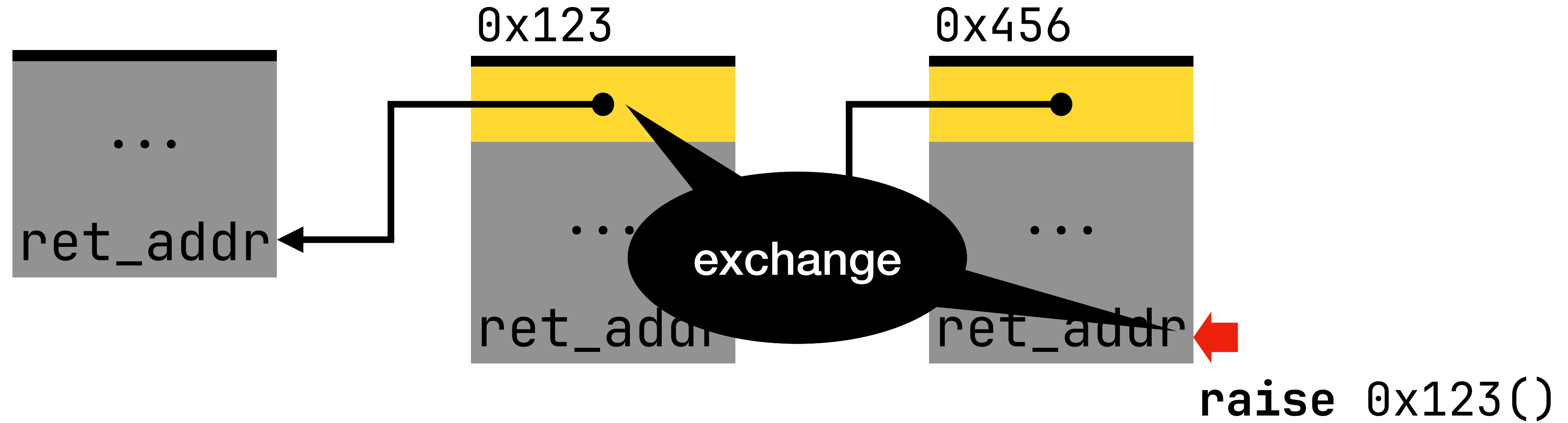


Direct Lexa Stack switching

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INSTALL
RAISE
RESUME

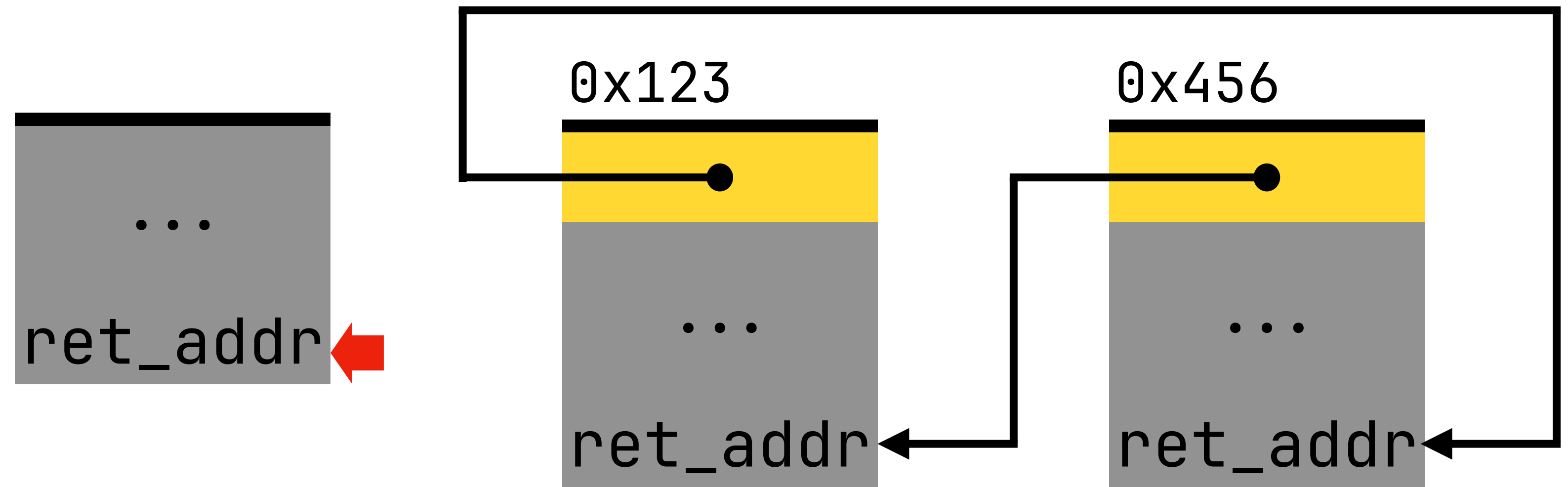


Direct Lexa Stack switching

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```

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  ...  
  raise yield()  
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INSTALL
RAISE
RESUME

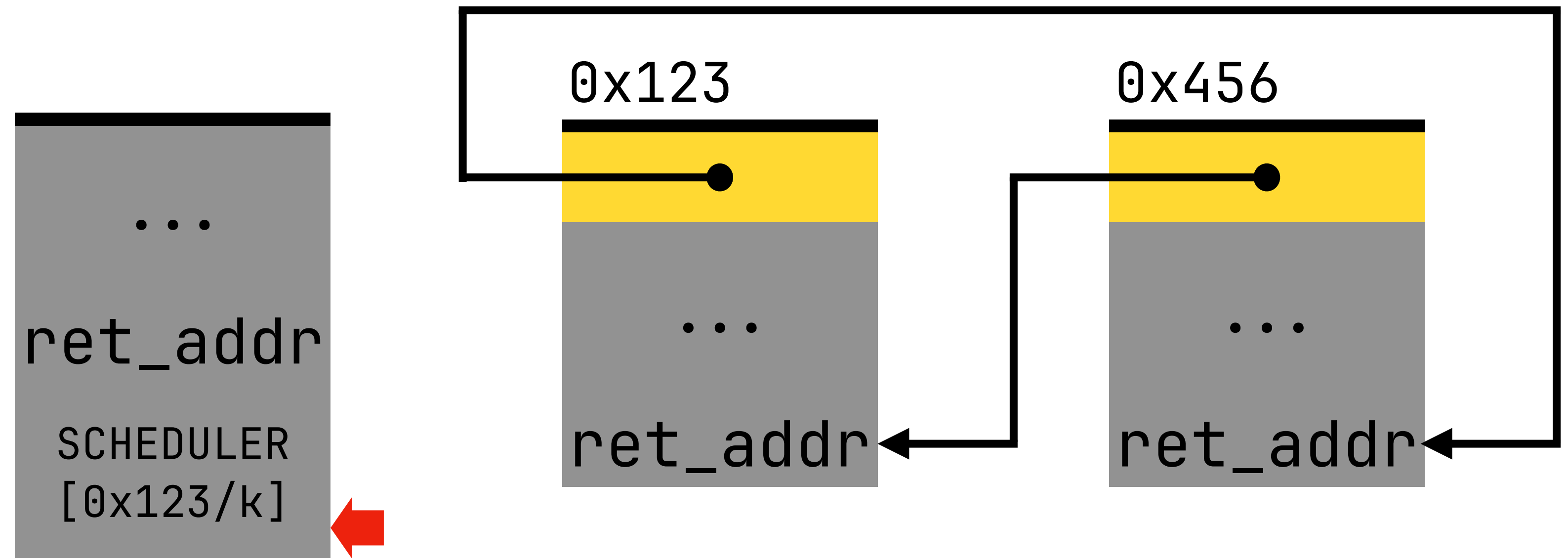


Direct Lexa Stack switching

```
handle  
  while true  
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  with yield(k) ⇒  
  SCHEDULER
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handle  
  ...  
  raise yield()  
  with ...
```

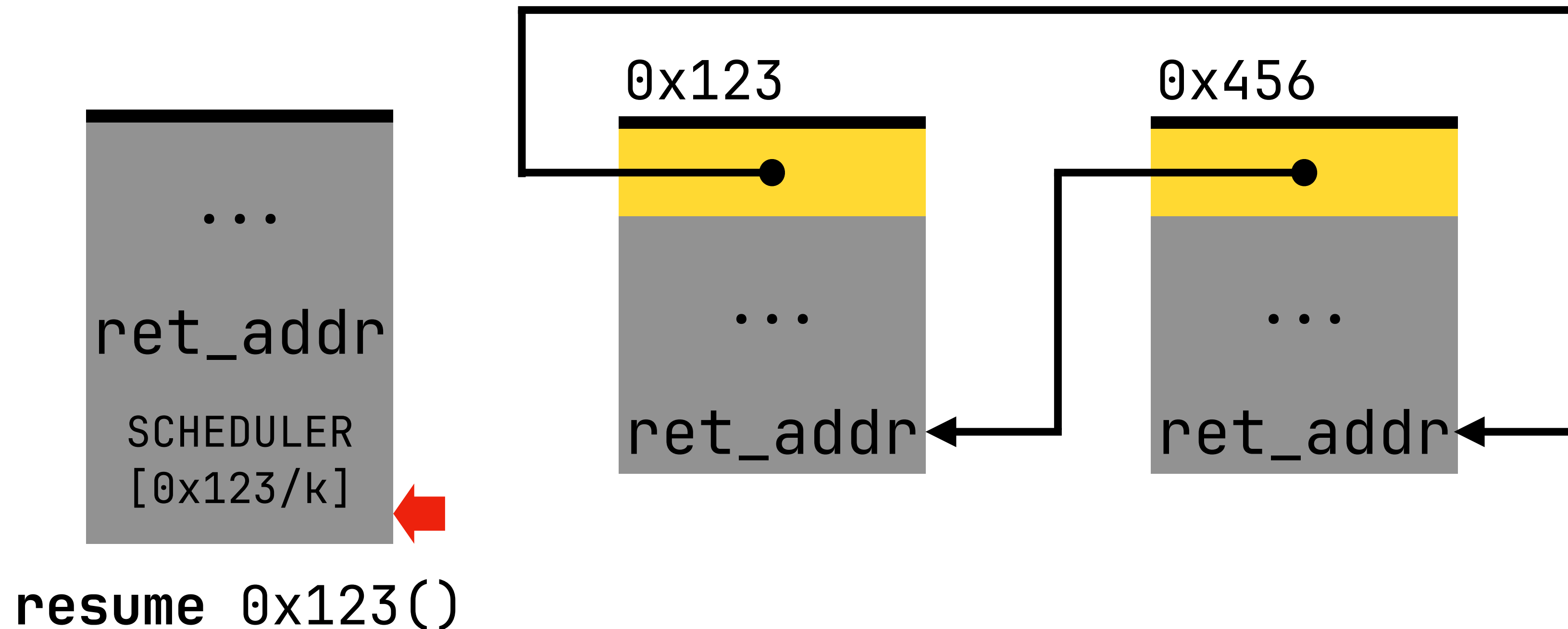
INSTALL
RAISE
RESUME



Direct Lexa Stack switching

```
handle
  while true
    DO_SOMETHING
  with yield(k) =>
    ...
  resume peer ()
```

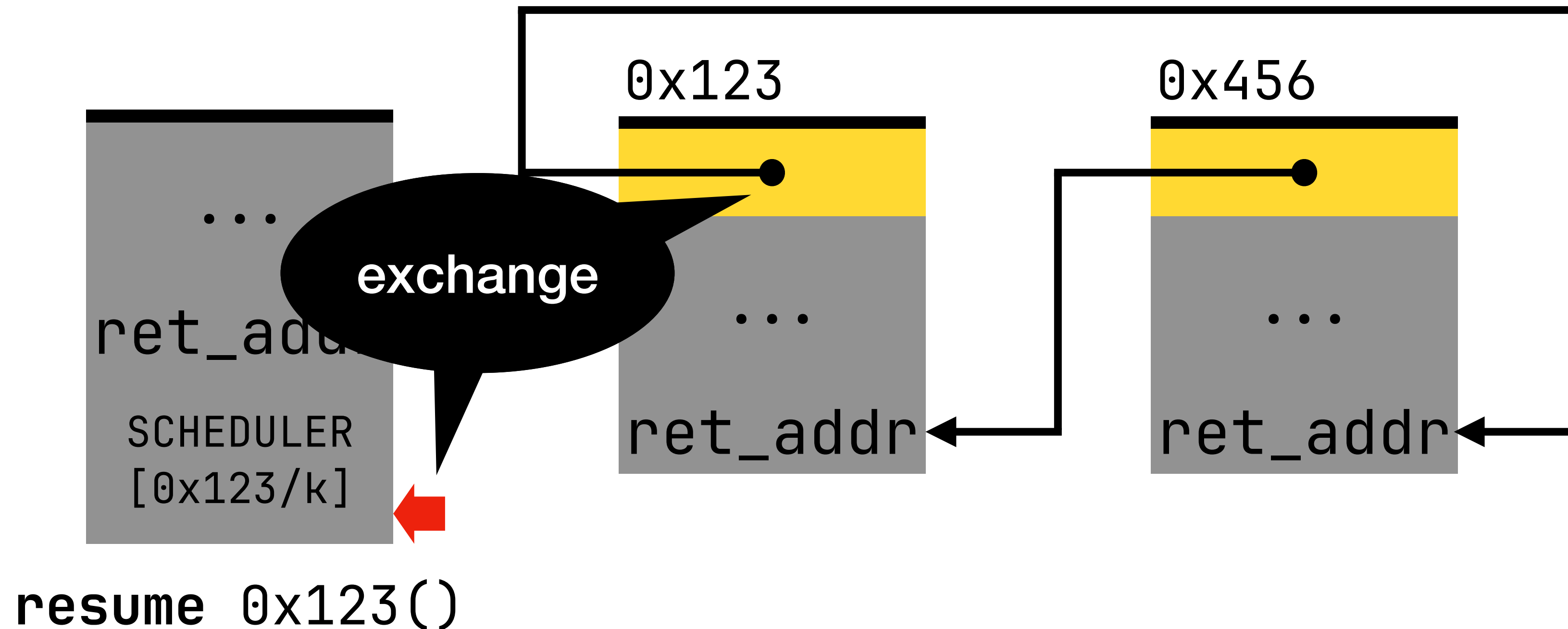
INSTALL
RAISE
RESUME



Direct Lexa Stack switching

```
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  while true  
    DO_SOMETHING  
  with yield(k) =>  
    ...  
  resume peer ()
```

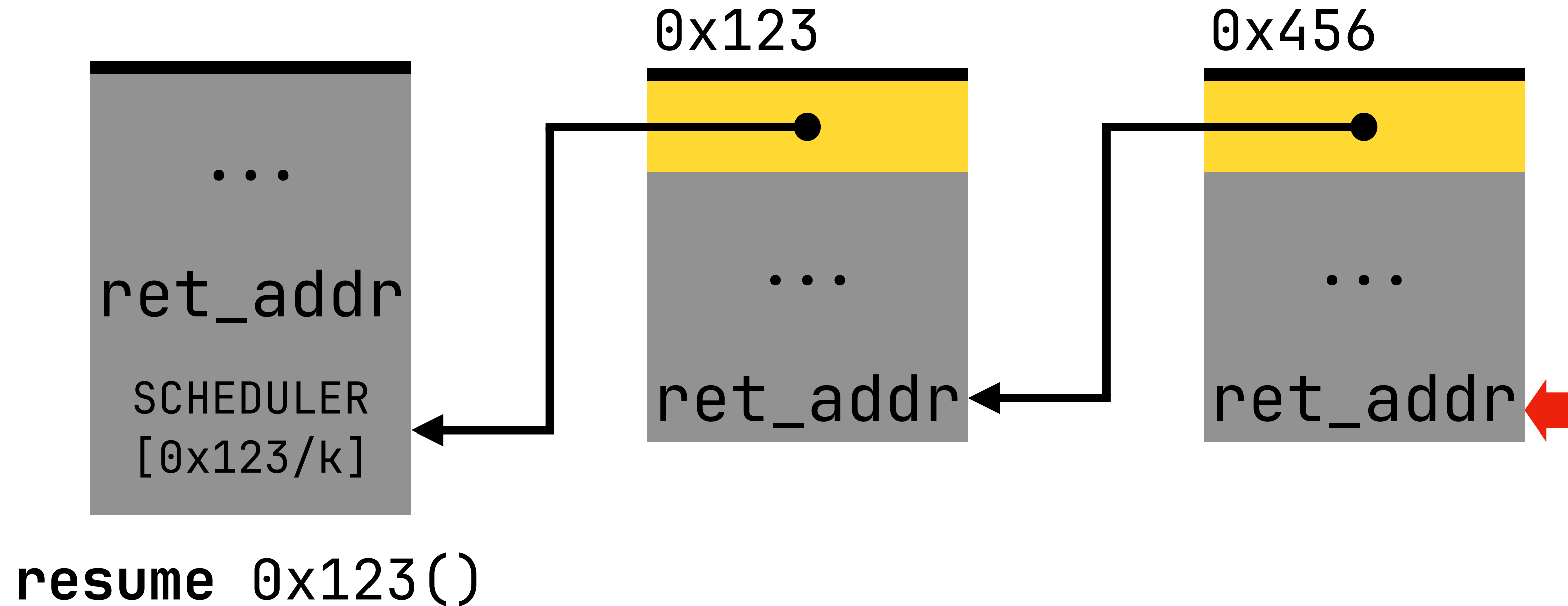
INSTALL
RAISE
RESUME



Direct Lexa Stack switching

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  with yield(k) =>
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  resume peer ()
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INSTALL
RAISE
RESUME

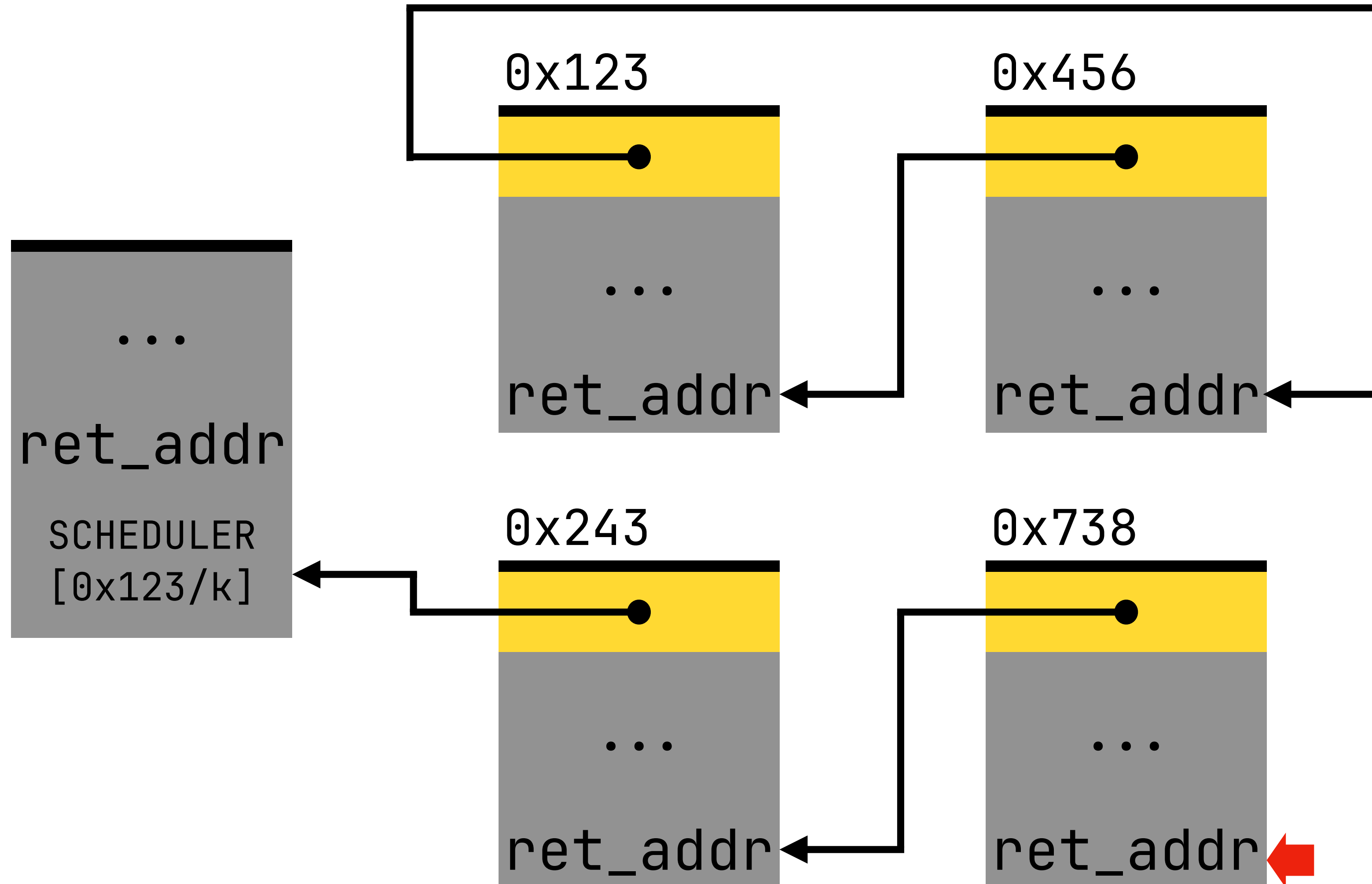


Direct Lexa

Stack switching

Running Example

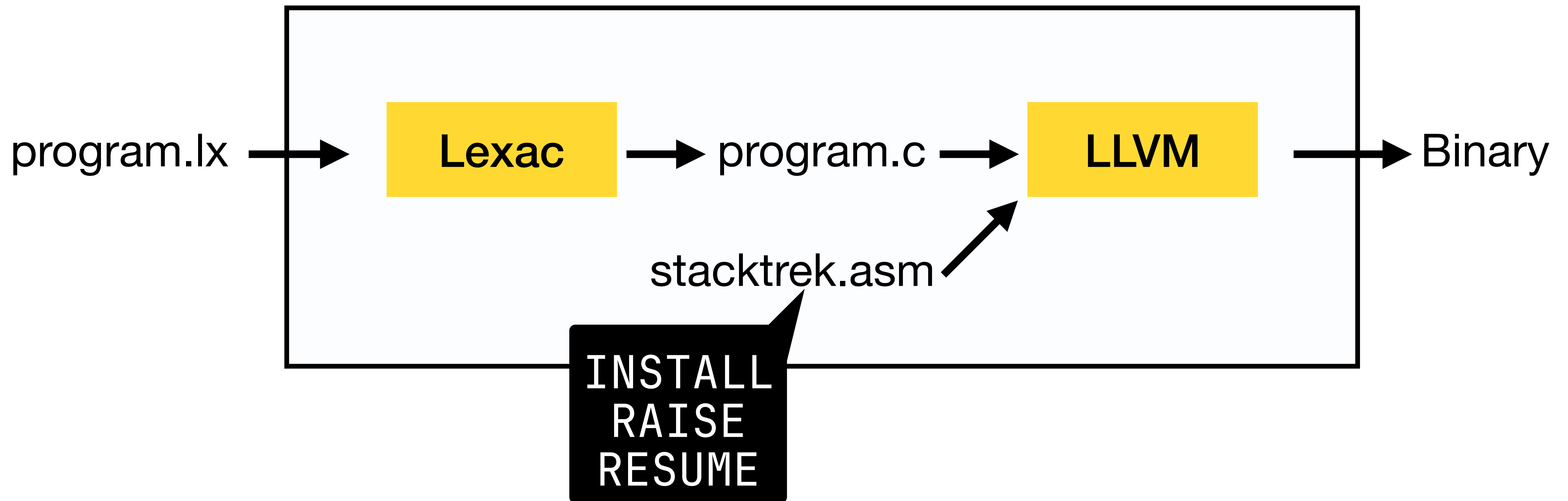
The Full Picture



Lexa compiler implementation

stacktrek Library

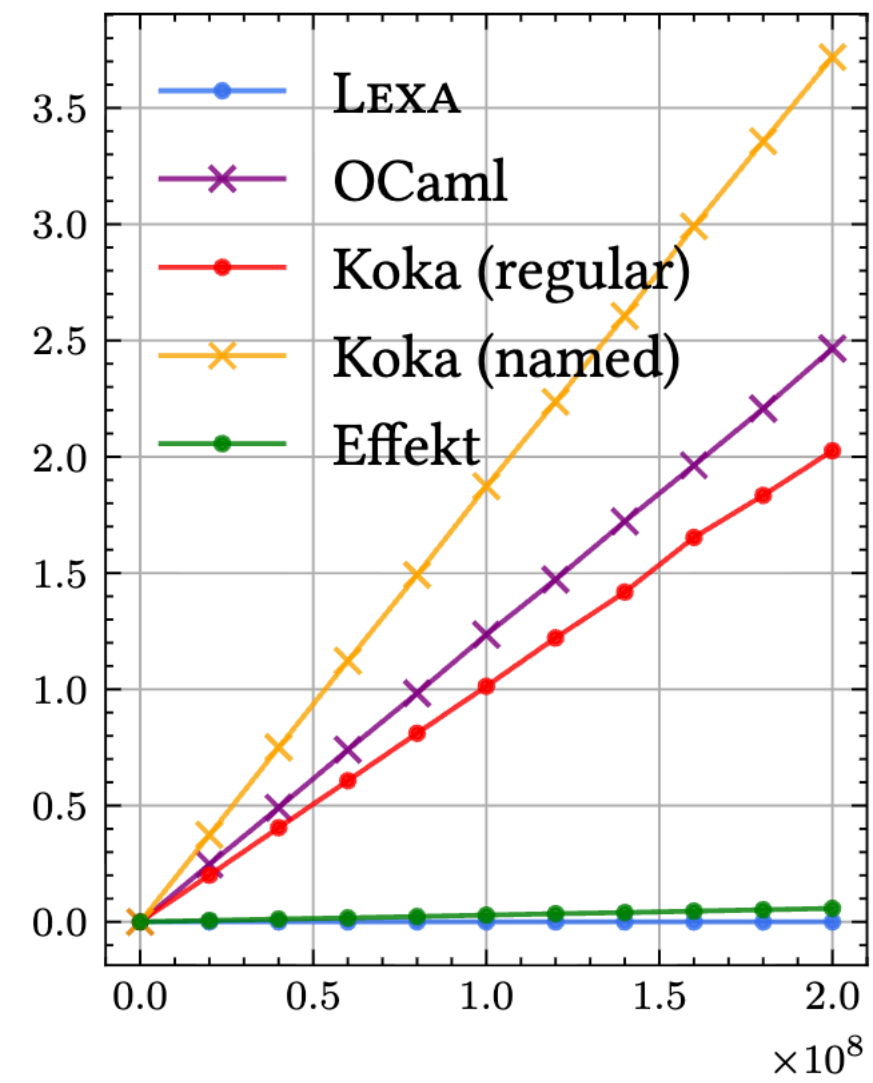
Lexa Compiler



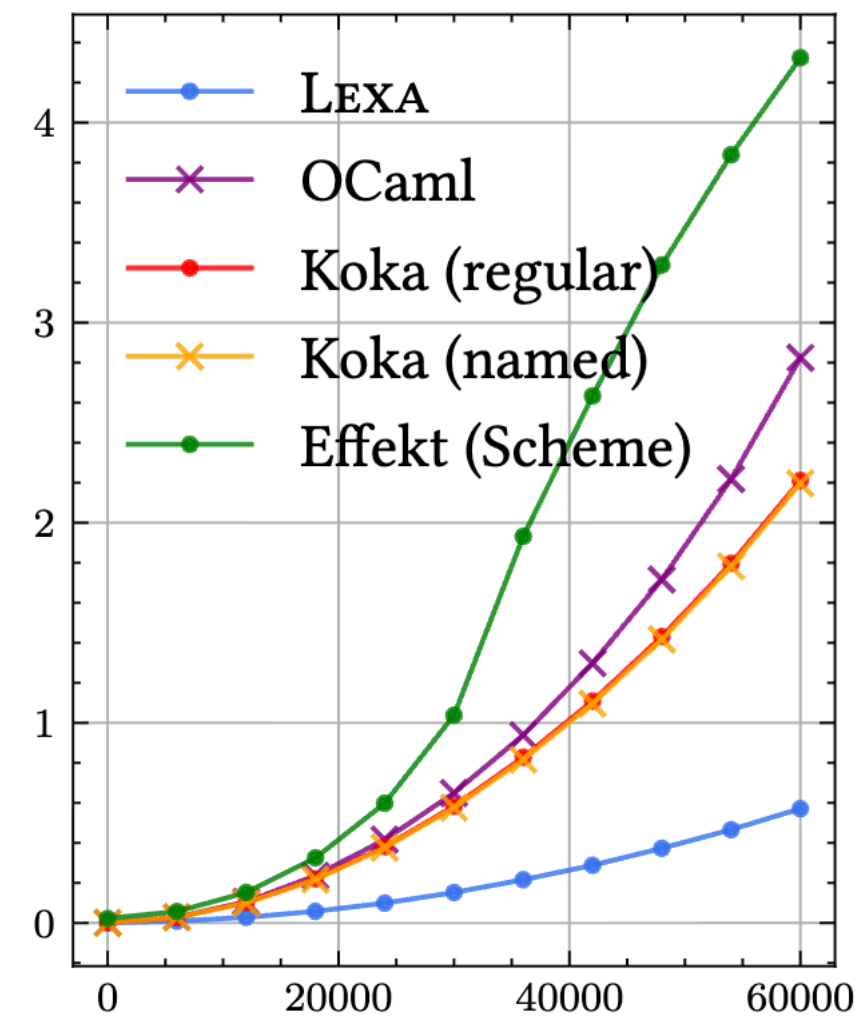
Lexa is fast

Selected benchmark results

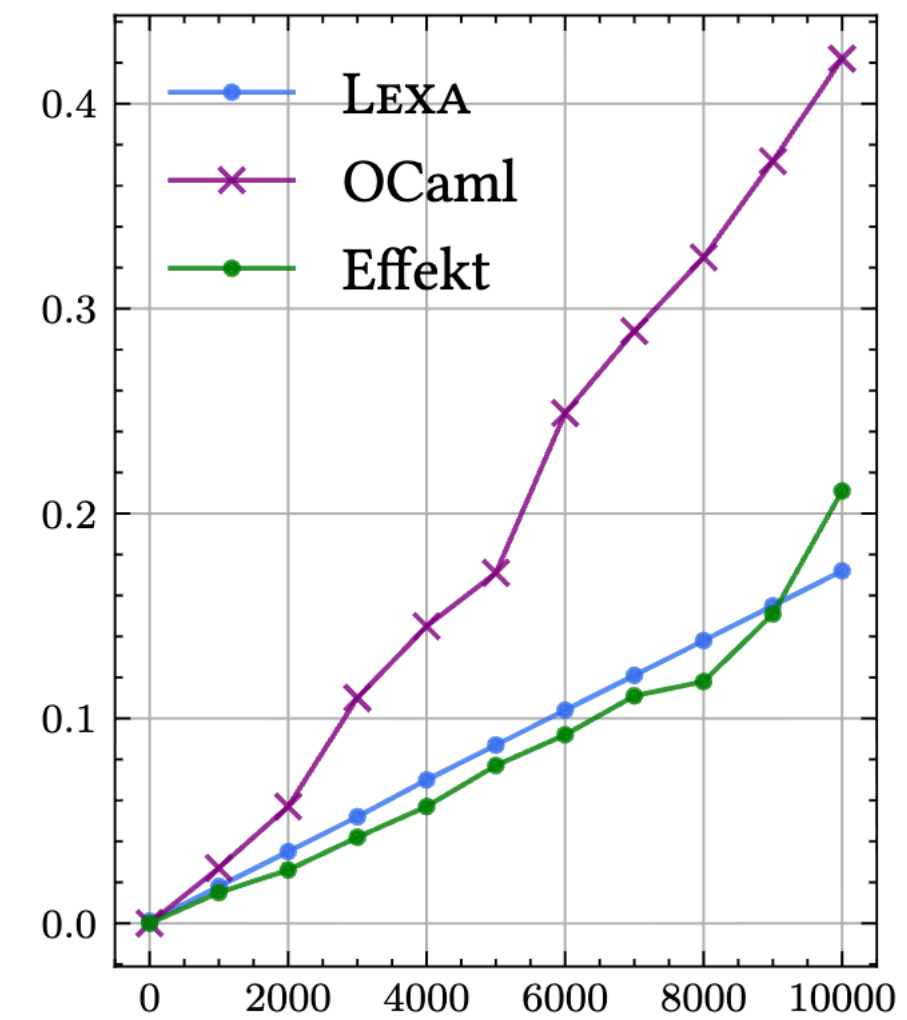
Countdown



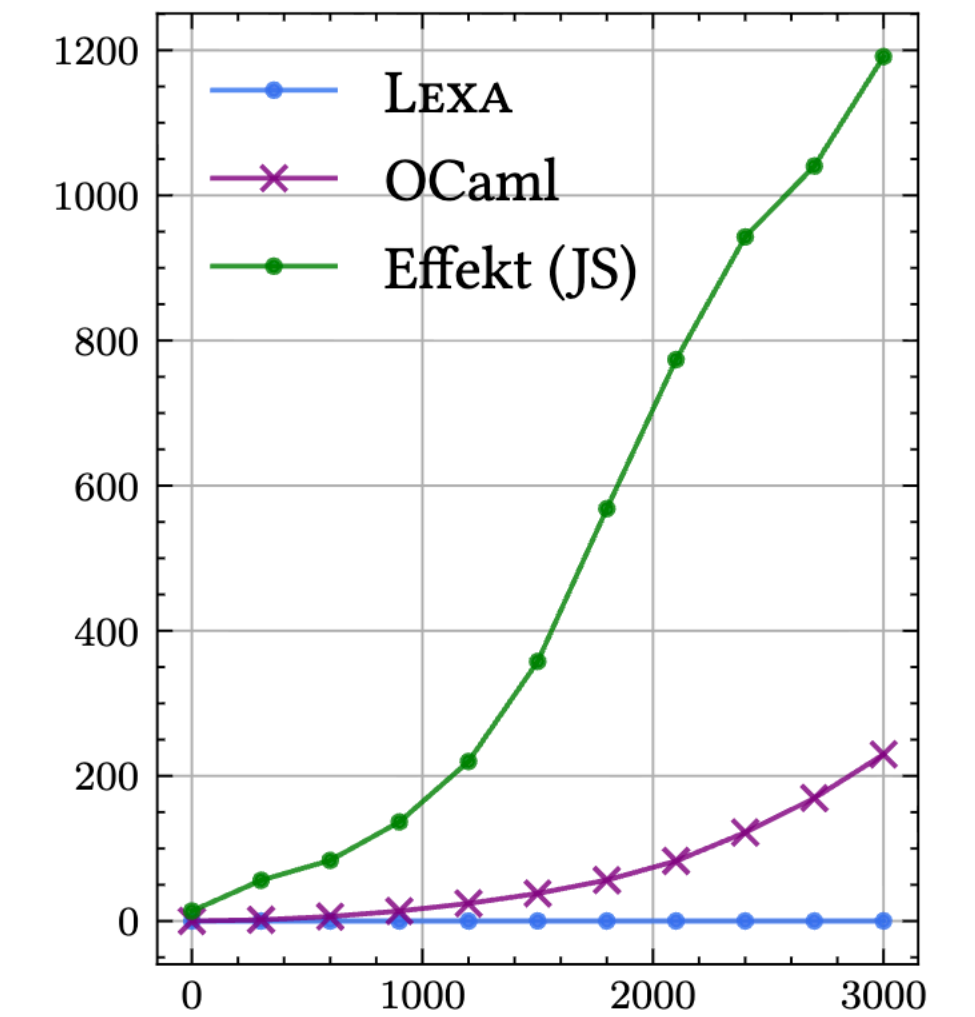
Handler Sieve



Resume Nontail 2

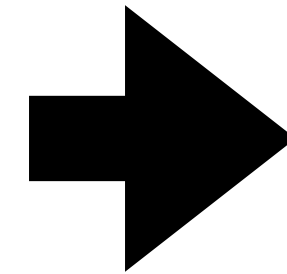
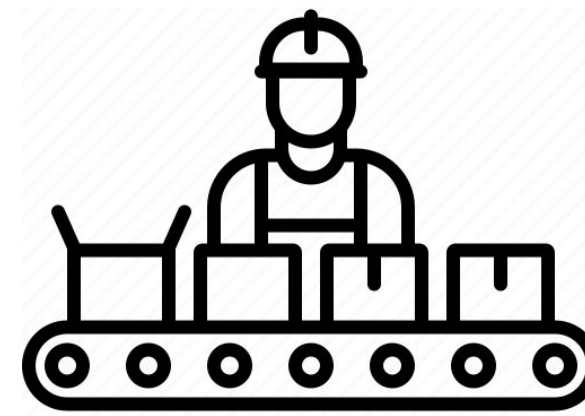


Interruptible Iterator



high-level effect handlers in Lexa

```
handle E with H  
  
raise ...  
  
resume ...
```



low-level stack switching in assembly

```
ENTER  
  
RAISE  
  
RESUME
```

Direct Lexa

A large, hollow black triangle pointing upwards. Inside the triangle, the text "tradeoffs between performance and expressivity" is centered.

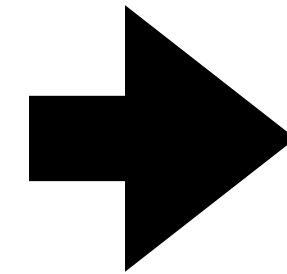
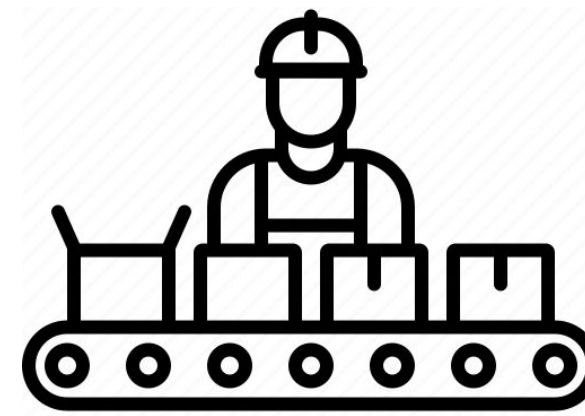
tradeoffs
between
performance
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Multi Lexa

Zero Lexa

high-level effect handlers in Lexa

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tradeoffs
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Multi Lexa

Zero Lexa

Multi Lexa

Multishot resumptions are expressive.

Multi Lexa

Multishot resumptions are expressive.

```
def checkpointing(program):  
  var cp = None;  
  handle  
    program(save, retry)  
  with save =  
    λx,k. cp = Some(k); resume k ()  
  with retry =  
    match cp  
    | Some k ⇒ resume k ()  
    | None ⇒ checkpointing(program)
```



program

Multi Lexa

Multishot resumptions are expressive.

```
def checkpointing(program):  
  var cp = None;  
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  with save =  
    λx,k. cp = Some(k); resume k ()  
  with retry =  
    match cp  
    | Some k ⇒ resume k ()  
    | None ⇒ checkpointing(program)
```



save
checkpoint



retry
checkpoint

Multi Lexa

Direct Lexa's use of stack address comes back and bites us.

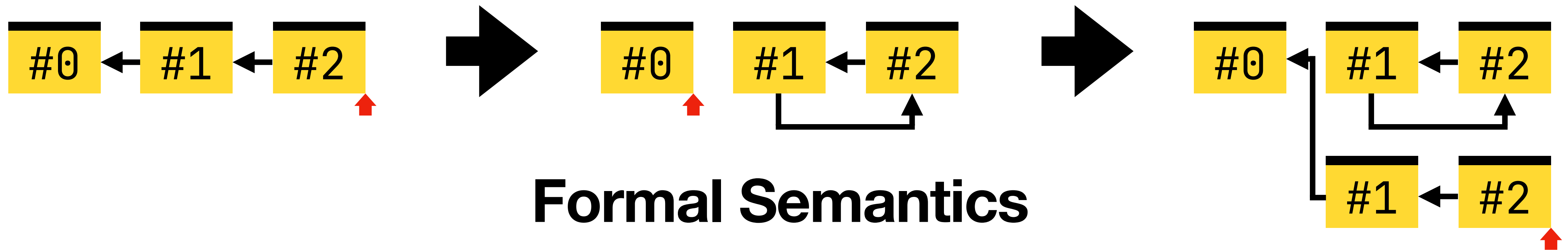
Direct Lexa does not support multishot resumptions.

Formal Semantics

Direct Lexa Implementation

Multi Lexa

Direct Lexa does not support multishot resumptions.

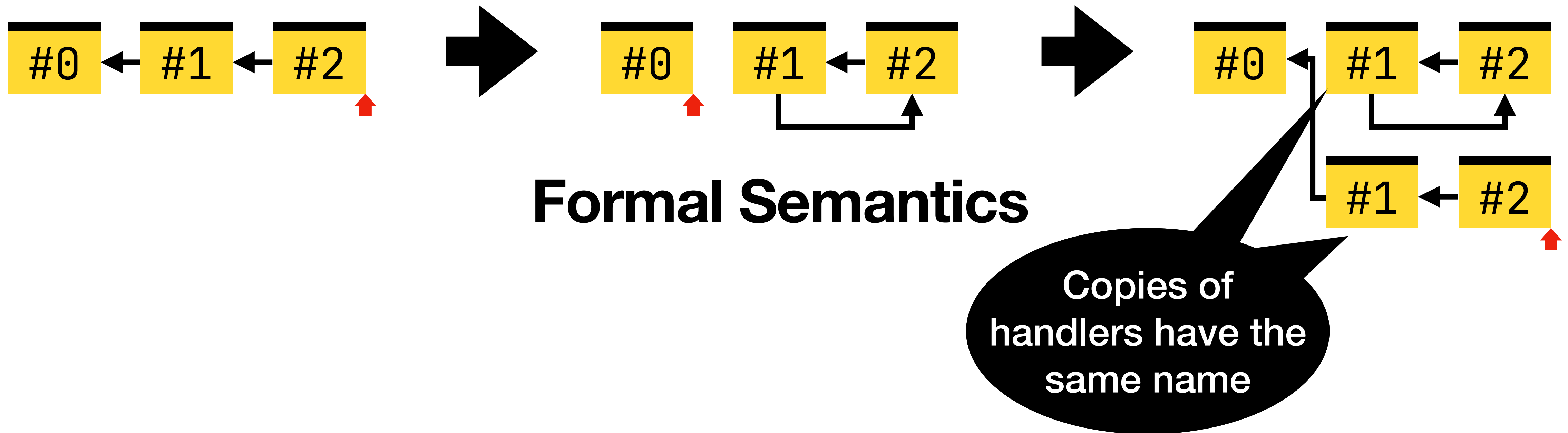


Formal Semantics

**Direct Lexa
Implementation**

Multi Lexa

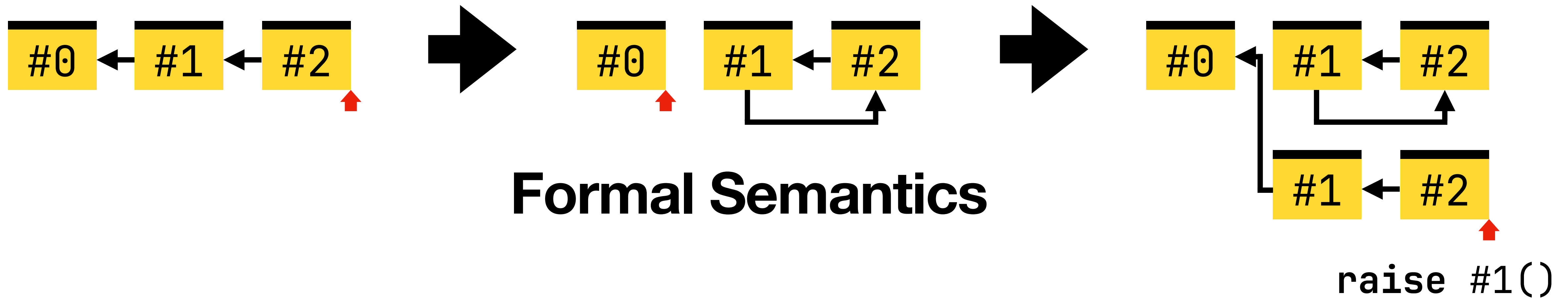
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**Direct Lexa
Implementation**

Multi Lexa

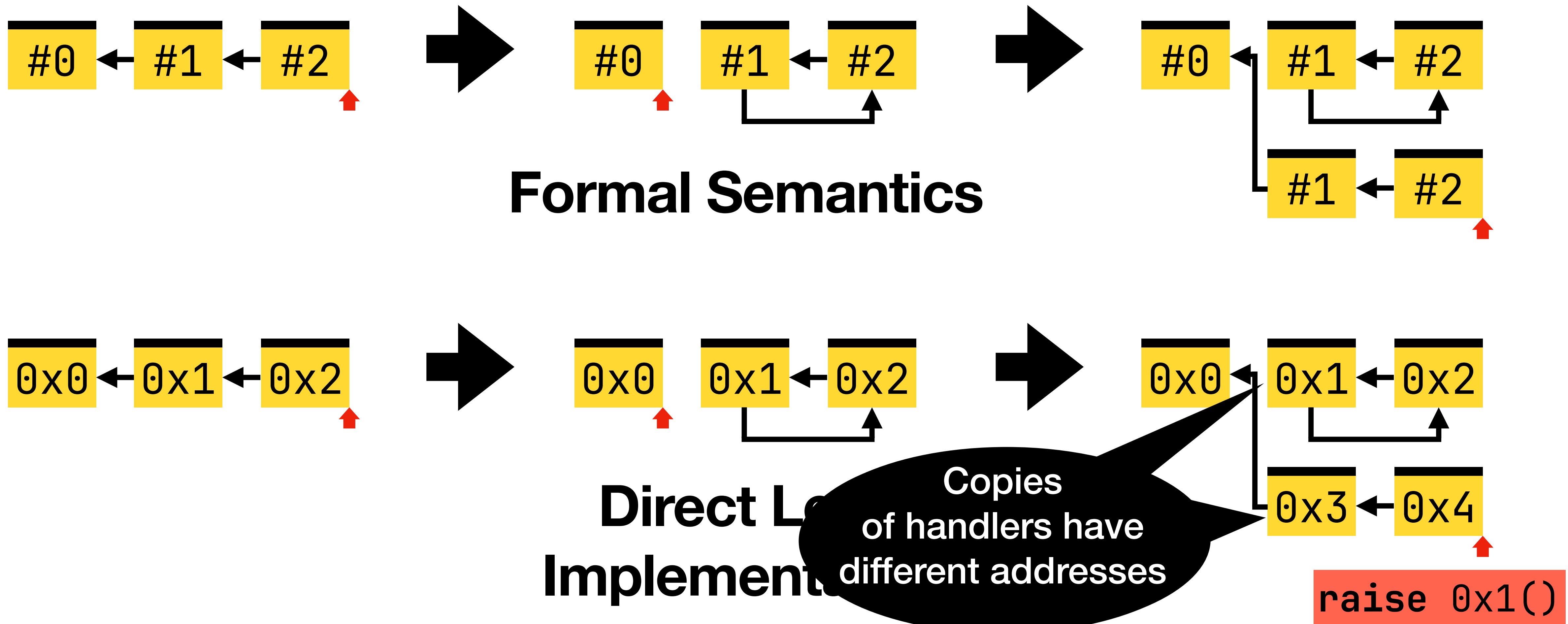
Direct Lexa does not support multishot resumptions.



**Direct Lexa
Implementation**

Multi Lexa

Direct Lexa does not support multishot resumptions.



Multi Lexa

Main idea: virtual resumption

Multi Lexa

Main idea: virtual resumption

- Stacks and handlers live in virtual memory space, and all addresses are virtual addresses
- Each continuation lives in a different zone; copies of the same handlers in different continuation live at the same relative address within zones

0x73A

Zone number

Handler ID

Stack Offset

Zone 0



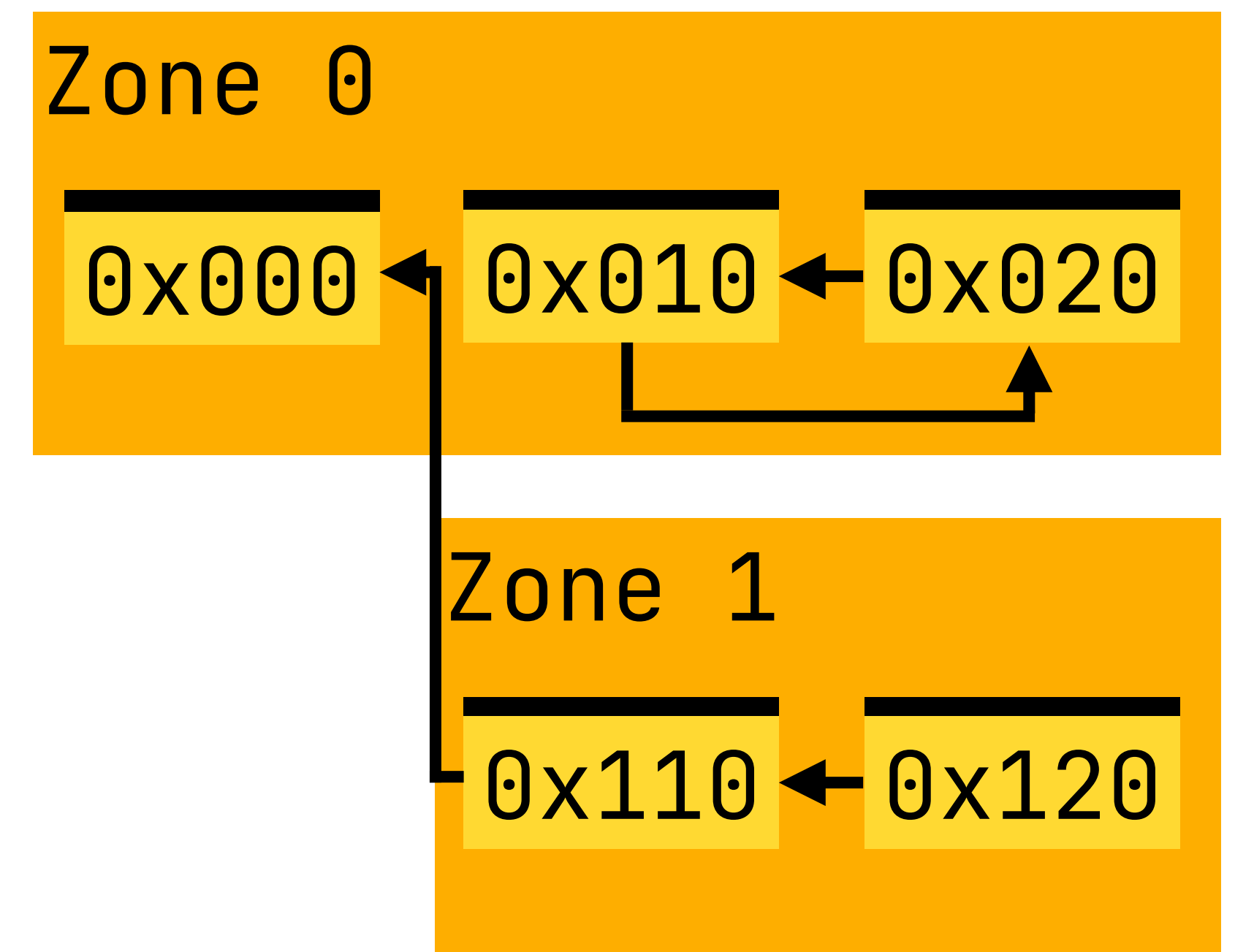
Multi Lexa

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0x73A

Zone number
Handler ID
Stack Offset



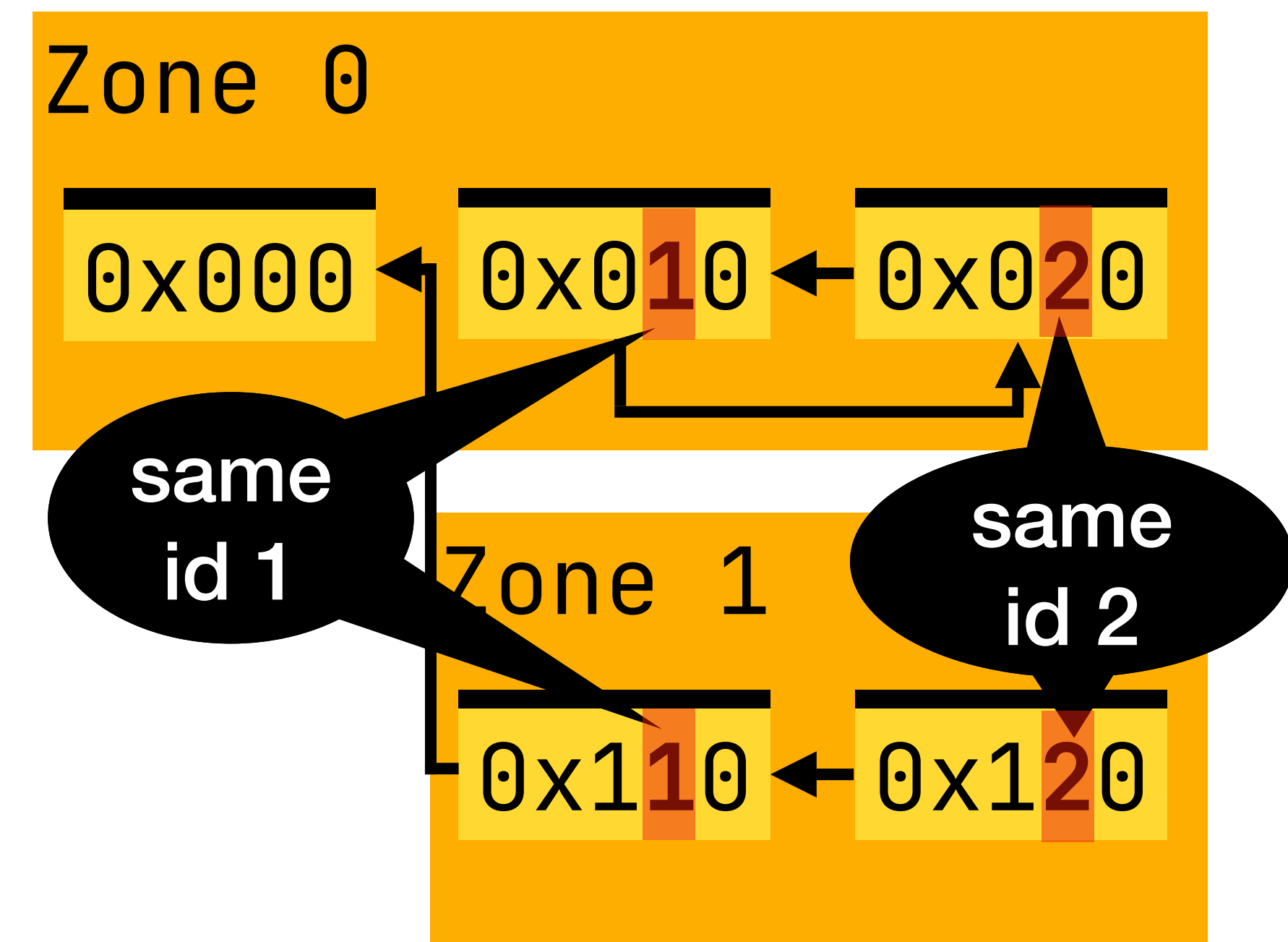
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0x73A

Zone number
Handler ID
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Multi Lexa

Main idea: virtual resumption

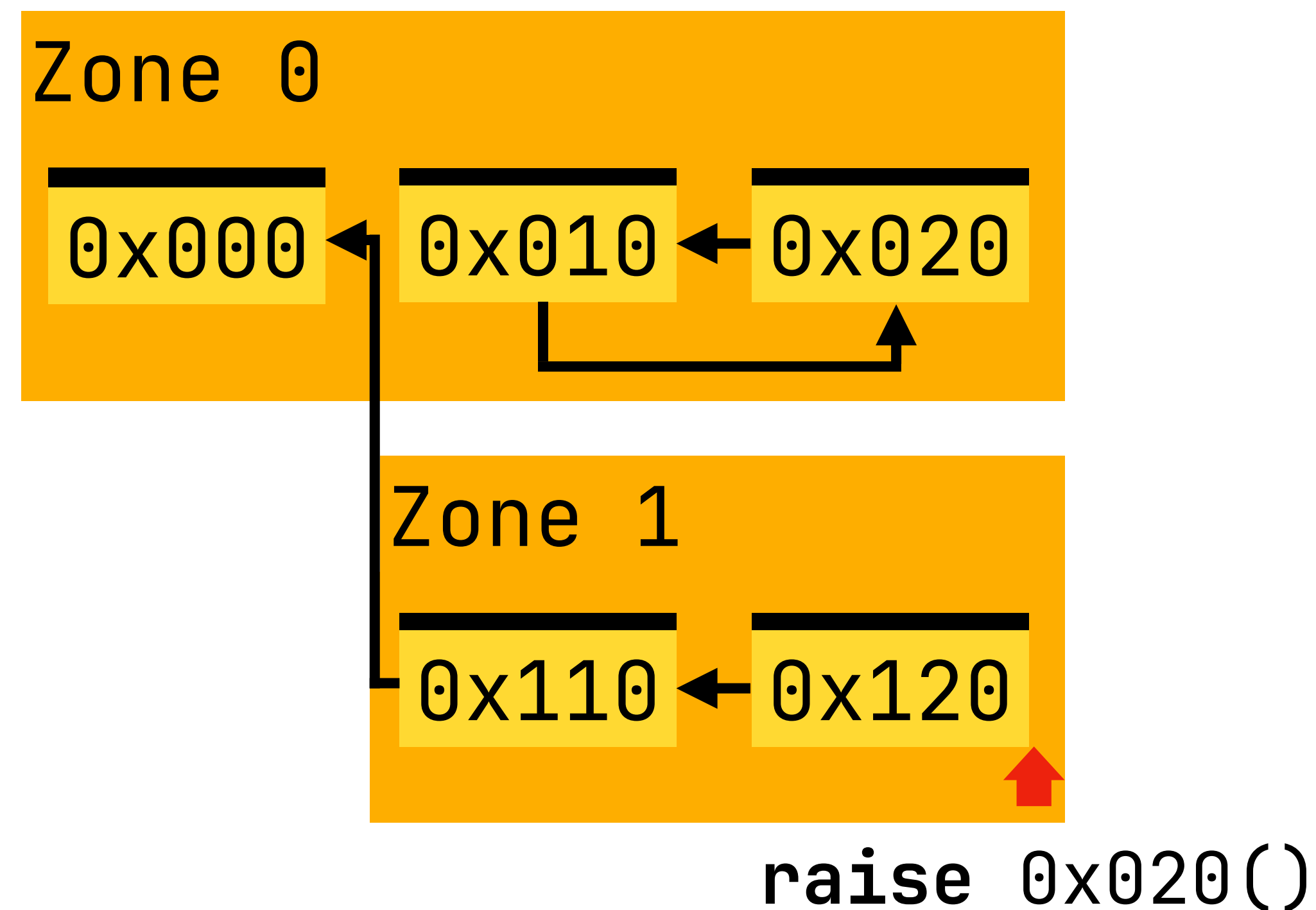
0x73A

Zone number

Handler ID

Stack Offset

- A software-level memory management unit(**sMMU**) is in charge of memory allocation and address translation



Multi Lexa

Main idea: virtual resumption

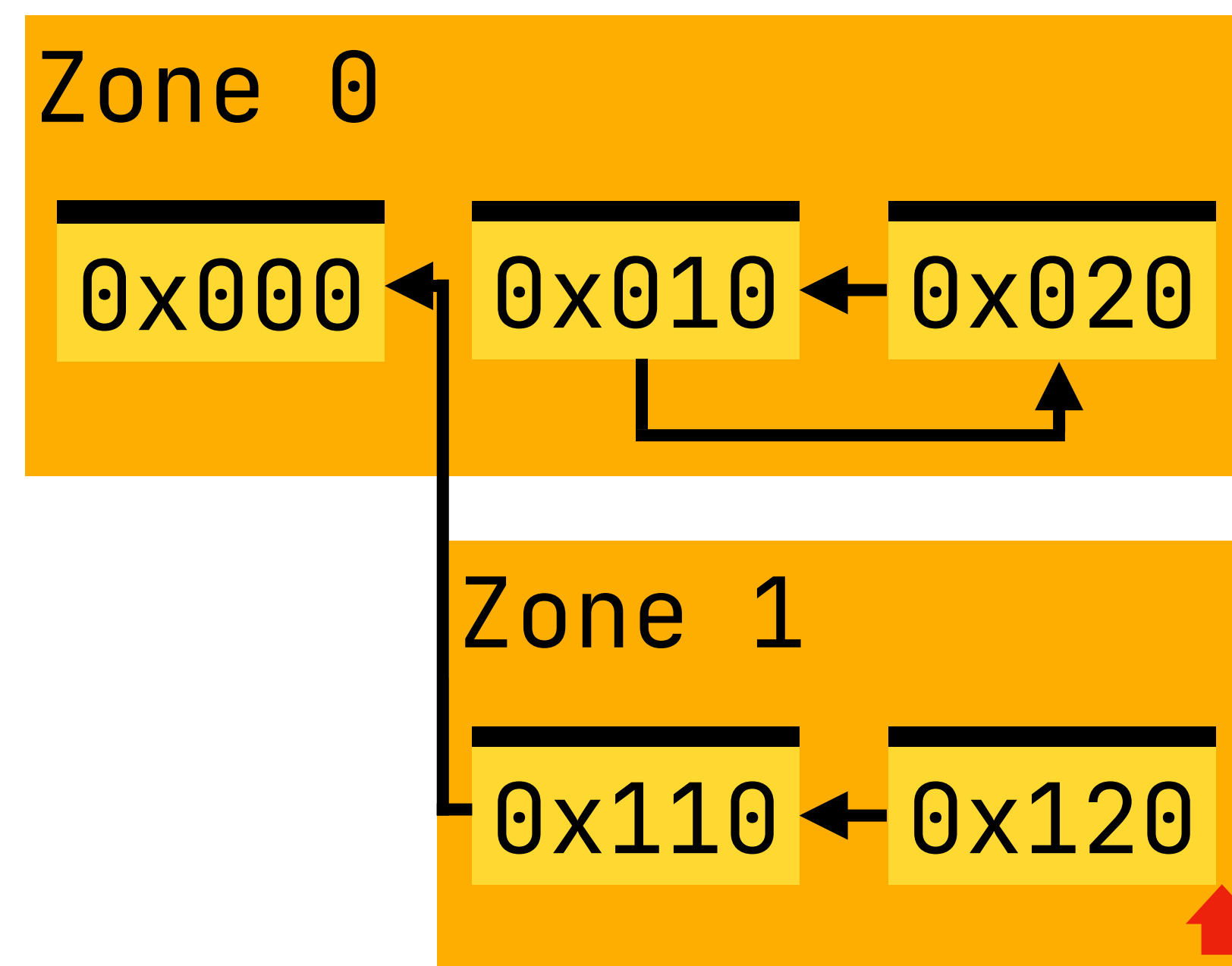
0x73A

Zone number

Handler ID

Stack Offset

- A software-level memory management unit (**sMMU**) is in charge of memory allocation and address translation



raise 0x020()

sMMU:

- Start from Zone 1, find handler with ID 2

Computed from sp

Computed from handler address

Multi Lexa

Main idea: virtual resumption

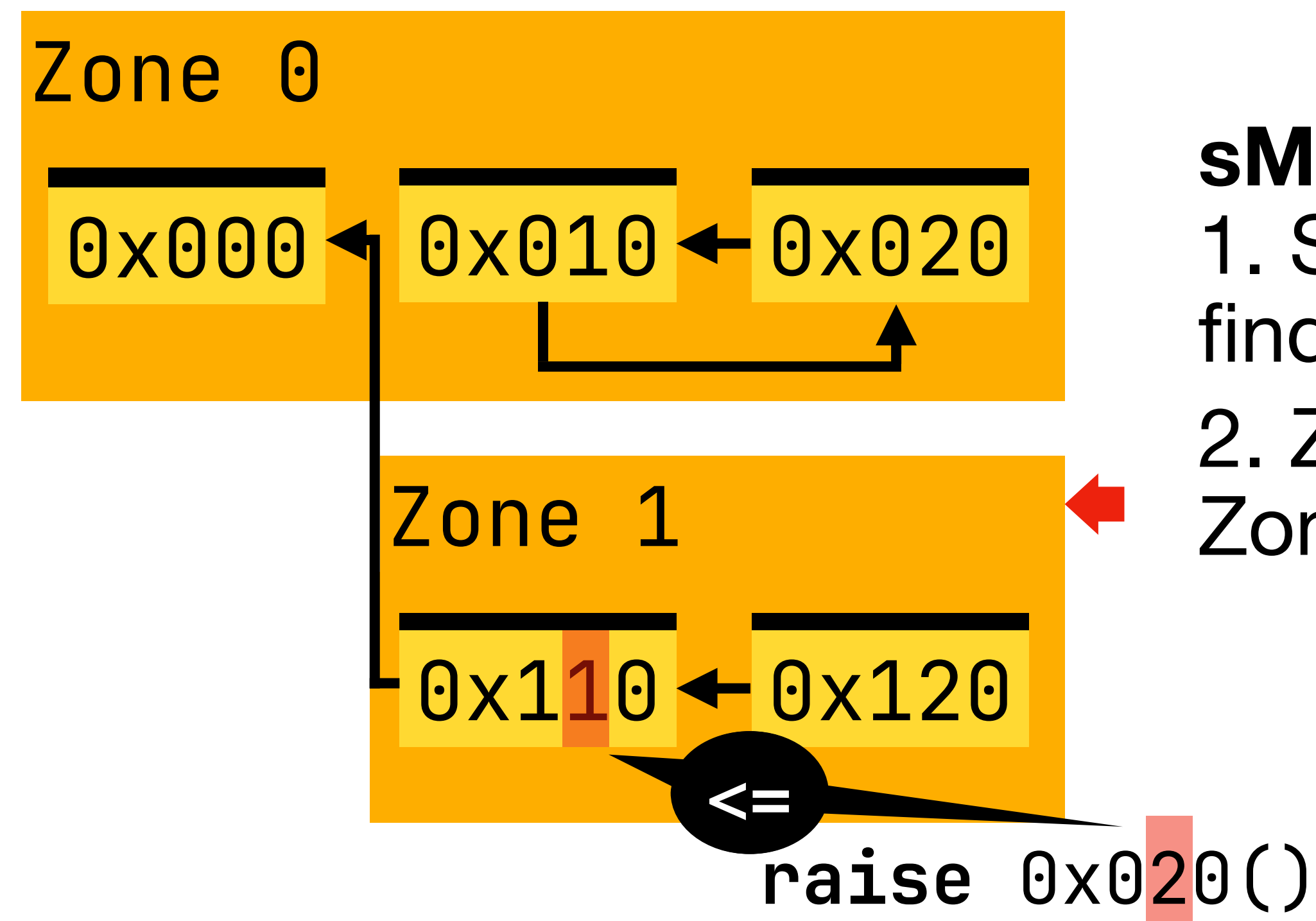
0x73A

Zone number

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- A software-level memory management unit(**sMMU**) is in charge of memory allocation and address translation



sMMU:

1. Start from Zone 1, find handler with ID 2
2. Zone walk: check if Zone 1 has Handler 2

Multi Lexa

Main idea: virtual resumption

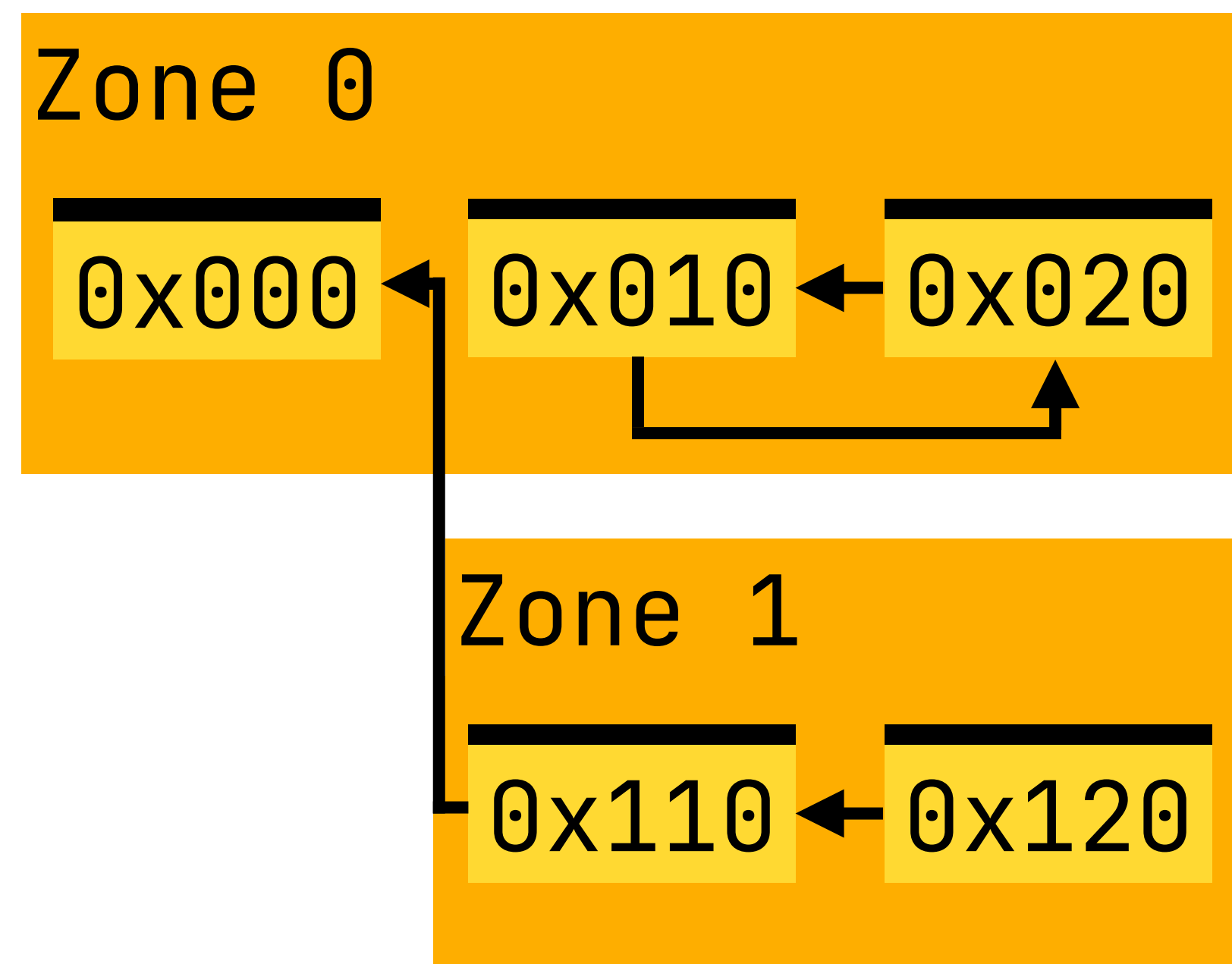
0x73A

Zone number

Handler ID

Stack Offset

- A software-level memory management unit(**sMMU**) is in charge of memory allocation and address translation



sMMU:

1. Start from Zone 1, find handler with ID 2
2. Zone walk: check if Zone 1 has Handler 2
3. Yes. Modify the zone number in virtual address to the target zone.

0x020 → 0x120

raise 0x020()

Multi Lexa

Main idea: virtual resumption

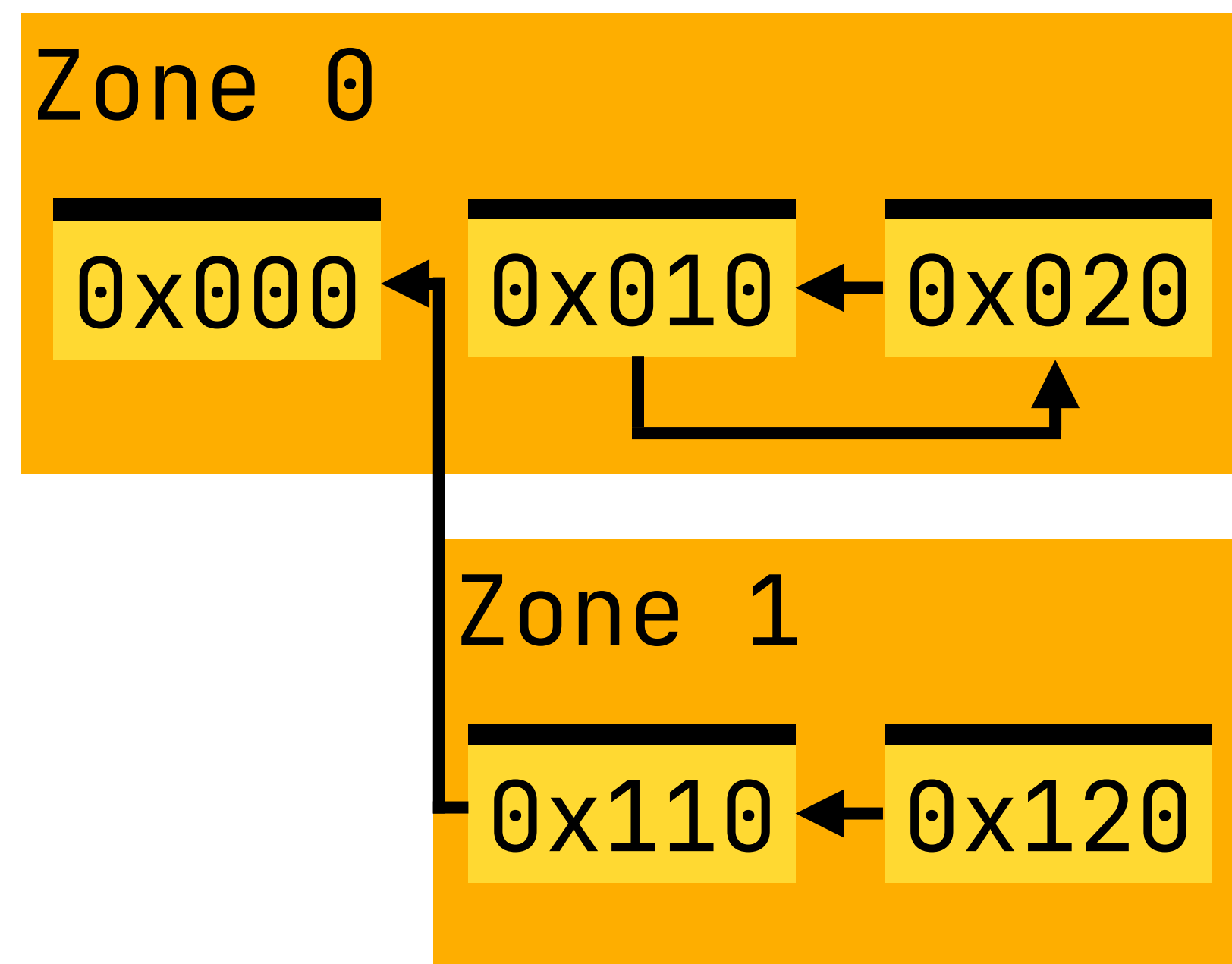
0x73A

Zone number

Handler ID

Stack Offset

- A software-level memory management unit(**sMMU**) is in charge of memory allocation and address translation



sMMU:

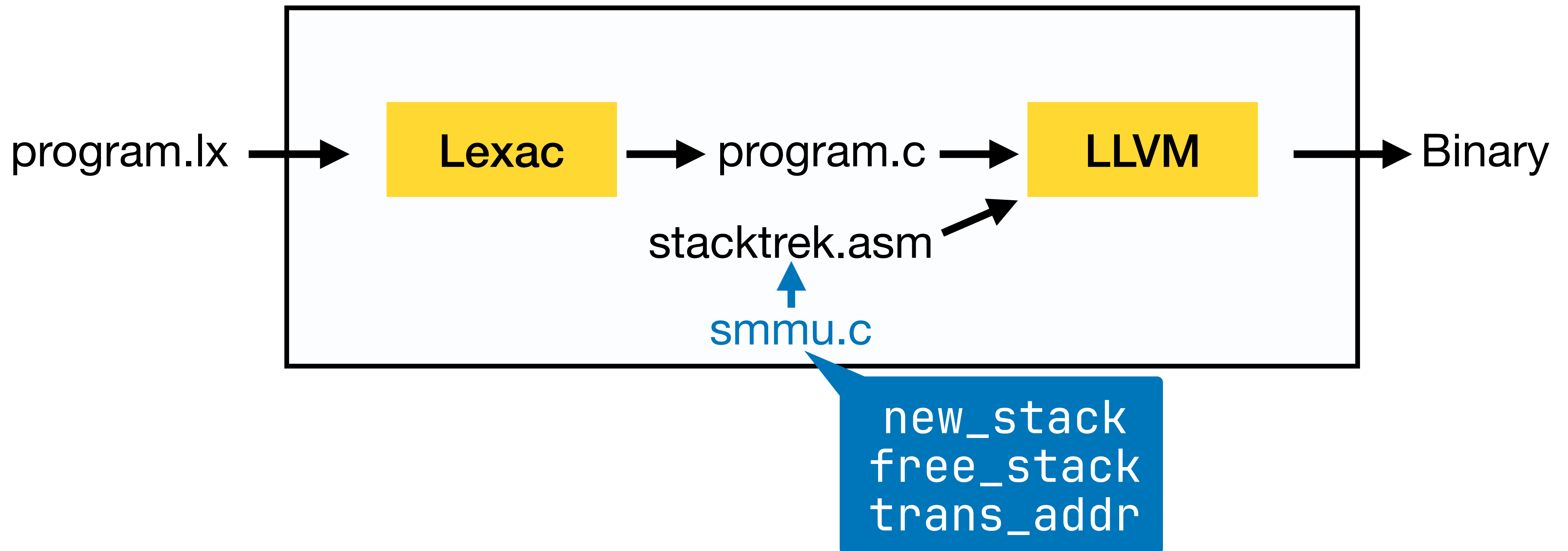
1. Start from Zone 1, find handler with ID 2
2. Zone walk: check if Zone 1 has Handler 2
3. Yes. Modify the zone number in virtual address to the target zone.
4. Return 0x120

raise 0x020()

Multi Lexa

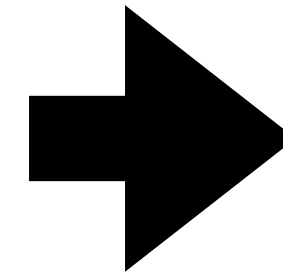
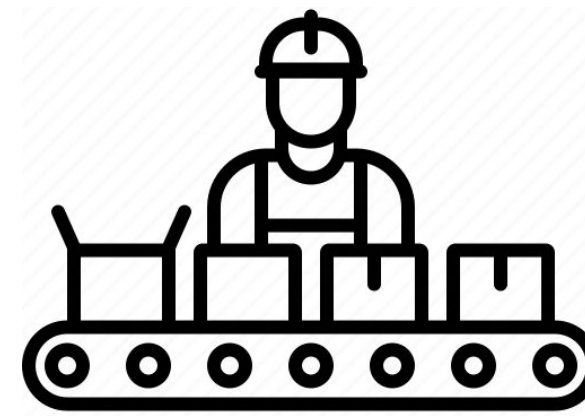
Library sMMU

Lexa Compiler



high-level effect handlers in Lexa

```
handle E with H  
  
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resume ...
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low-level stack switching in assembly

```
ENTER  
  
RAISE  
  
RESUME
```

Direct Lexa

A large, hollow black triangle pointing upwards. Inside the triangle, the text "tradeoffs between performance and expressivity" is centered.

tradeoffs
between
performance
and expressivity

Multi Lexa

Zero Lexa

Effect handler

A powerful language feature

Language Definition

Syntax

`handle E with H`

`raise ...`

`resume ...`

+

active research

Semantics

`Install a handler`

`Find handler; Capture continuation`

`Resume continuation`

A glimpse of two semantics

Lexical Semantics

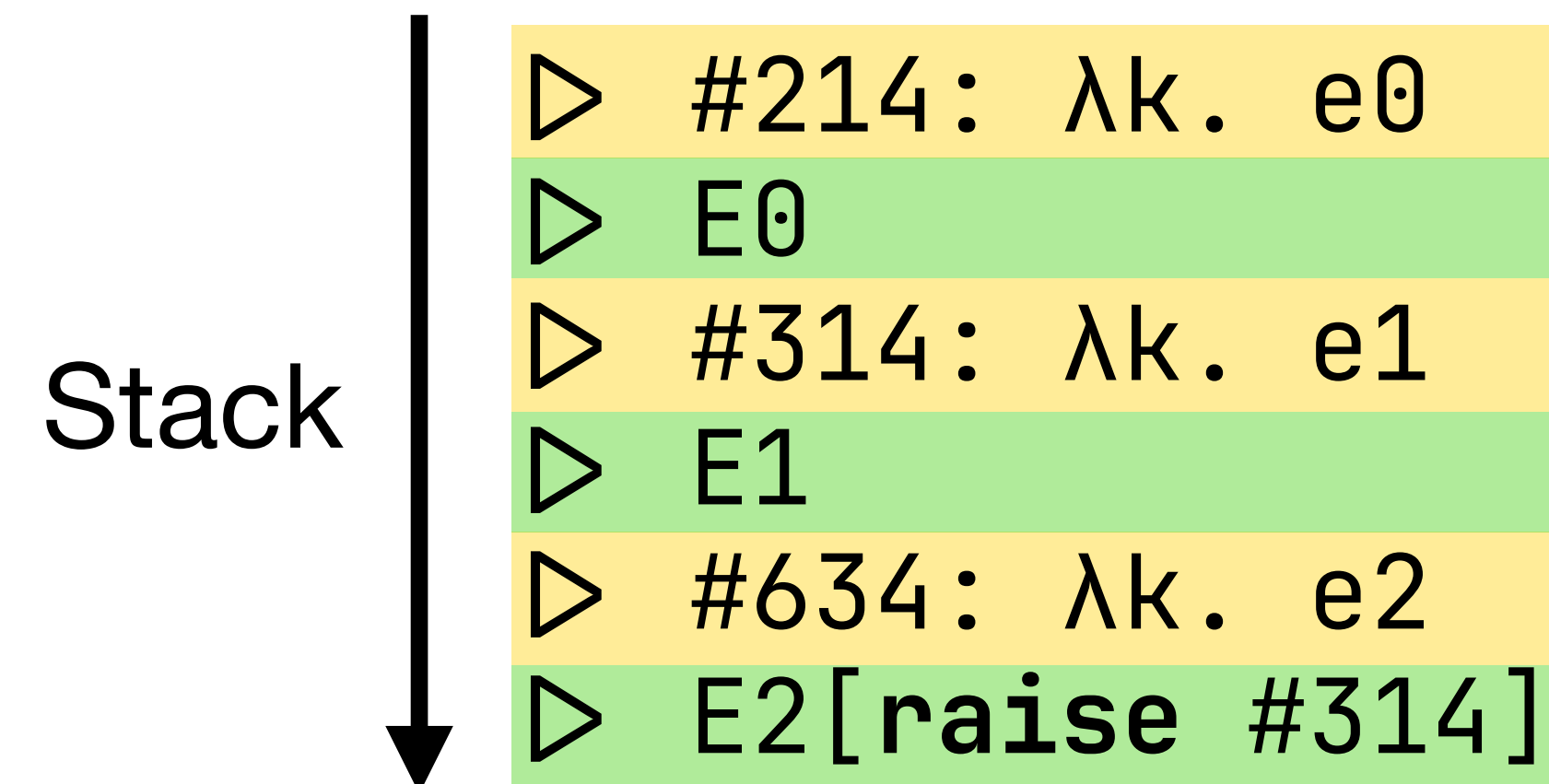
Stack

▷ #214: $\lambda k. e0$
▷ E0
▷ #314: $\lambda k. e1$
▷ E1
▷ #634: $\lambda k. e2$
▷ E2[raise #314]

A glimpse of two semantics

```
def f(n, exception_handler) =  
  ...  
  if (condition)  
    raise exception_handler(...);  
  ...
```

Lexical Semantics



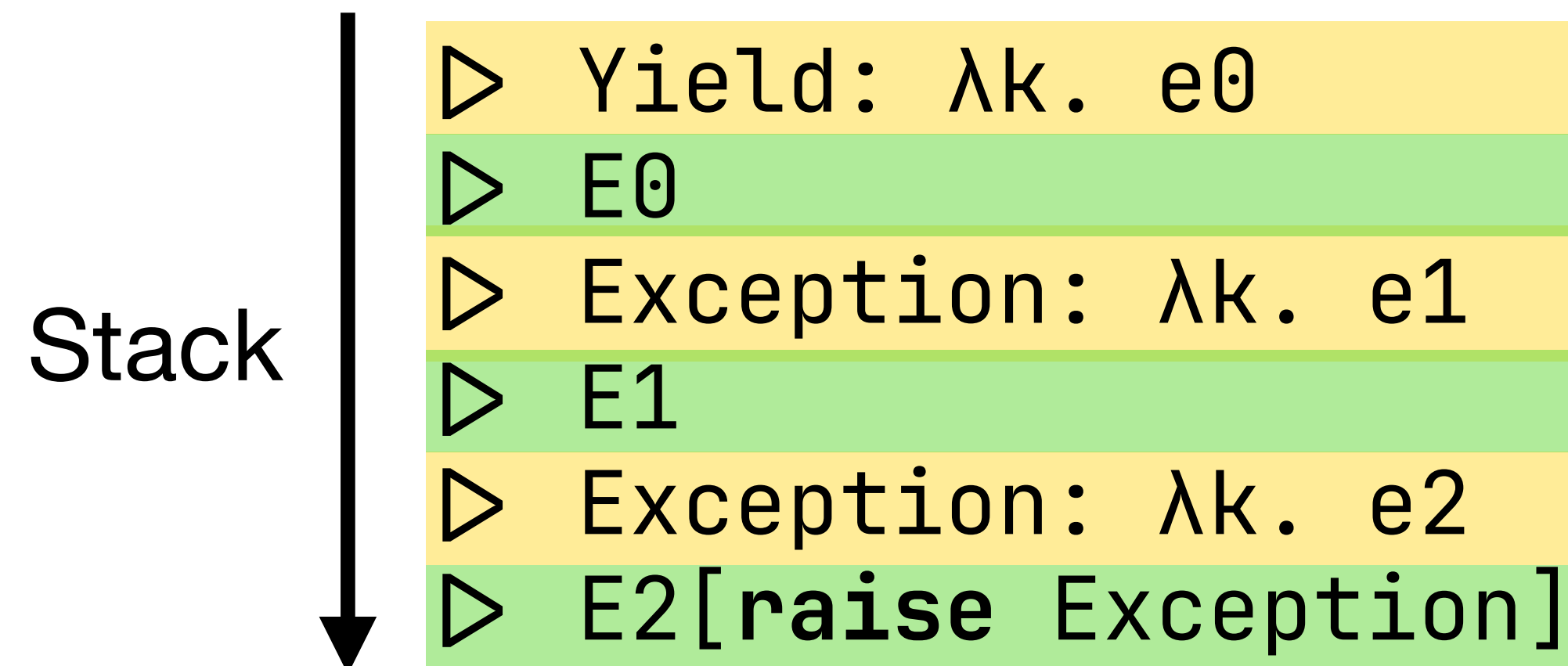
A glimpse of two semantics

```
def f(n) =  
  ...  
  if (condition)  
    raise Exception(...);  
  ...
```

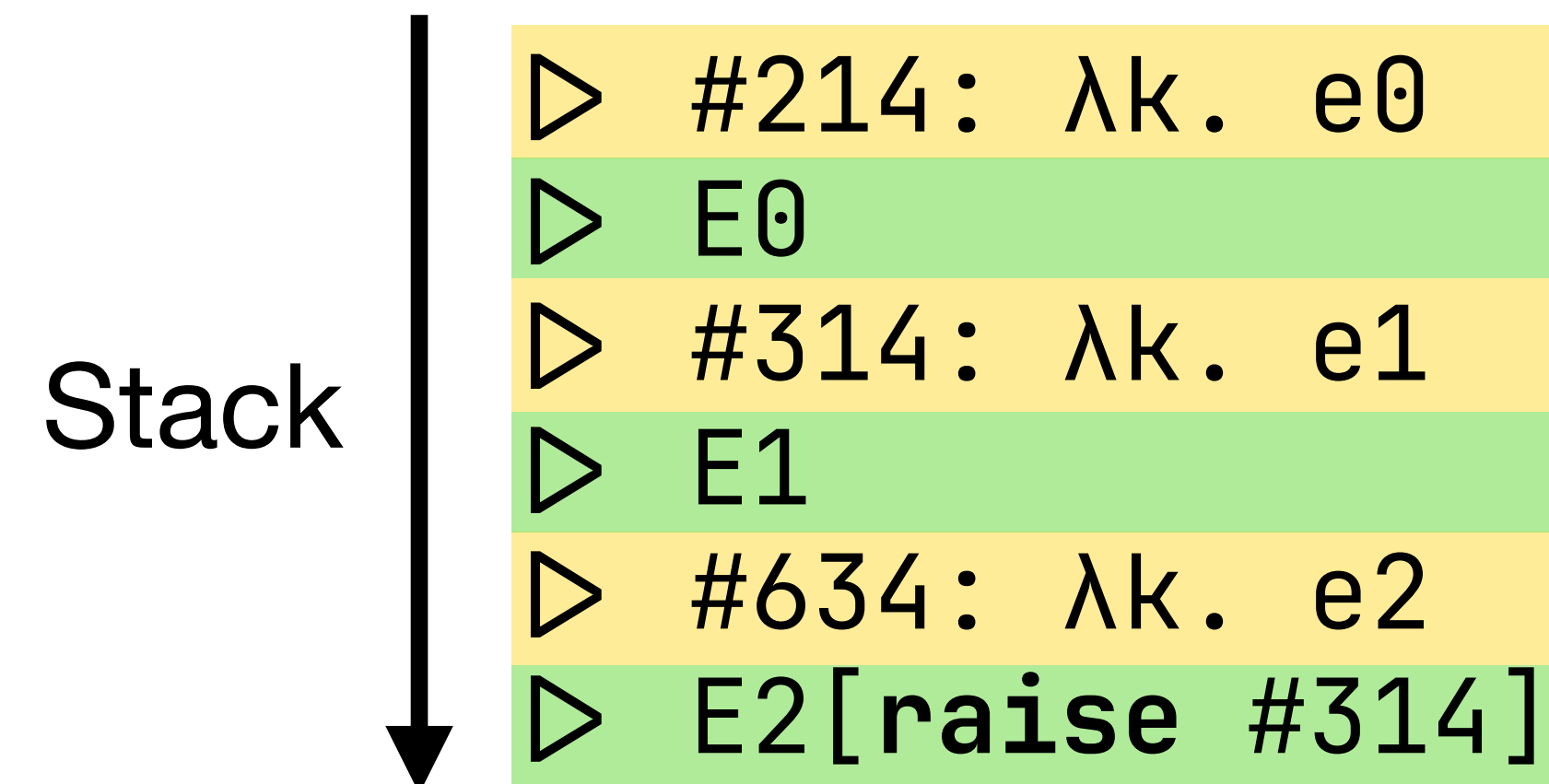
V.S.

```
def f(n, exception_handler) =  
  ...  
  if (condition)  
    raise exception_handler(...);  
  ...
```

Traditional Semantics



Lexical Semantics



Problem

All existing implementations impose a runtime cost

extra
parameters

```
def f(n) =  
  ...  
  if (condition)  
    raise Exception(...);  
  ...
```

Traditional Semantics

V.S.

```
def f(n, exception_handler) =  
  ...  
  if (condition)  
    raise exception_handler(...);  
  ...
```

Lexical Semantics

Problem

All existing implementations impose a runtime cost

extra
parameters

```
def f(n) =  
  ...  
  if (condition)  
    raise Exception(...);  
  ...
```

V.S.

```
def f(n, exception_handler) =  
  ...  
  if (condition)  
    raise exception_handler(...);  
  ...
```

Traditional Semantics

Lexical Semantics

Programmers are used to the
zero-overhead handlers.

Problem

All existing implementations impose a runtime cost

extra parameters

```
def f(n) =  
  ...  
  if (condition)  
    raise Exception(...);  
  ...
```

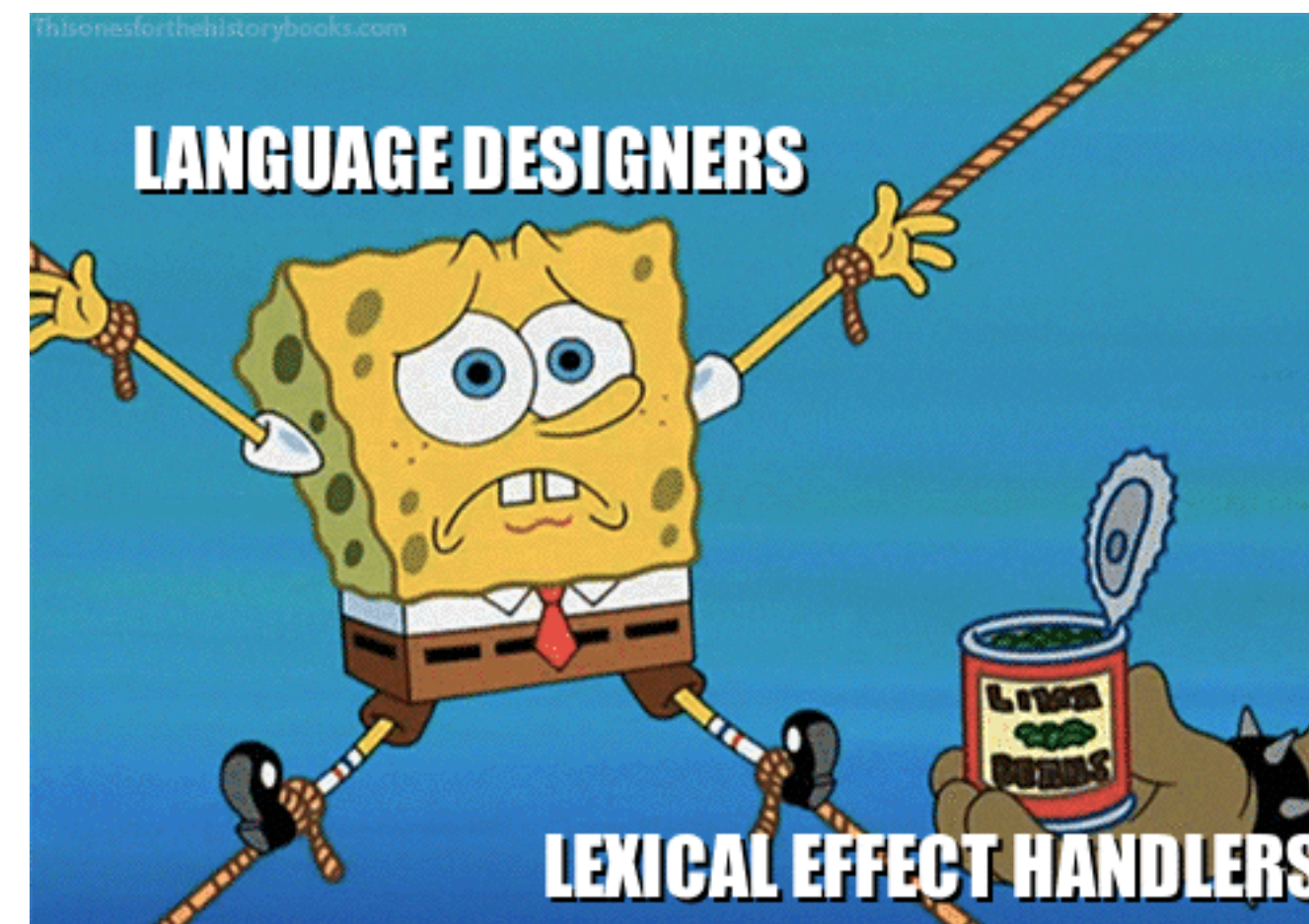
V.S.

```
def f(n, exception_handler) =  
  ...  
  if (condition)  
    raise exception_handler(...);  
  ...
```

Traditional Semantics

Lexical Semantics

Programmers are used to the zero-overhead handlers.



It's a hard sell.

Zero Lexa

Zero-overhead Principle

“What you don’t use, you don’t pay for.”

- Bjarne Stroustrup [1995, 2012]

“Normal case execution efficiency should not be impaired at all.”

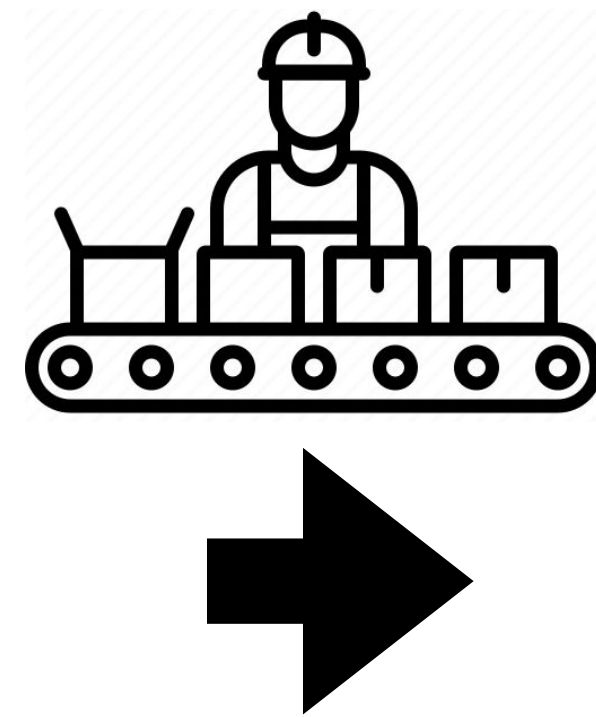
- Atkinson, Liskov, and Scheifler [1978]

- Liskov and Snyder [1979]

Zero Lexa

Zero-overhead Lexical effect handler

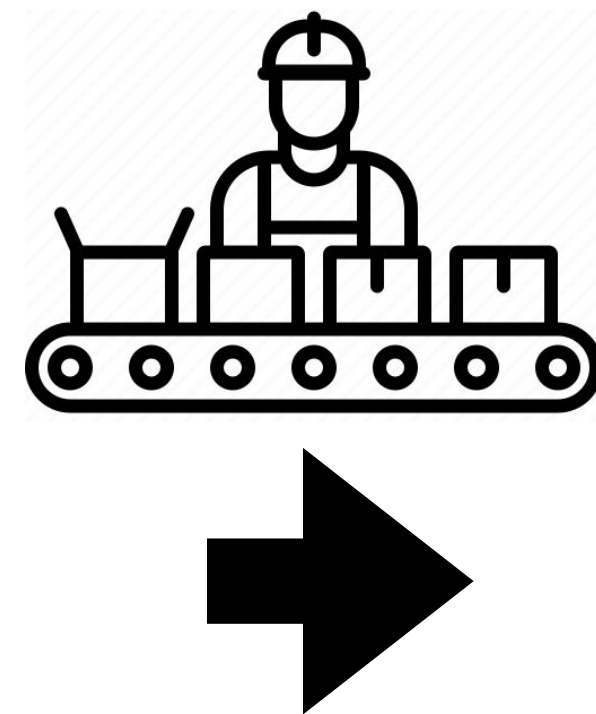
```
def f(n, handler) =  
  ...  
  if (condition)  
    raise handler  
  ...
```



Zero Lexa

Zero-overhead Lexical effect handler

```
def f(n, handler) =  
  ...  
  if (condition)  
    raise handler  
  ...
```



```
def f(n) =  
  ...  
  if (condition)  
    raise _  
  ...
```

Compile away handler passing, yet
retain the same behavior

Zero Lexa

Zero-overhead Lexical effect handler

We will first see a non-zero-overhead execution.

We will then figure out how to make it zero-overhead.

Zero Lexa, Example 1

```
val parked = Null
def work() {
  handle
    while true
      DO_SOMETHING
      raise yield()
  with yield(k) =
    val peer = parked
    parked = k
    if peer then
      resume peer ()
}
work()
work()
```

Zero Lexa, Example 1

```
def work() {  
  handle  
    DO_SOMETHING  
  with yield = ...  
}
```

Zero Lexa, Example 1

```
def work() {  
  handle  
  foo(yield)  
  with yield = ...  
}
```

Zero Lexa, Example 1

```
def work() {  
  handle  
    foo(yield)  
  with yield = ...  
}
```

```
def foo(yield):  
  handle  
    bar(yield, exc)  
  with exc = ...
```

Zero Lexa, Example 1

```
def work() {  
  handle  
    foo(yield)  
  with yield = ...  
}
```

```
def foo(yield):  
  handle  
    bar(yield, exc)  
  with exc = ...
```

```
def bar(yield, exc):  
  raise yield()  
  raise exc()
```

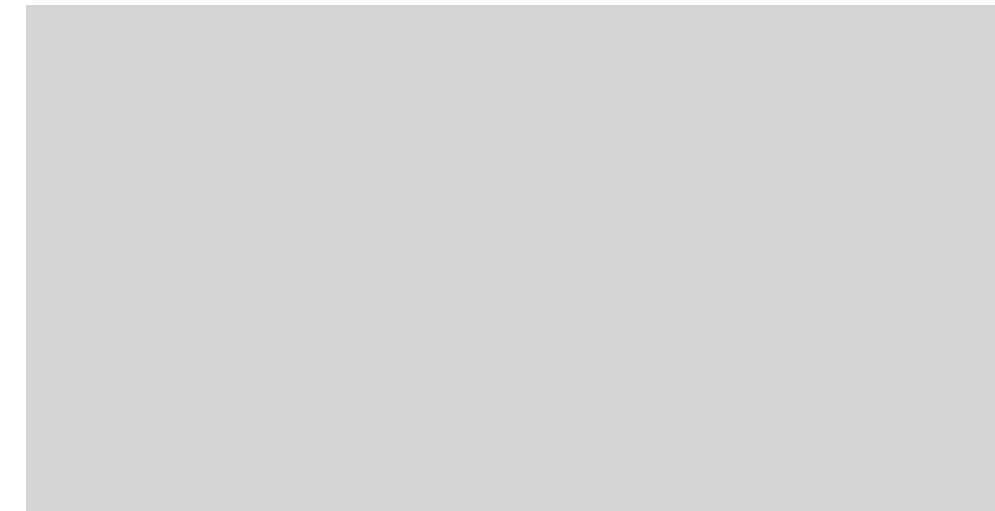
Zero Lexa, Example 1

```
def work() {  
  handle  
    foo(yield)  
  with yield = ...  
}
```

```
def foo(yield):  
  handle  
    bar(yield, exc)  
  with exc = ...
```

```
def bar(yield, exc):  
  raise yield()  
  raise exc()
```

work:



Zero Lexa, Example 1

```
def work() {  
  handle  
    foo(yield)  
  with yield = ...  
}
```

```
def foo(yield):  
  handle  
    bar(yield, exc)  
  with exc = ...
```

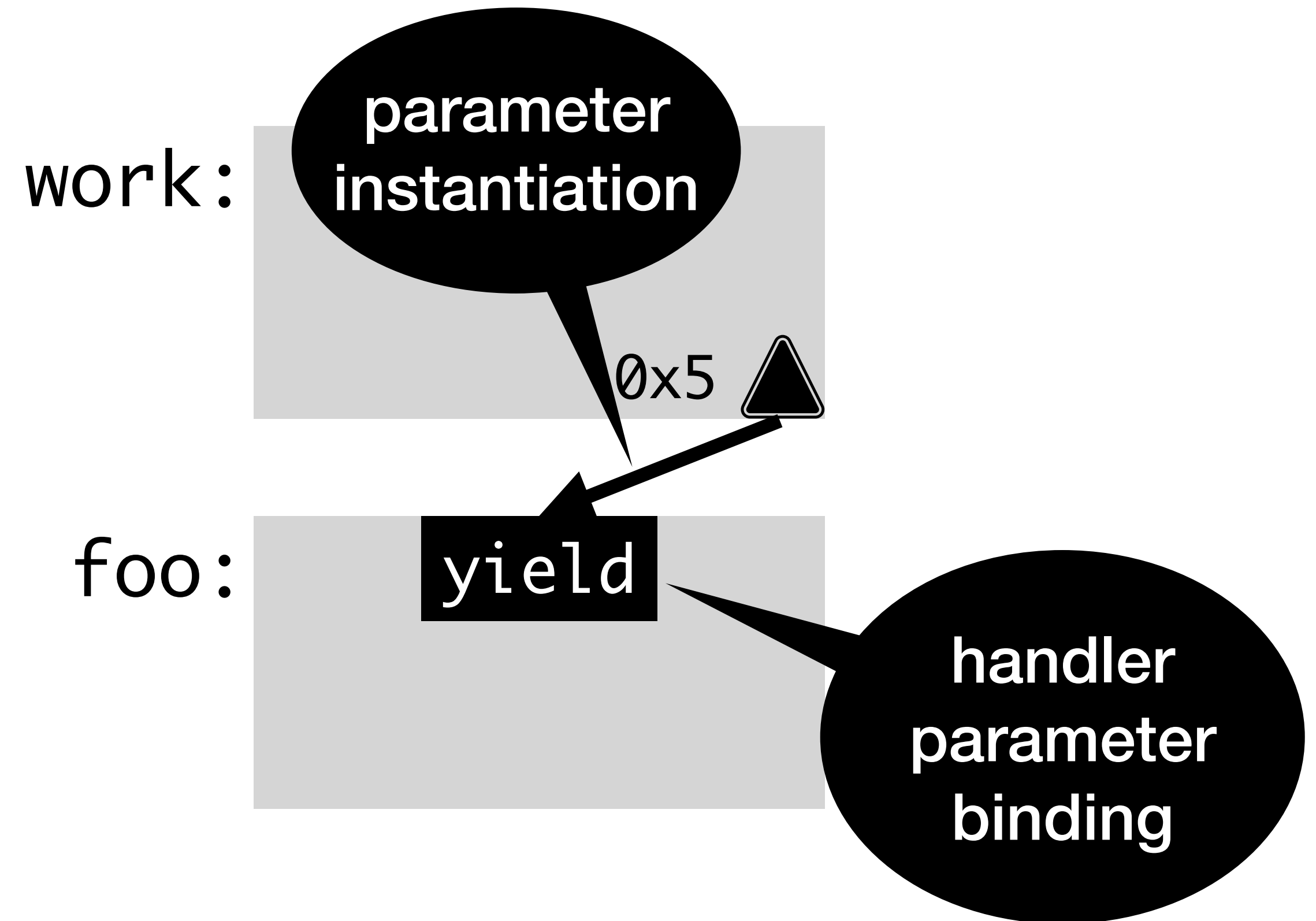
```
def bar(yield, exc):  
  raise yield()  
  raise exc()
```

work:



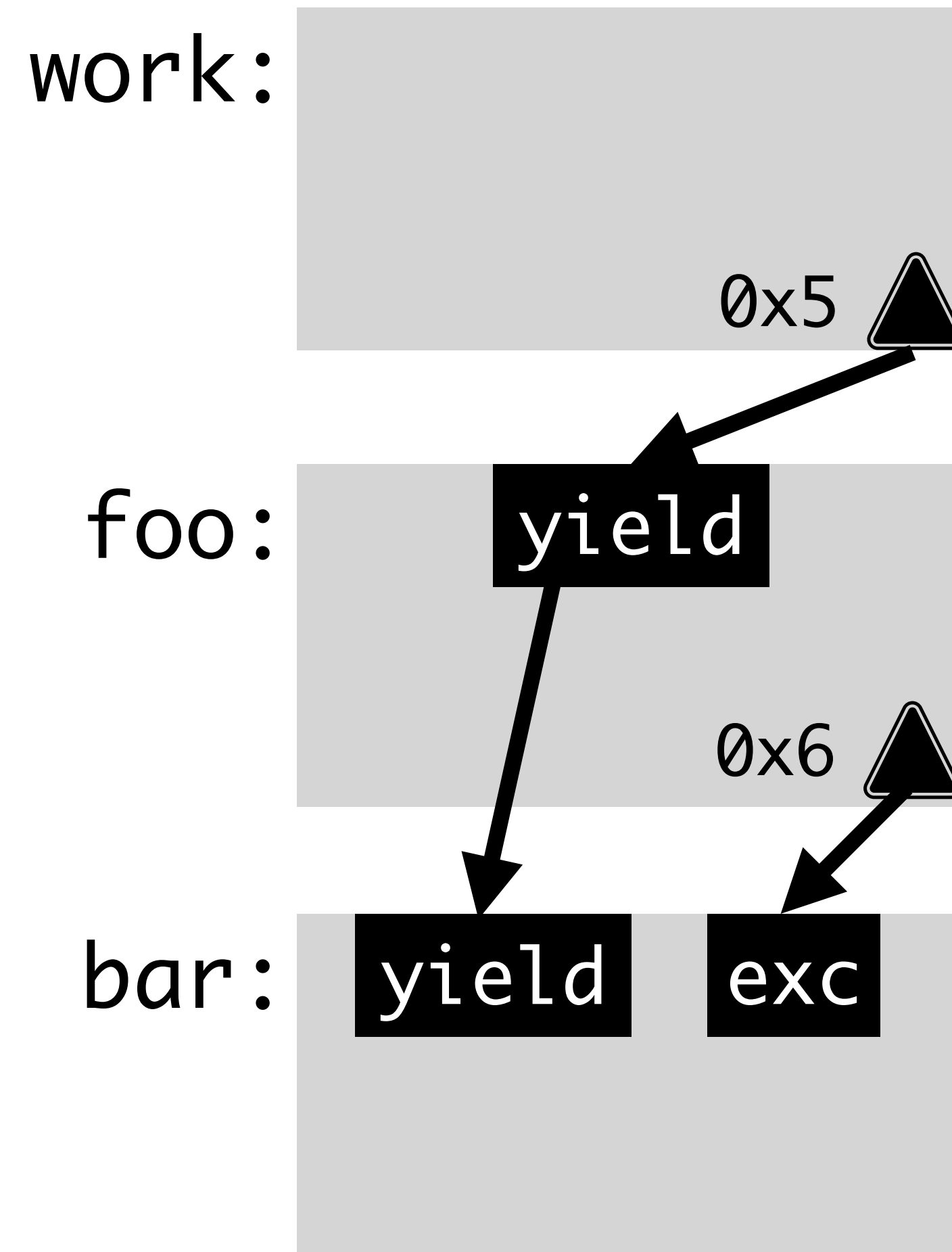
Zero Lexa, Example 1

```
def work() {  
  handle  
  foo(yield)  
  with yield = ...  
}  
  
def foo(yield):  
  handle  
  bar(yield, exc)  
  with exc = ...  
  
def bar(yield, exc):  
  raise yield()  
  raise exc()
```



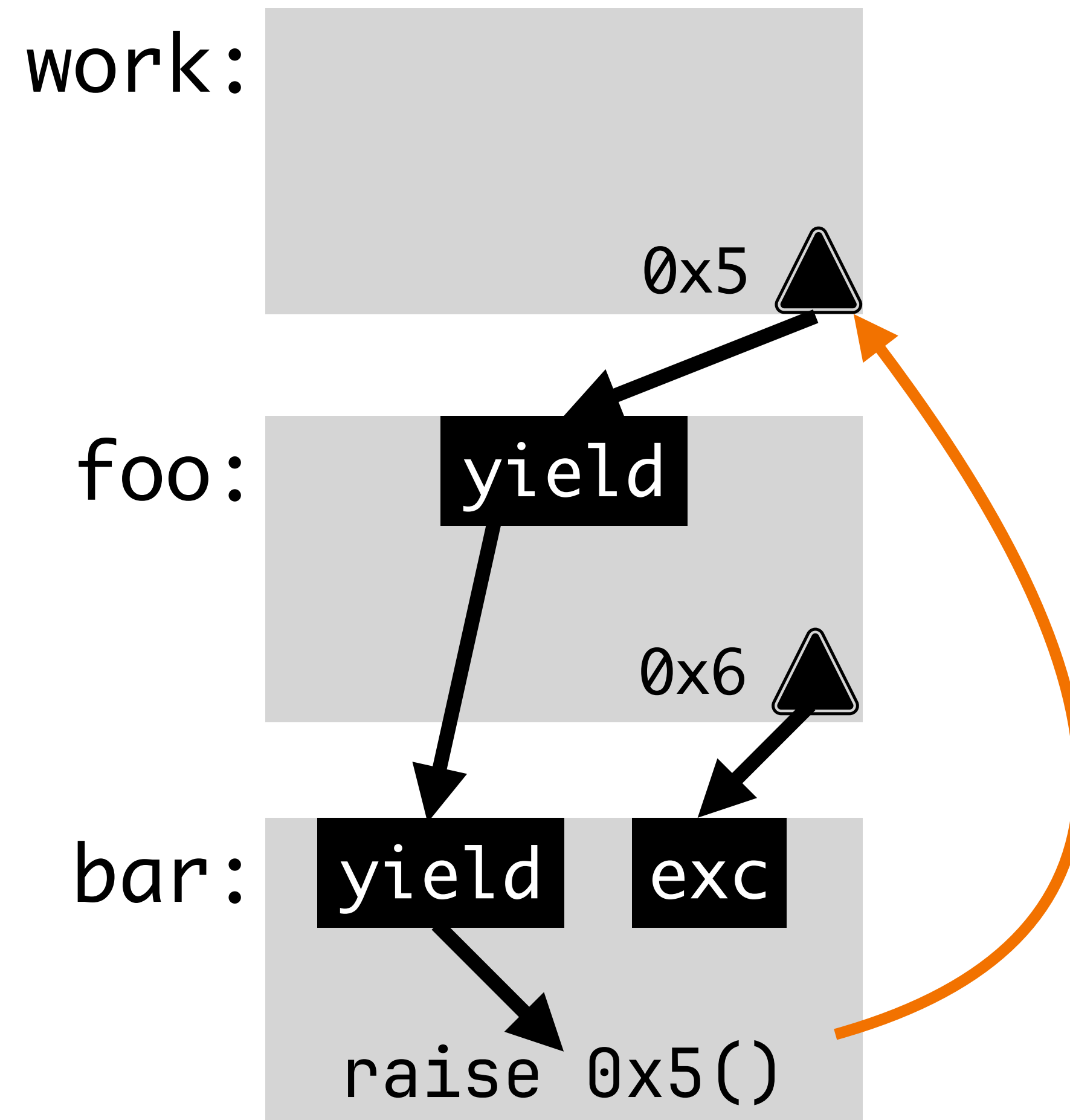
Zero Lexa, Example 1

```
def work() {  
  handle  
  foo(yield)  
  with yield = ...  
}  
  
def foo(yield):  
  handle  
  bar(yield, exc)  
  with exc = ...  
  
def bar(yield, exc):  
  raise yield()  
  raise exc()
```



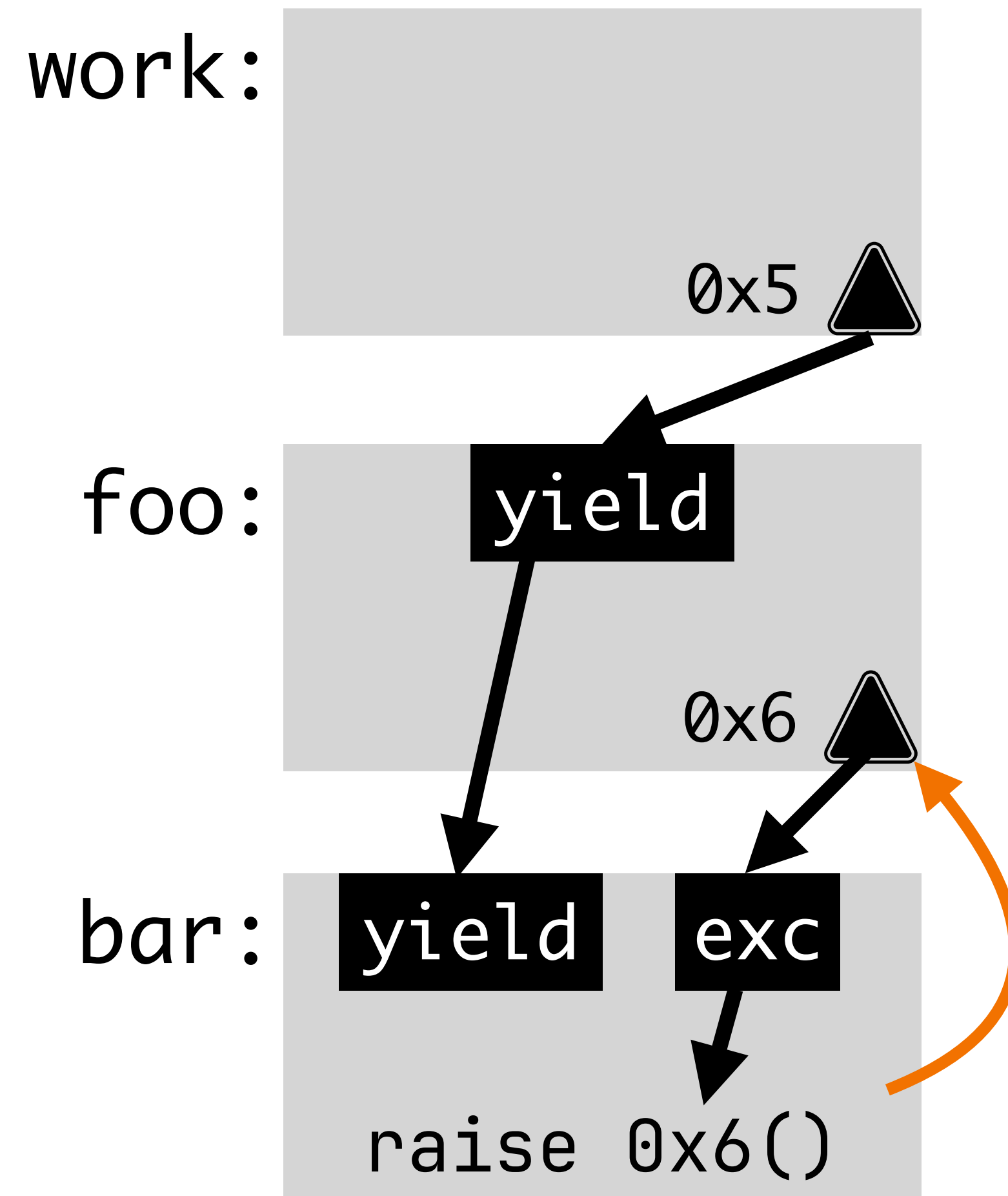
Zero Lexa, Example 1

```
def work() {  
  handle  
  foo(yield)  
  with yield = ...  
}  
  
def foo(yield):  
  handle  
  bar(yield, exc)  
  with exc = ...  
  
def bar(yield, exc):  
  raise yield()  
  raise exc()
```



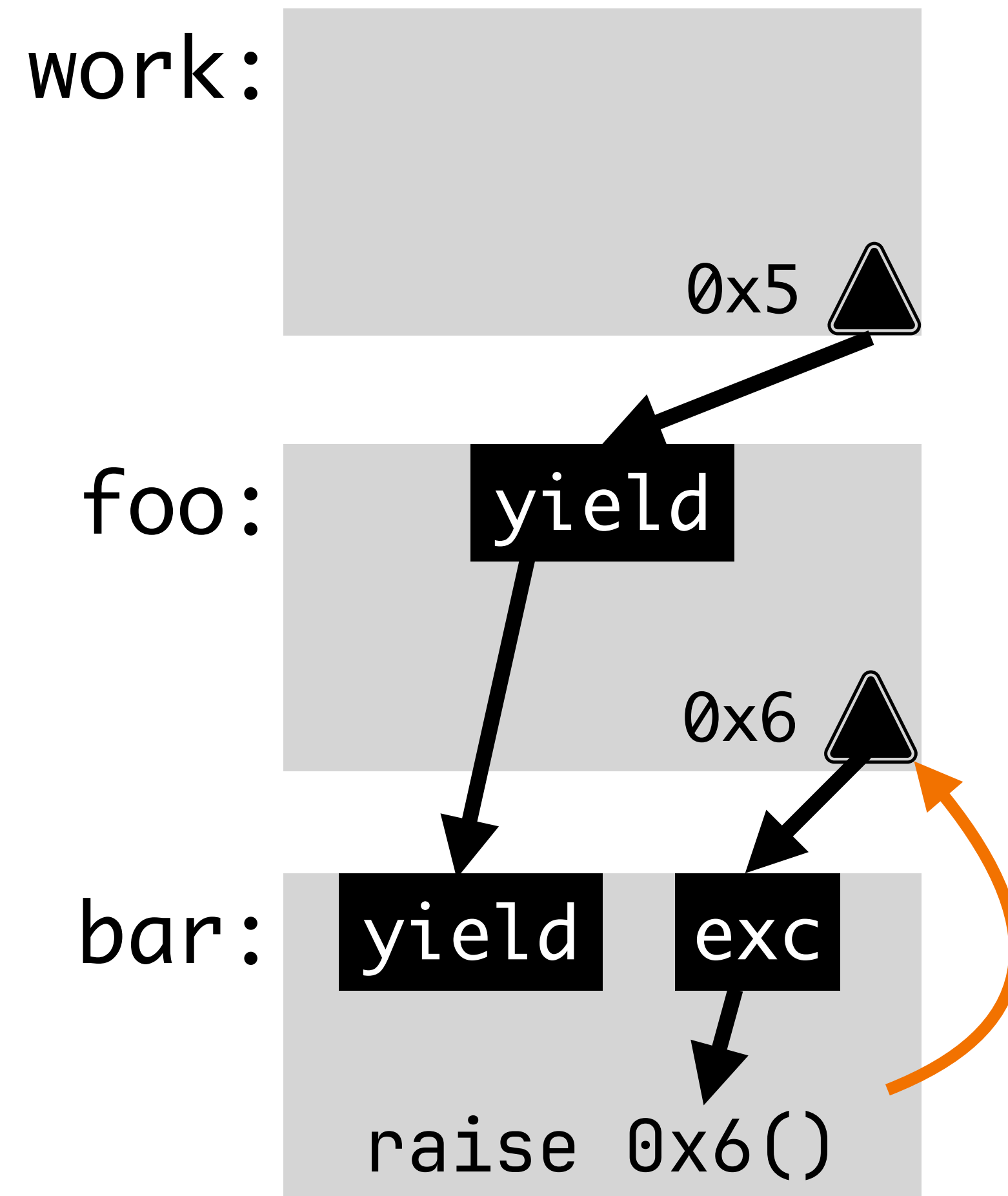
Zero Lexa, Example 1

```
def work() {  
  handle  
  foo(yield)  
  with yield = ...  
}  
  
def foo(yield):  
  handle  
  bar(yield, exc)  
  with exc = ...  
  
def bar(yield, exc):  
  raise yield()  
  raise exc()
```



Zero Lexa, Example 1

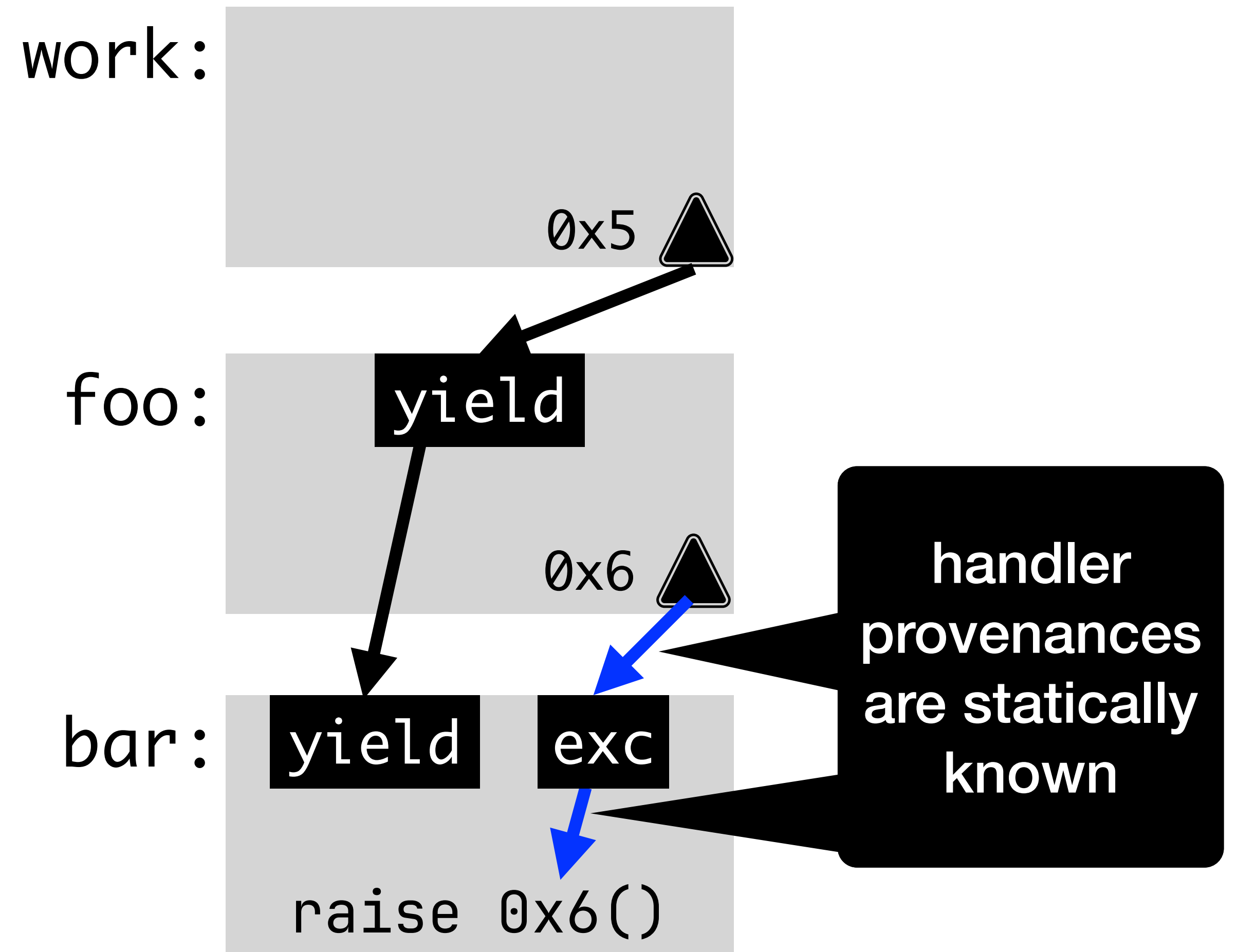
```
def work() {  
  handle  
  foo(yield)  
  with yield = ...  
}  
  
def foo(yield):  
  handle  
  bar(yield, exc)  
  with exc = ...  
  
def bar(yield, exc):  
  raise yield()  
  raise exc()
```



Can we locate the intended handler without passing down addresses?

Zero Lexa, Example 1

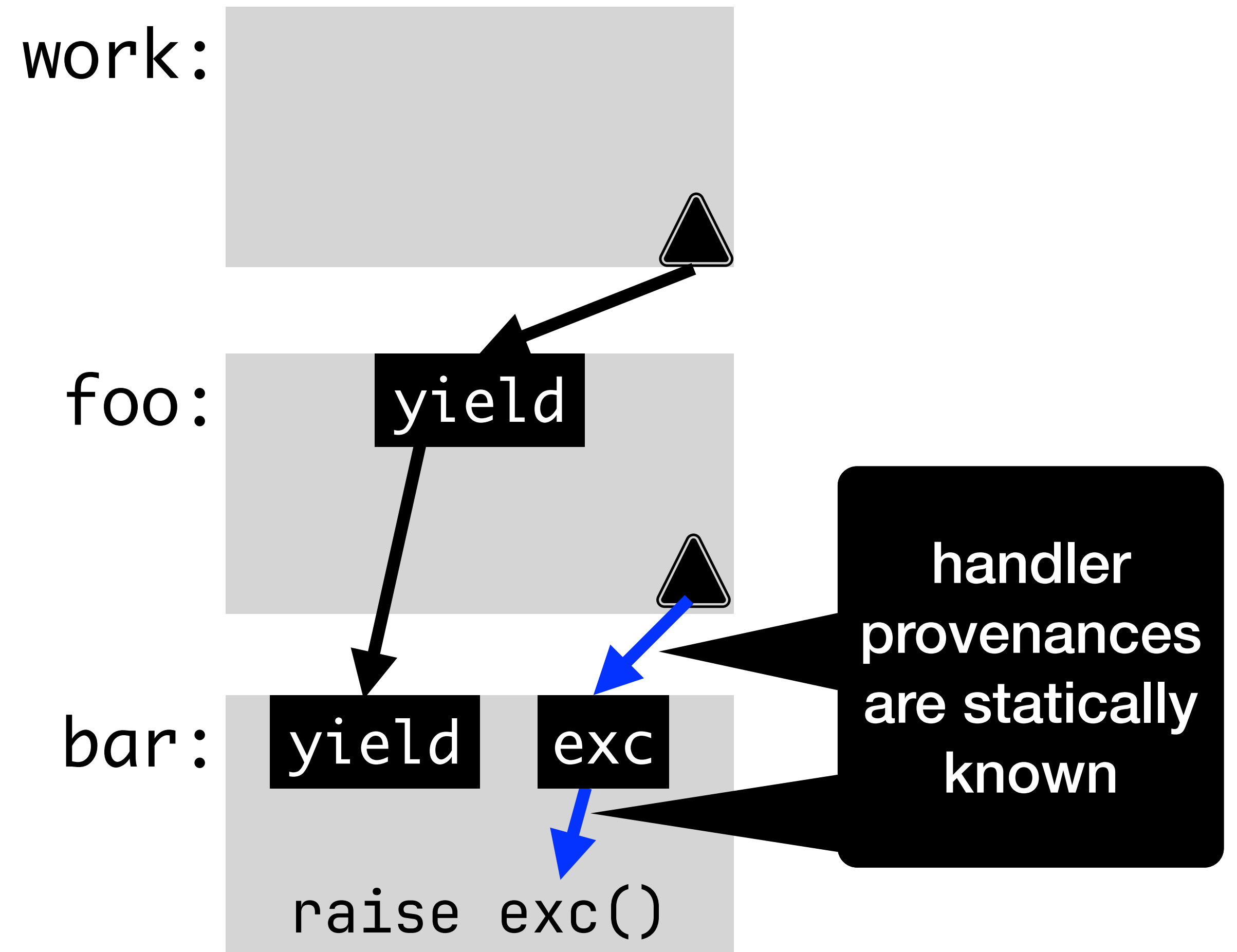
```
def work() {  
  handle  
  foo(yield)  
  with yield = ...  
}  
  
def foo(yield):  
  handle  
  bar(yield, exc)  
  with exc = ...  
  
def bar(yield, exc):  
  raise yield()  
  raise exc()
```



Can we locate the intended handler without passing down addresses?

Zero Lexa, Example 1

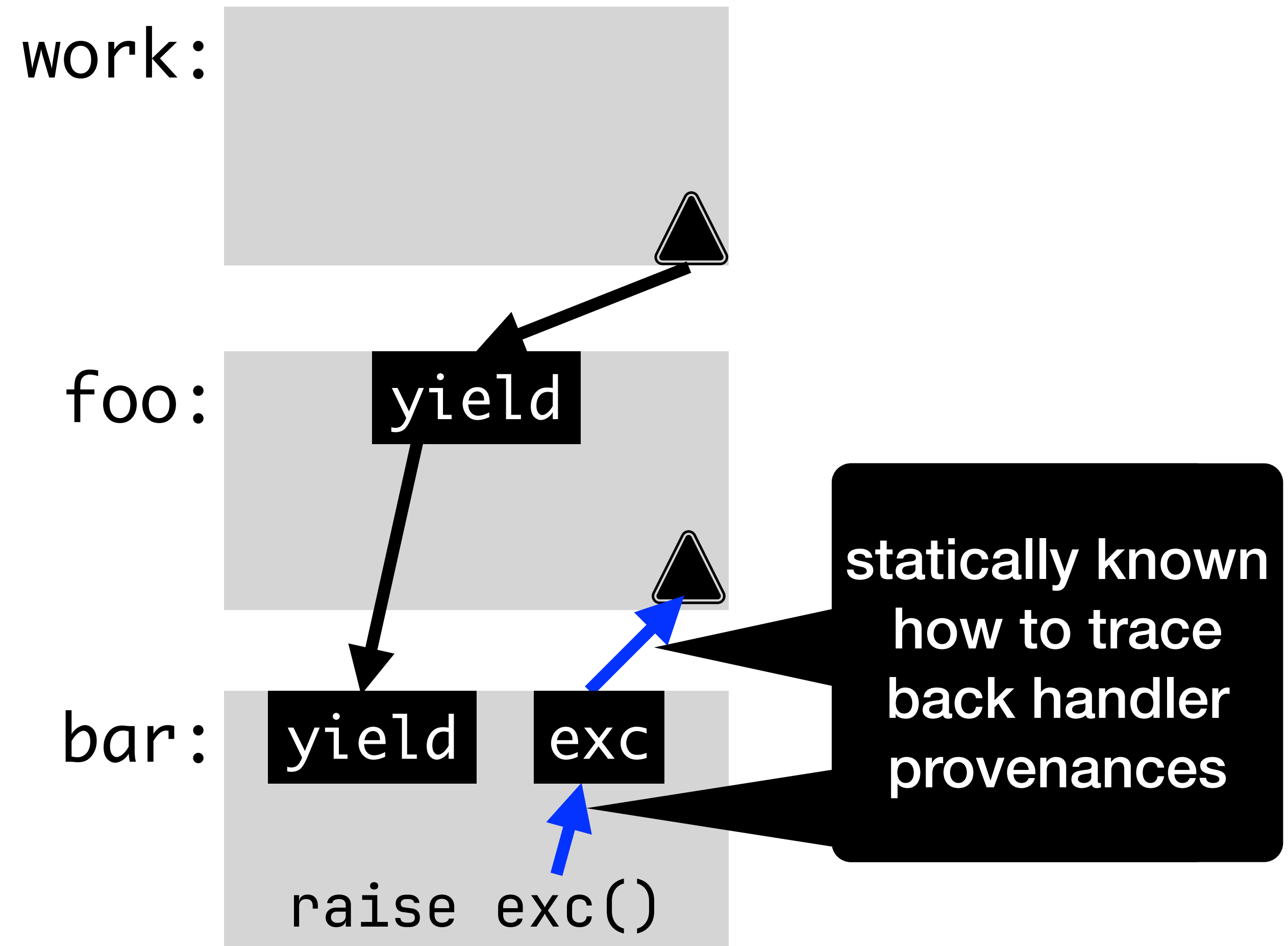
```
def work() {  
  handle  
  foo(yield)  
  with yield = ...  
}  
  
def foo(yield):  
  handle  
  bar(yield, exc)  
  with exc = ...  
  
def bar(yield, exc):  
  raise yield()  
  raise exc()
```



Can we locate the intended handler without passing down addresses?

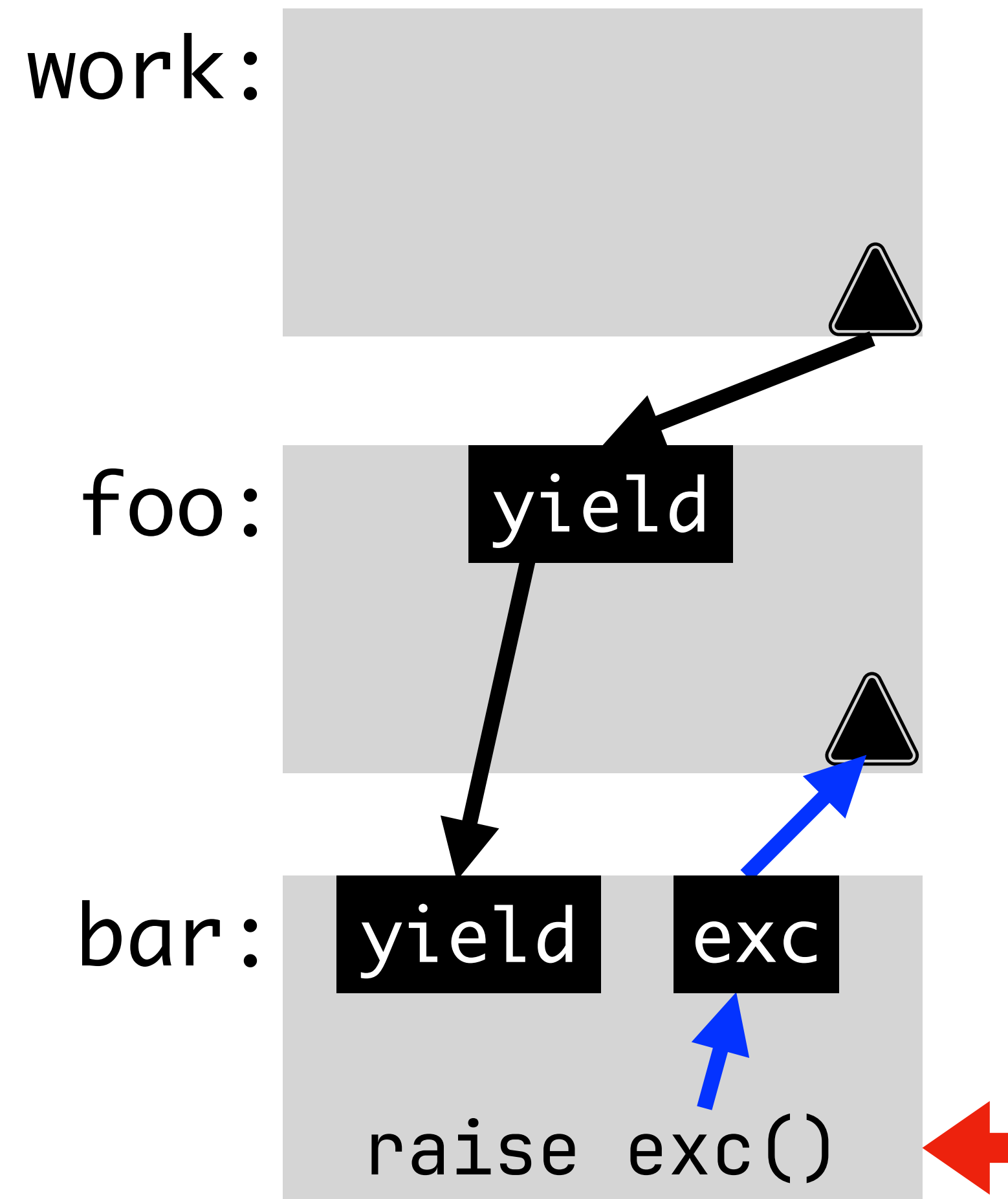
Zero Lexa, Example 1

```
def work() {  
  handle  
  foo(yield)  
  with yield = ...  
}  
  
def foo(yield):  
  handle  
  bar(yield, exc)  
  with exc = ...  
  
def bar(yield, exc):  
  raise yield()  
  raise exc()
```



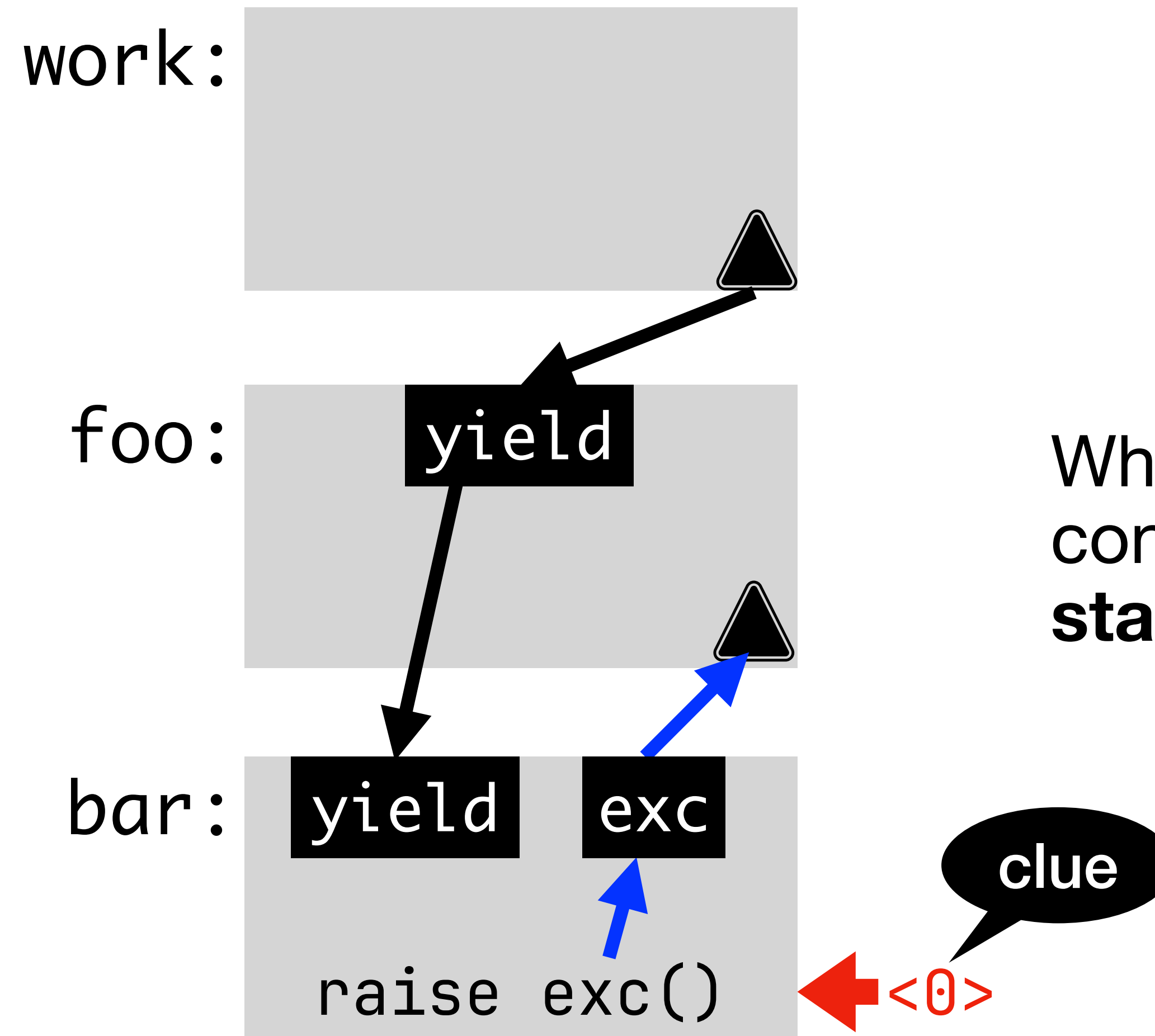
Can we locate the intended handler without passing down addresses?

Zero Lexa, Example 1



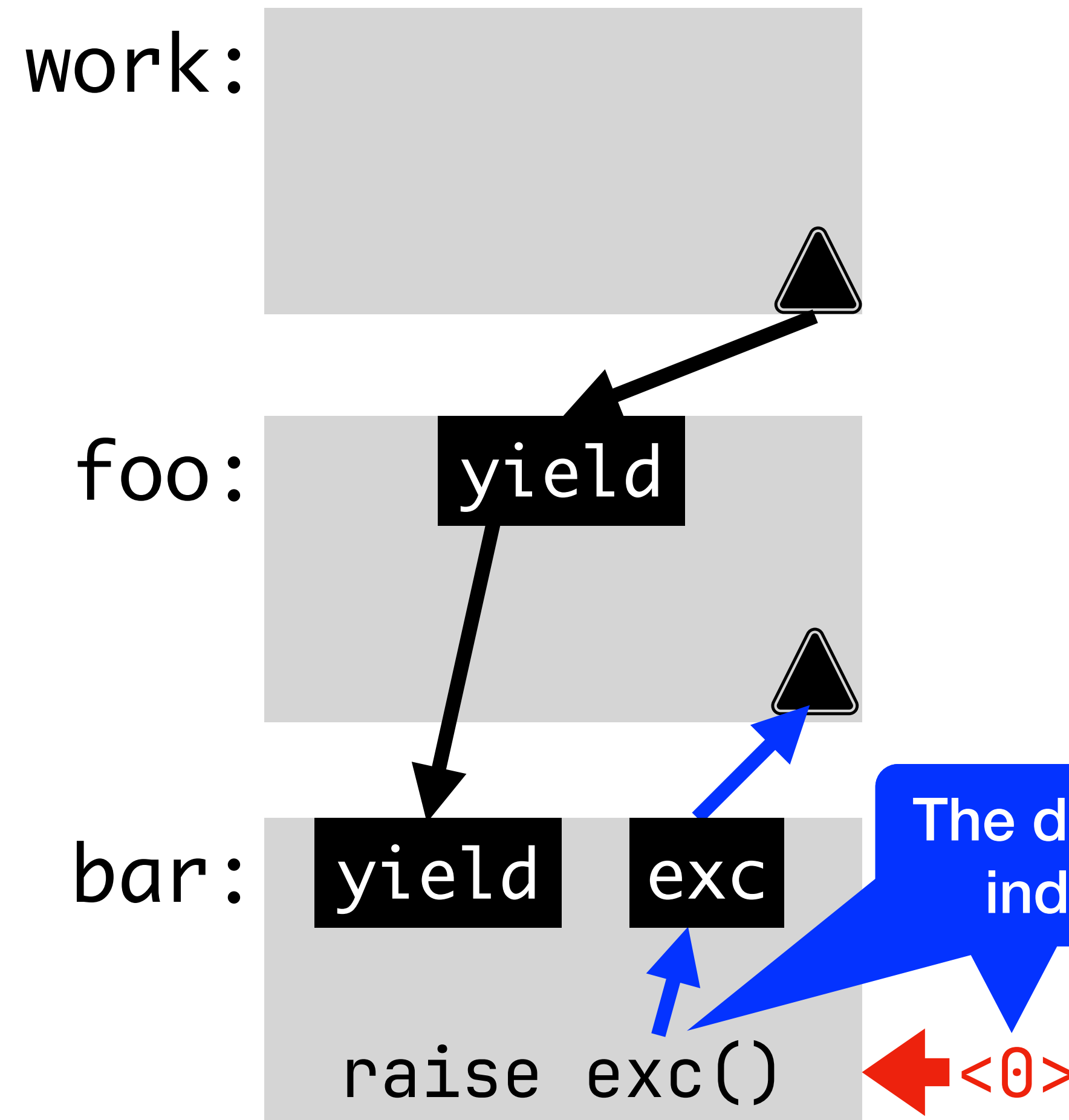
When an effect is raised, control goes to a runtime **stack-walker**.

Zero Lexa, Example 1



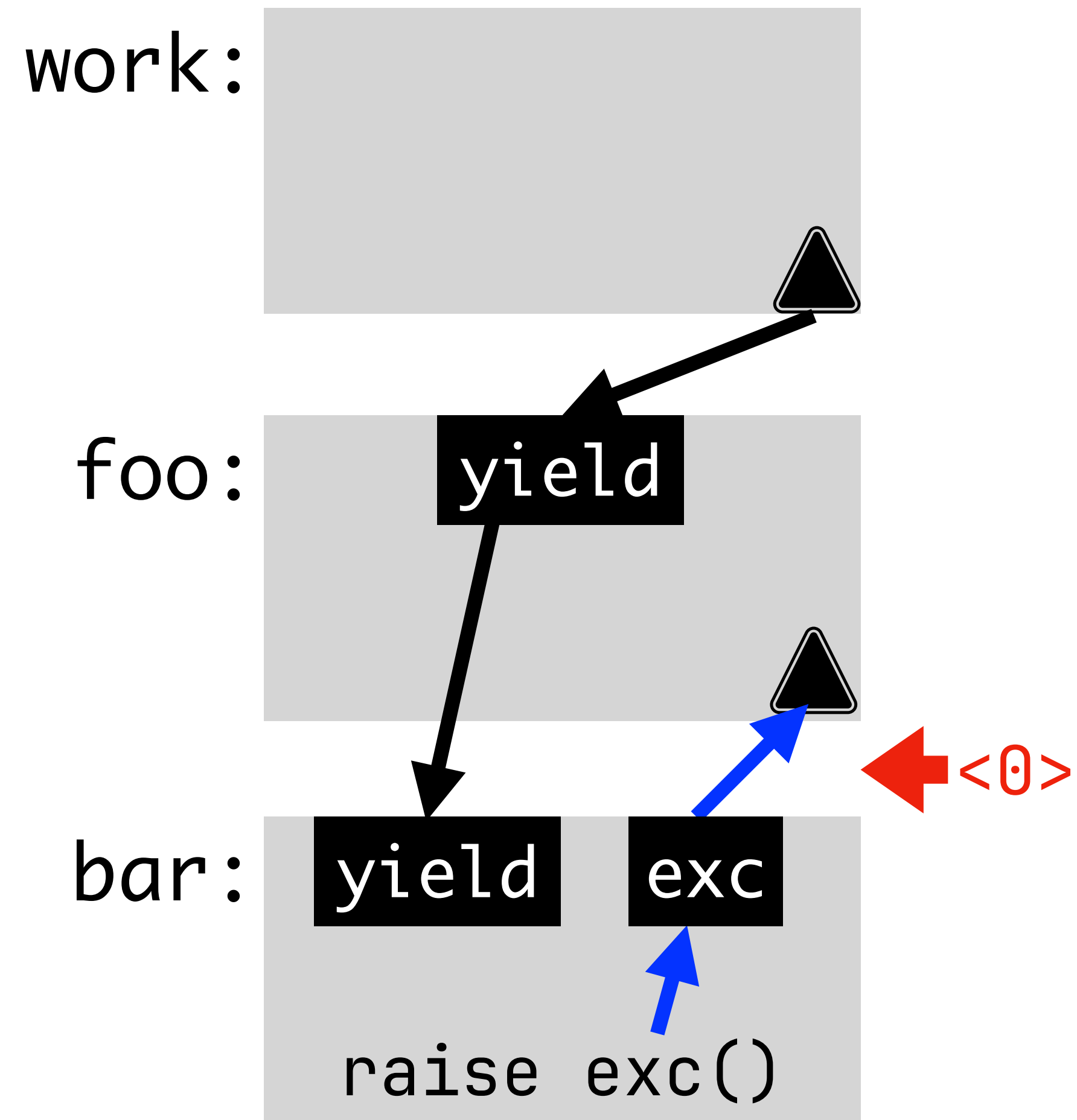
When an effect is raised, control goes to a runtime **stack-walker**.

Zero Lexa, Example 1

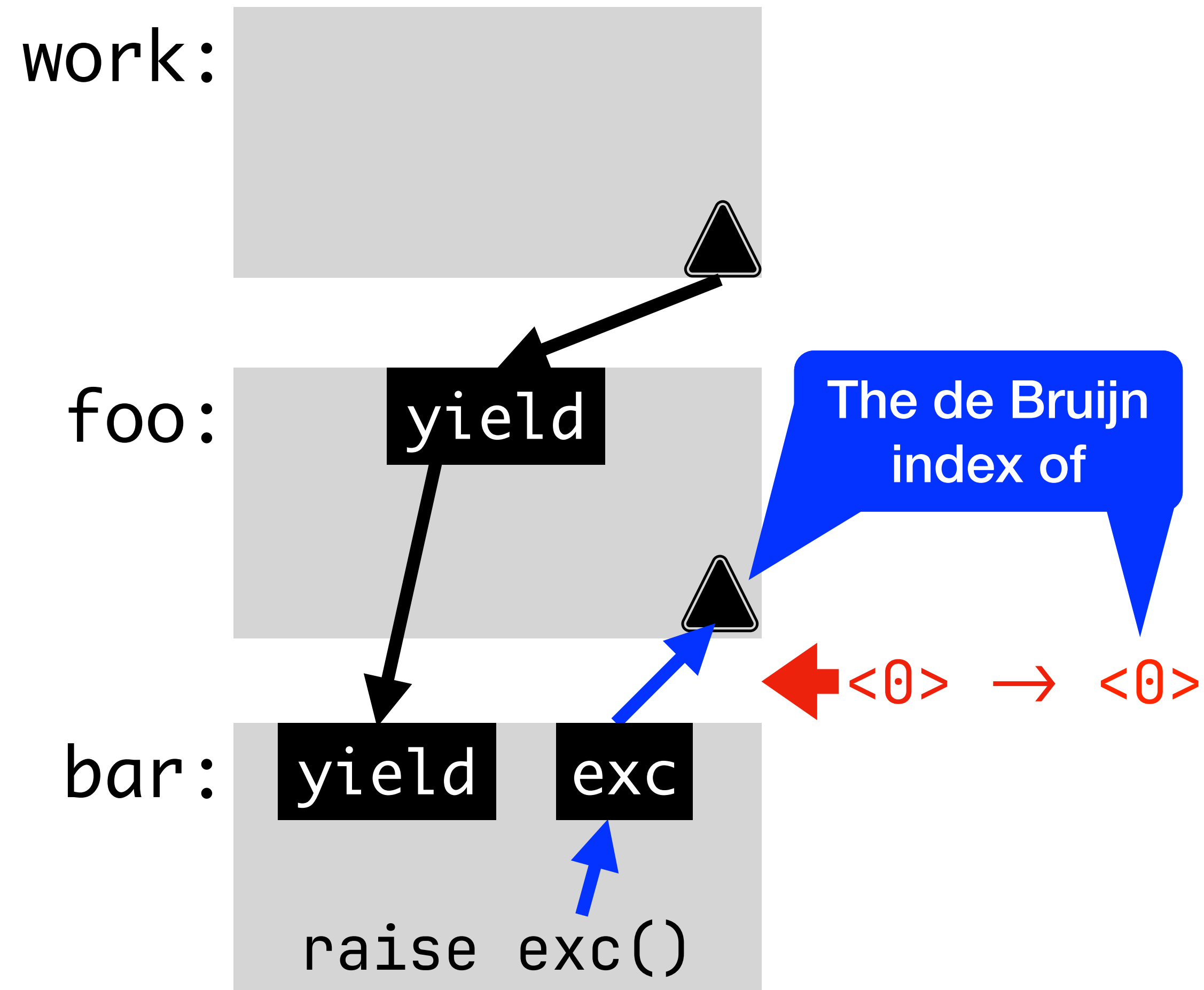


When an effect is raised, control goes to a runtime **stack-walker**.

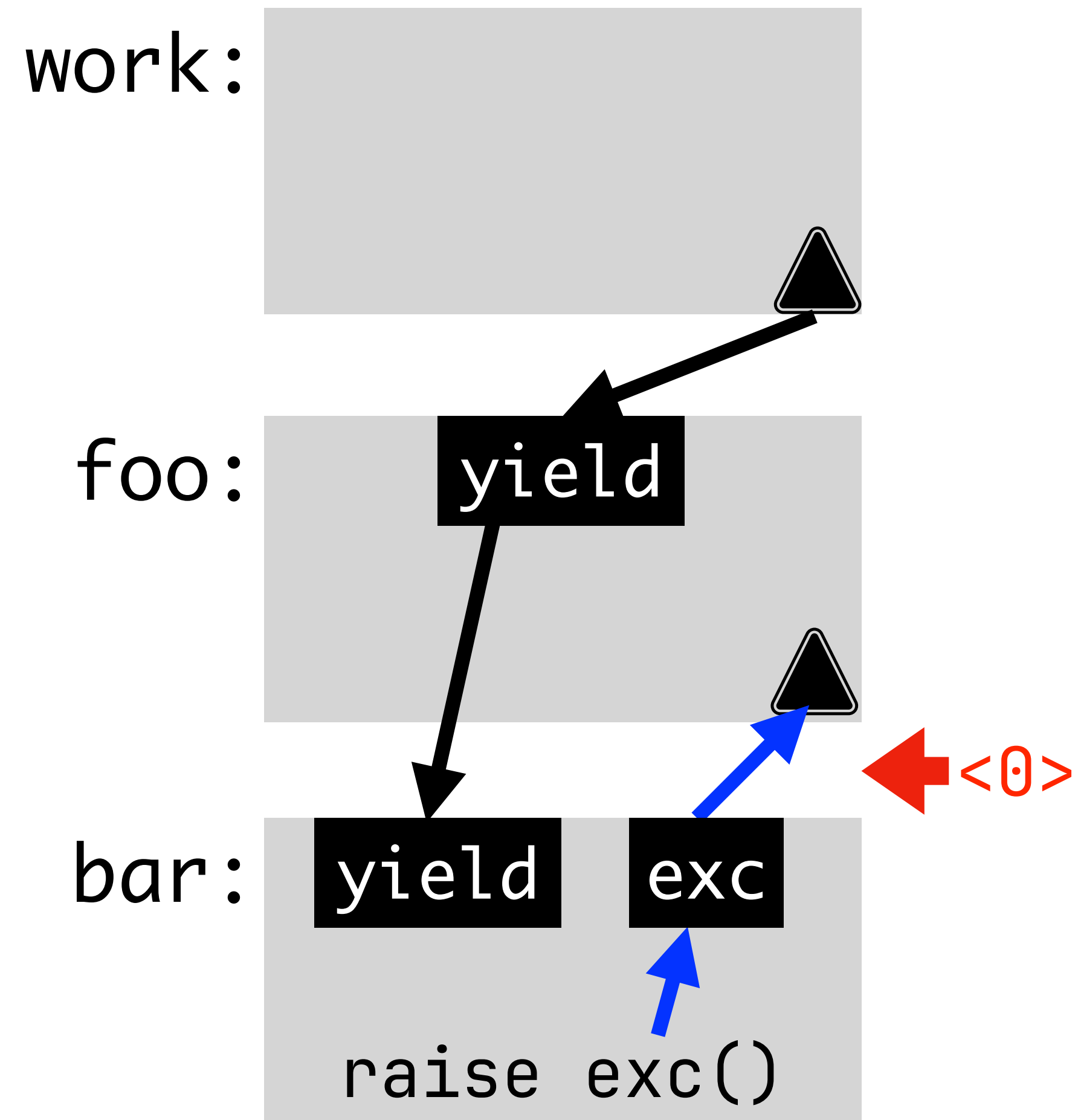
Zero Lexa, Example 1



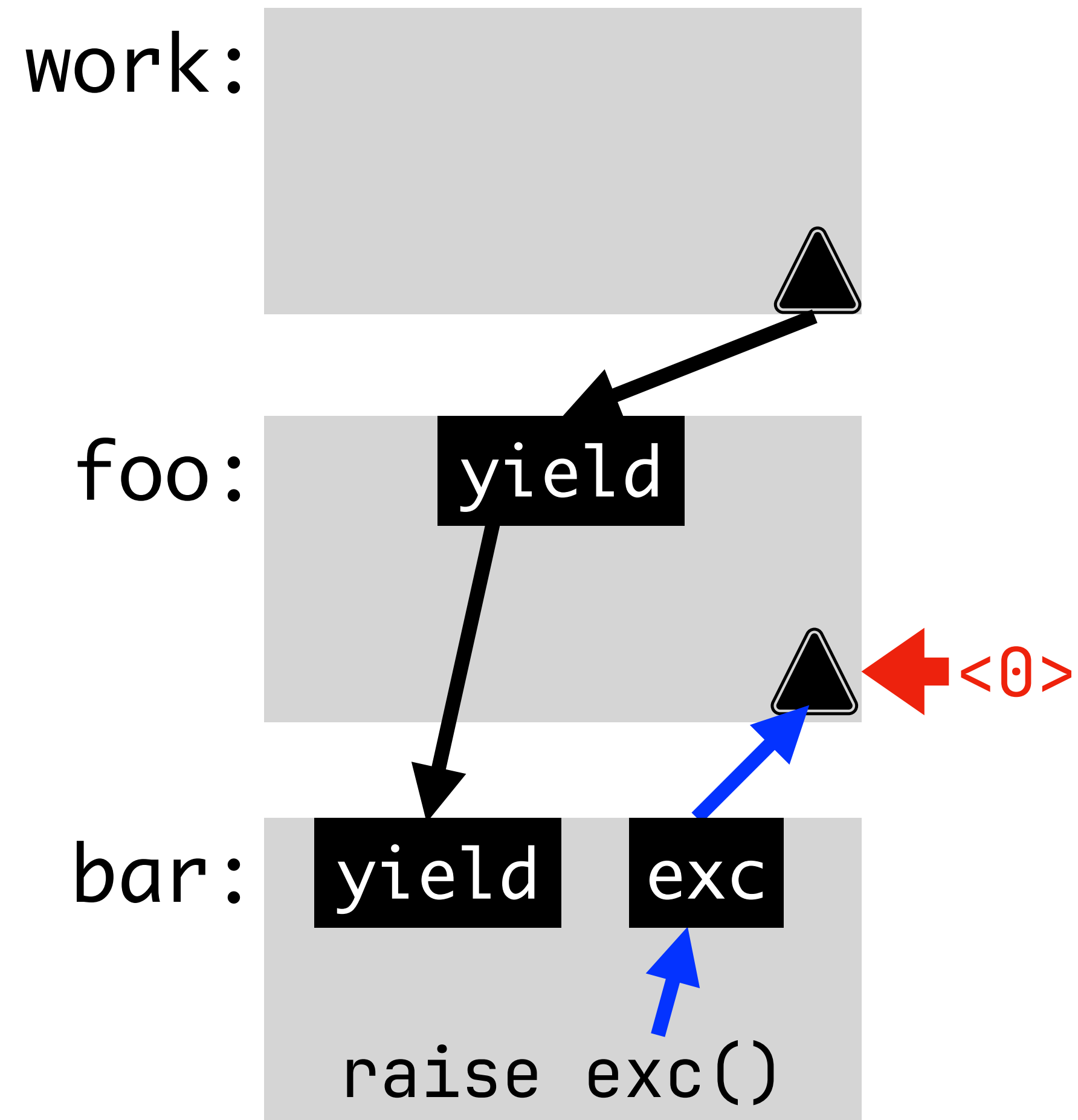
Zero Lexa, Example 1



Zero Lexa, Example 1



Zero Lexa, Example 1



Zero Lexa, Example 2

Our compilation can also deal with higher-order functions.

Zero Lexa, Example 2

```
def main():
    handle
    handle
    let
        bar =  $\lambda().\text{raise log}()$ ;
                raise exc()
    in
        foo[log, exc](bar)
    with log: Logging = ...
    with exc: Exception = ...

def foo[a](g):
    ...
    g()
    ...
```

Zero Lexa, Example 2

```
def main():
    handle
    handle
    let
        bar =  $\lambda().\text{raise log}();$ 
               $\text{raise exc}()$ 
    in
        foo[log, exc](bar)
    with log: Logging = ...
    with exc: Exception = ...

def foo[a](g):
    ...
    g()
    ...
```

Zero Lexa, Example 2

```
def main():  
    handle  
    handle  
    let  
        bar =  $\lambda().\text{raise log}()$ ;  
            raise exc()  
    in  
        foo[log, exc](bar)  
    with log: Logging = ...  
    with exc: Exception = ...
```

```
def foo[a](g):  
    ...  
    g()  
    ...
```

Zero Lexa, Example 2

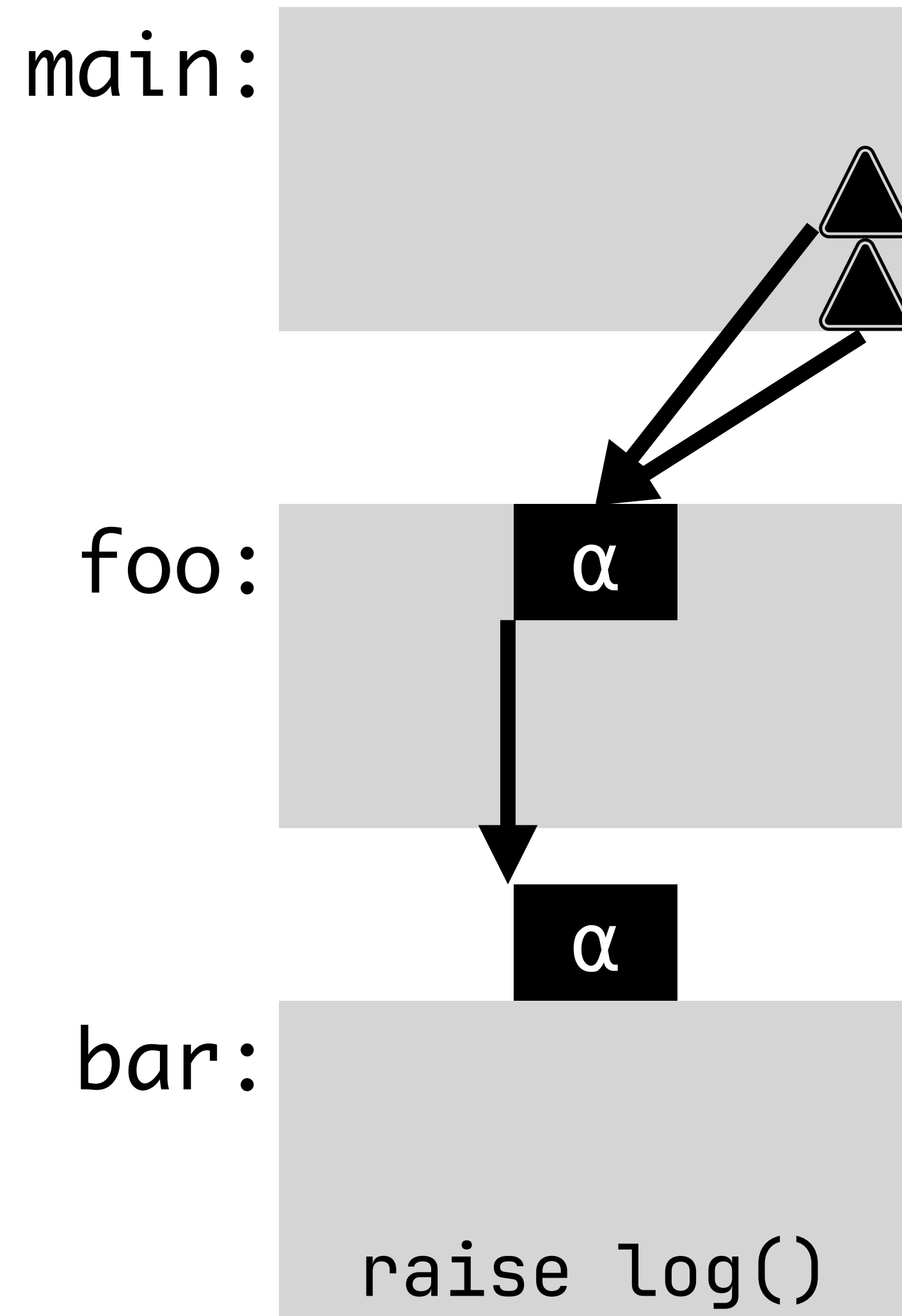
```
def main():
    handle
    handle
    let
        bar =  $\lambda().\text{raise log}();$ 
                raise exc()
    in
        foo[log, exc](bar)
    with log: Logging = ...
    with exc: Exception = ...

def foo[a](g):
    ...
    g()
    ...
```

Zero Lexa, Example 2

```
def main():
    handle
    handle
    let
        bar =  $\lambda().$ raise log();
              raise exc()
    in
        foo[log, exc](bar)
    with log: Logging = ...
    with exc: Exception = ...

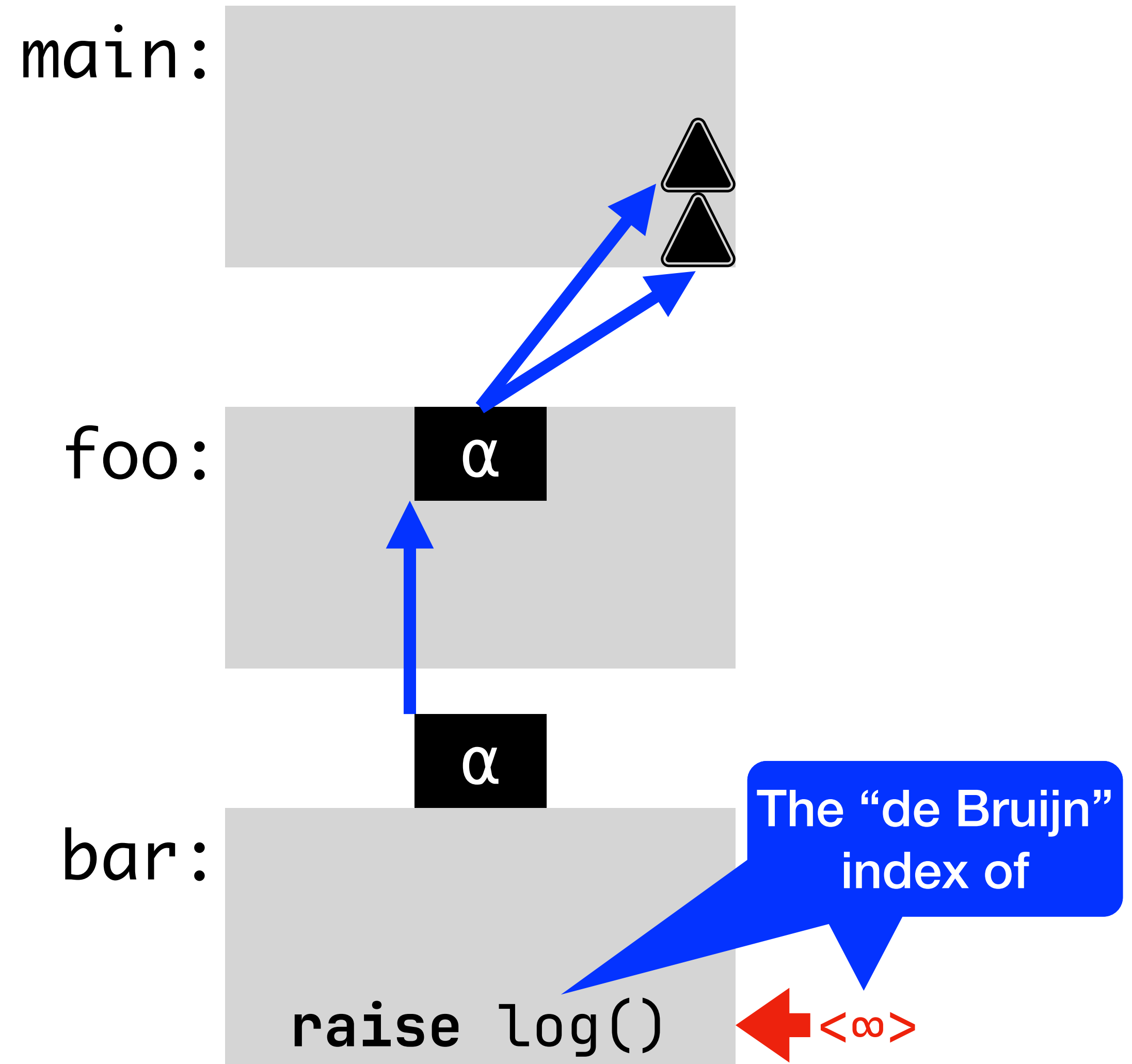
def foo[a](g):
    ...
    g()
    ...
```



Zero Lexa, Example 2

```
def main():
    handle
    handle
    let
        bar =  $\lambda().$ raise log();
              raise exc()
    in
        foo[log, exc](bar)
    with log: Logging = ...
    with exc: Exception = ...

def foo[a](g):
    ...
    g()
    ...
```

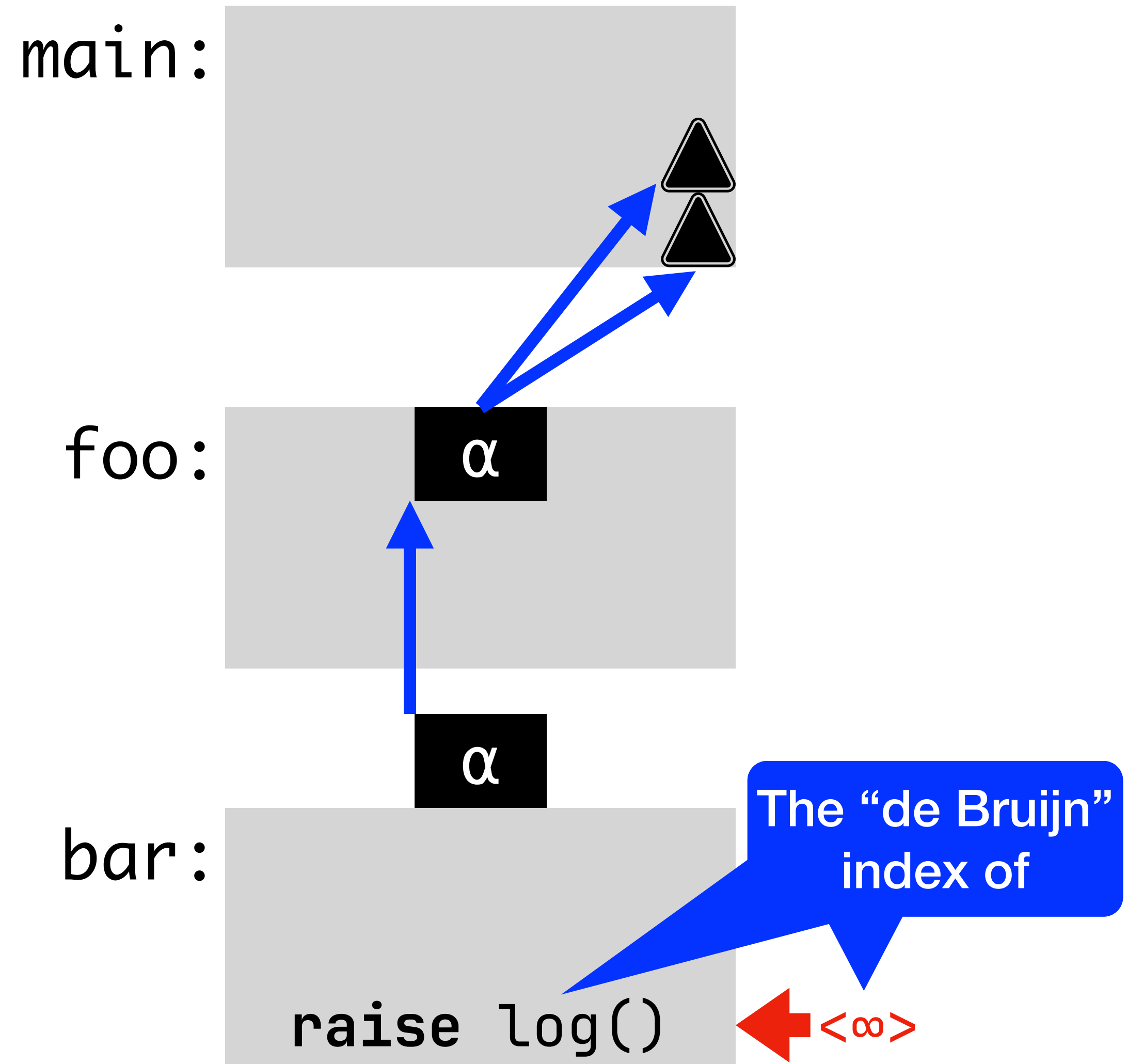


Zero Lexa, Example 2

```
def main():
    handle
    handle
    let
        bar =  $\lambda().$ raise log();
              raise exc()
    in
        foo[log, exc](bar)
    with log: Logging = ...
    with exc: Exception = ...

def foo[a](g):
    ...
    g()
    ...
```

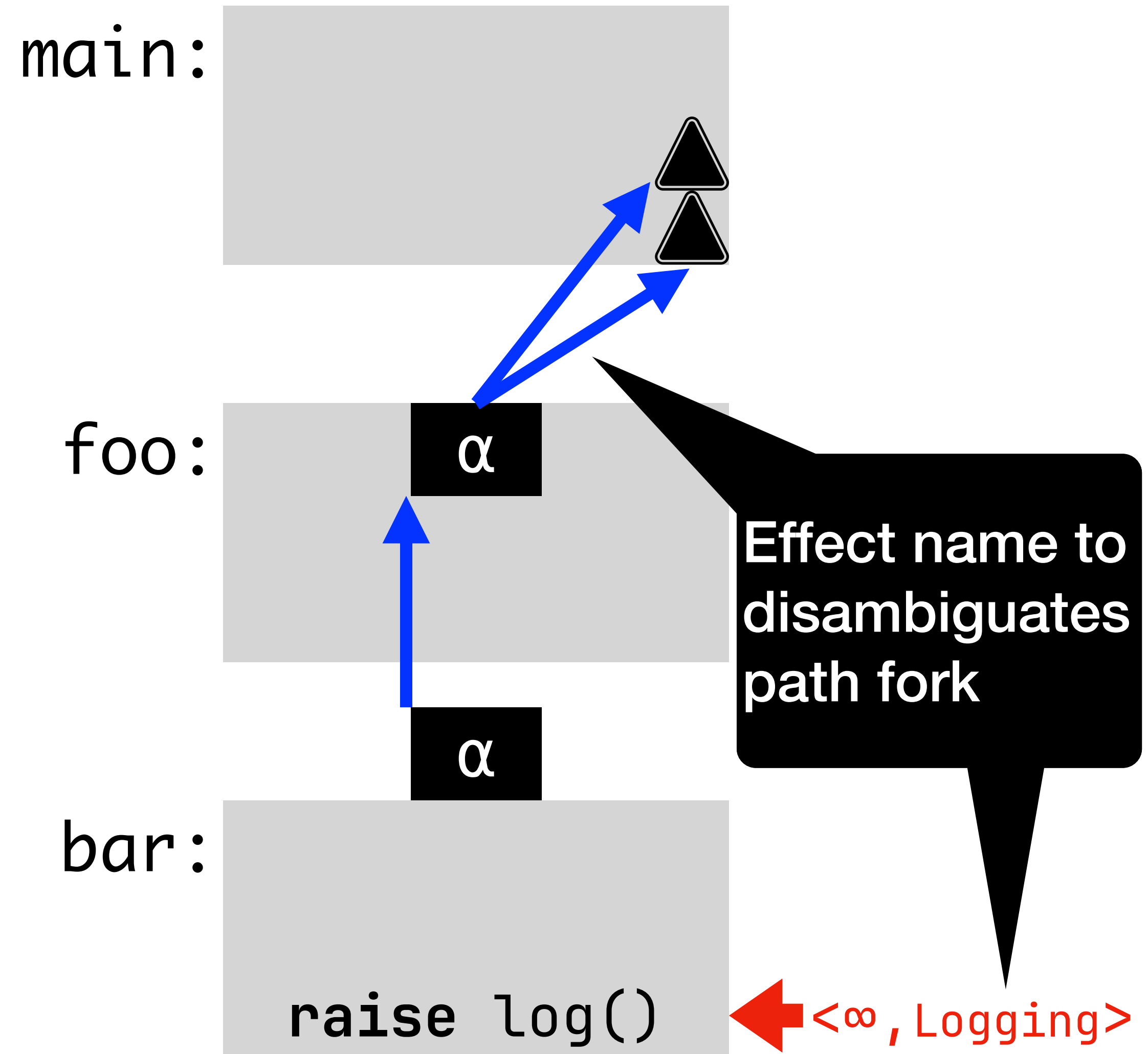
bound outside



Zero Lexa, Example 2

```
def main():
    handle
    handle
    let
        bar =  $\lambda().$ raise log();
              raise exc()
    in
        foo[log, exc](bar)
    with log: Logging = ...
    with exc: Exception = ...

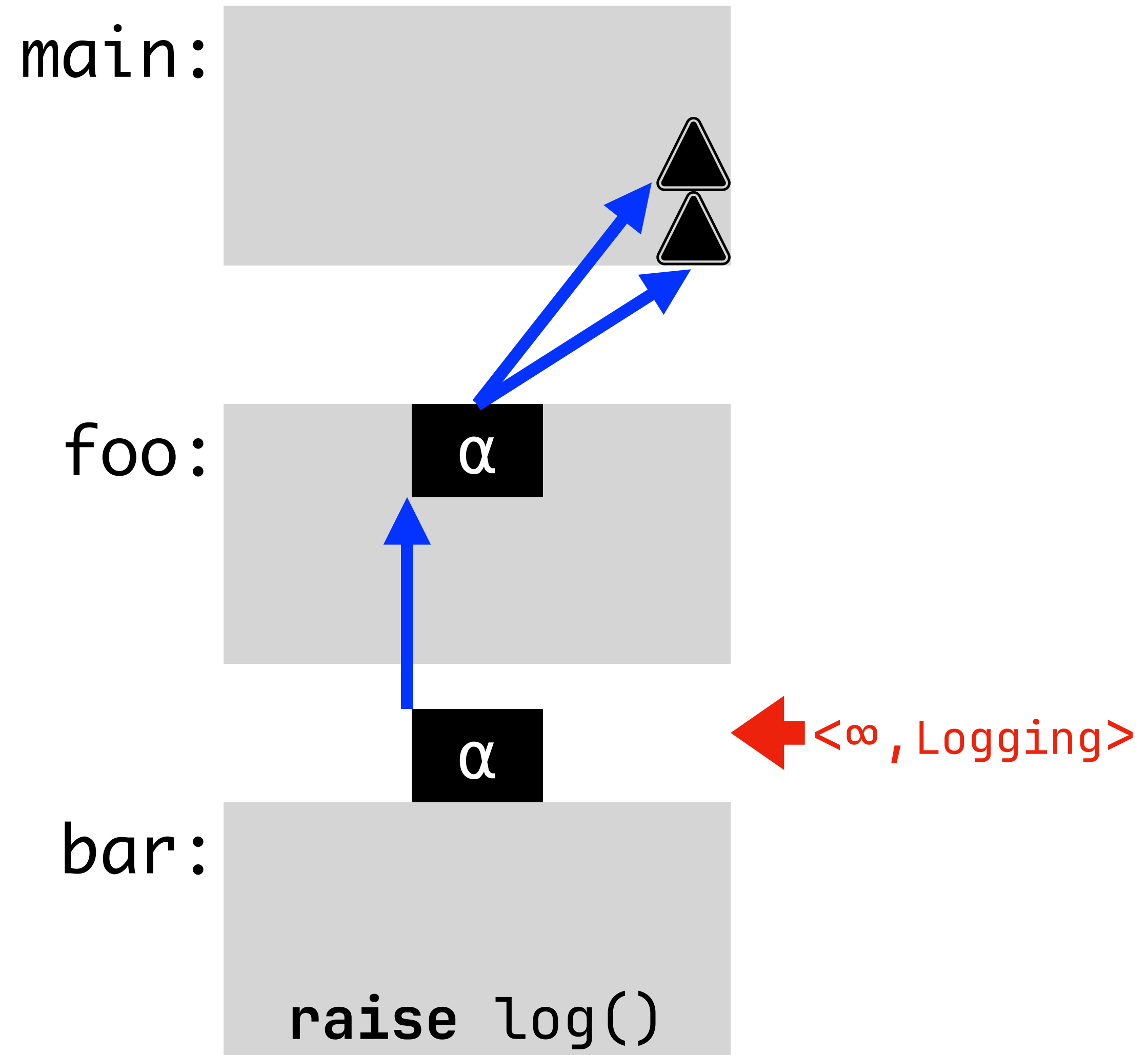
def foo[a](g):
    ...
    g()
    ...
```



Zero Lexa, Example 2

```
def main():
    handle
    handle
    let
        bar =  $\lambda().$ raise log();
              raise exc()
    in
        foo[log, exc](bar)
    with log: Logging = ...
    with exc: Exception = ...

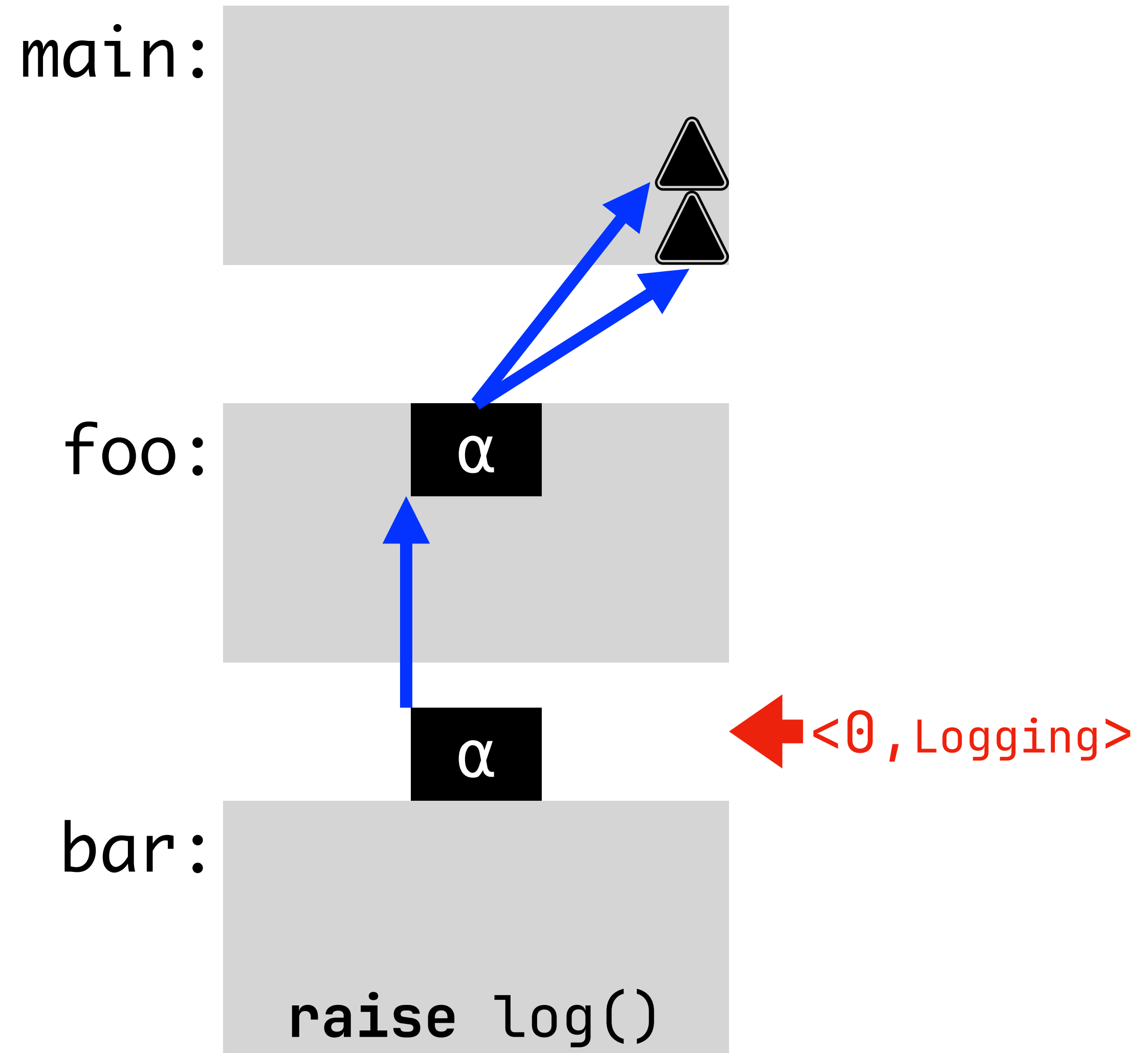
def foo[a](g):
    ...
    g()
    ...
```



Zero Lexa, Example 2

```
def main():
    handle
    handle
    let
        bar =  $\lambda().$ raise log();
              raise exc()
    in
        foo[log, exc](bar)
    with log: Logging = ...
    with exc: Exception = ...

def foo[a](g):
    ...
    g()
    ...
```

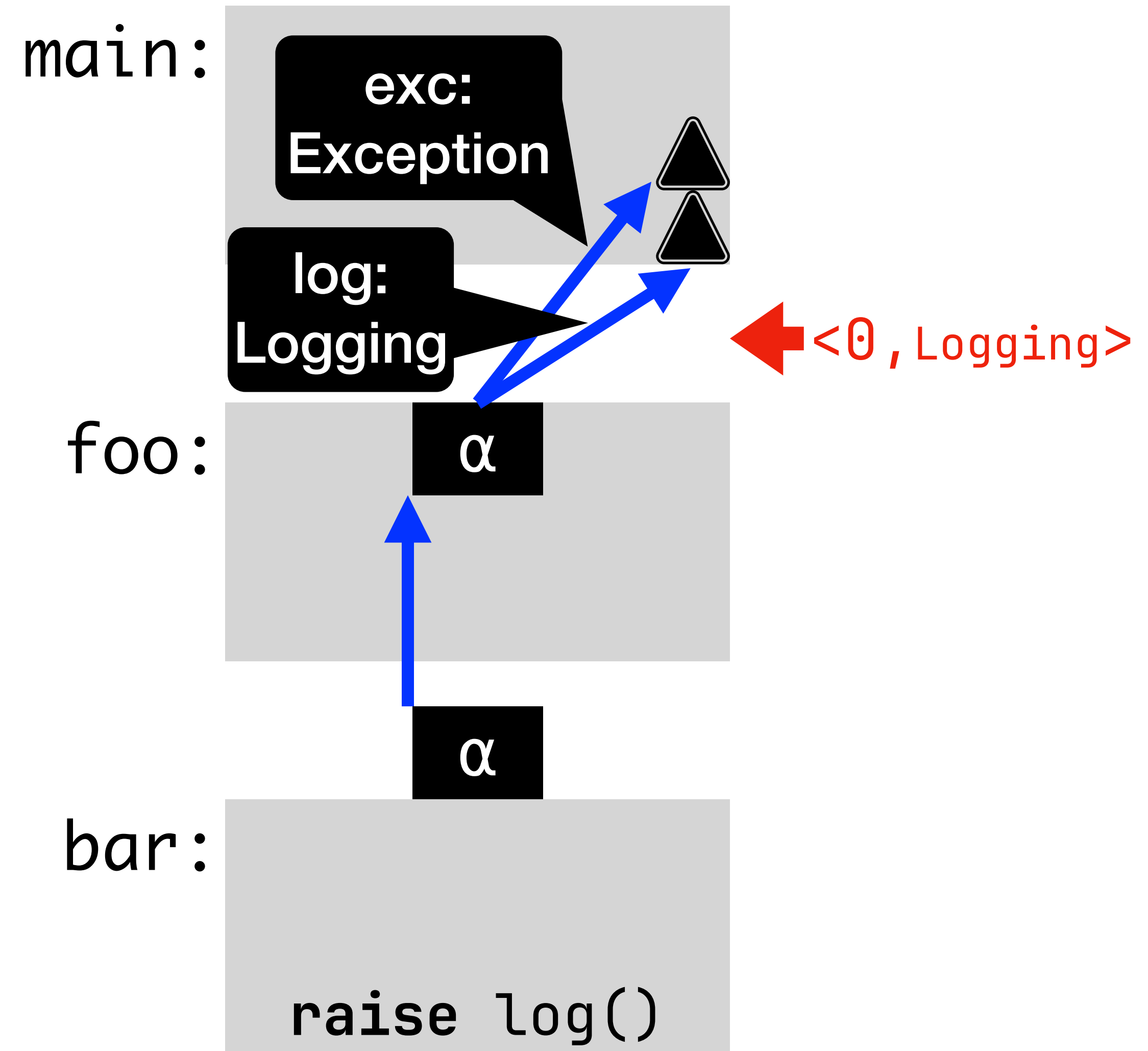


Zero Lexa, Example 2

```
def main():
    handle
    handle
    let
        bar =  $\lambda().\text{raise log}();$ 
              raise exc()

    in
        foo[log, exc](bar)
        with log: Logging = ...
        with exc: Exception = ...

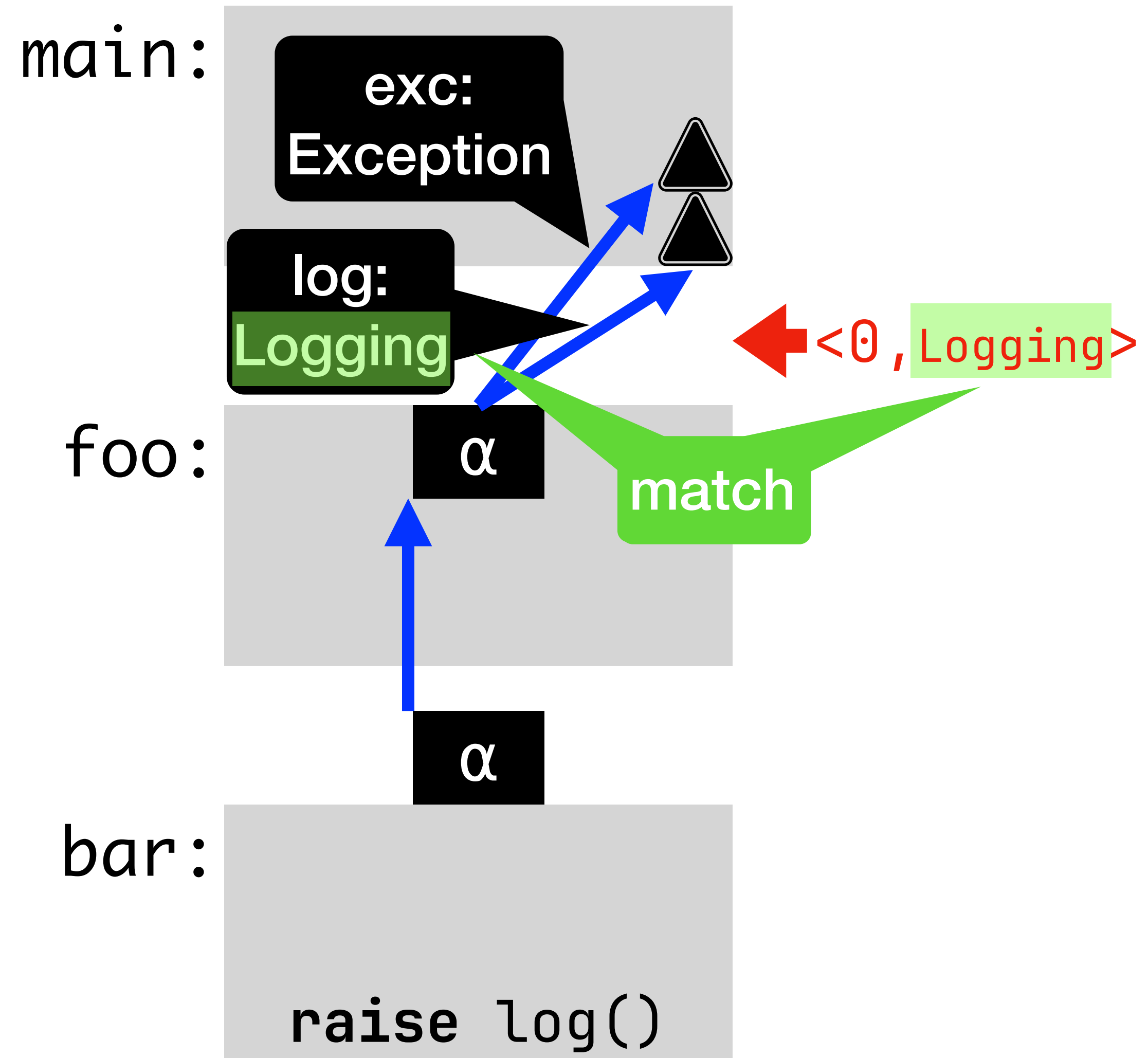
def foo[a](g):
    ...
    g()
    ...
```



Zero Lexa, Example 2

```
def main():
    handle
    handle
    let
        bar =  $\lambda().\text{raise log}();$ 
              raise exc()
    in
        foo[log, exc](bar)
        with log: Logging = ...
        with exc: Exception = ...

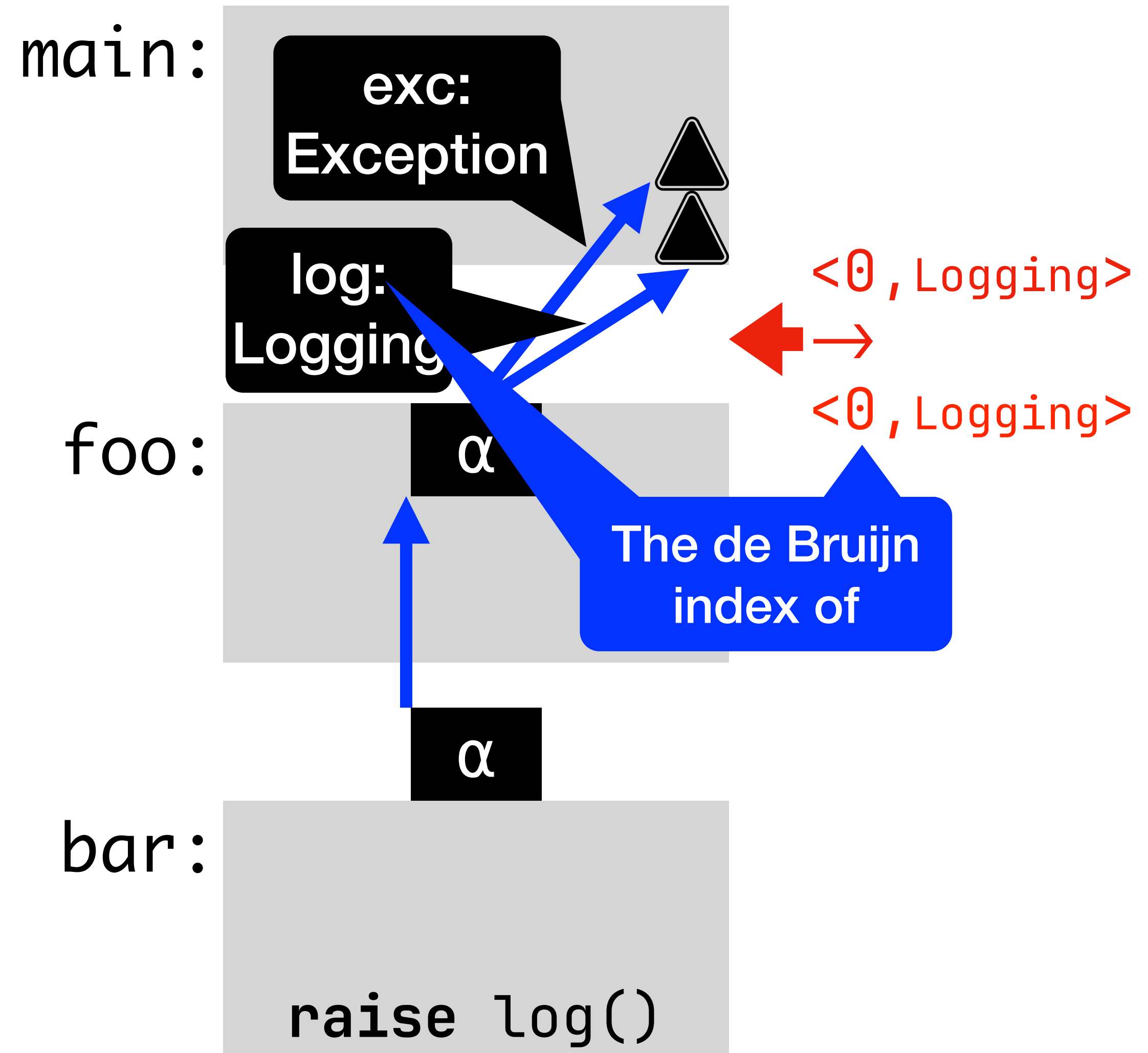
def foo[a](g):
    ...
    g()
    ...
```



Zero Lexa, Example 2

```
def main():
    handle
    handle
    let
        bar =  $\lambda().\text{raise log}();$ 
              raise exc()
    in
        foo[log, exc](bar)
        with log: Logging = ...
        with exc: Exception = ...

def foo[a](g):
    ...
    g()
    ...
```

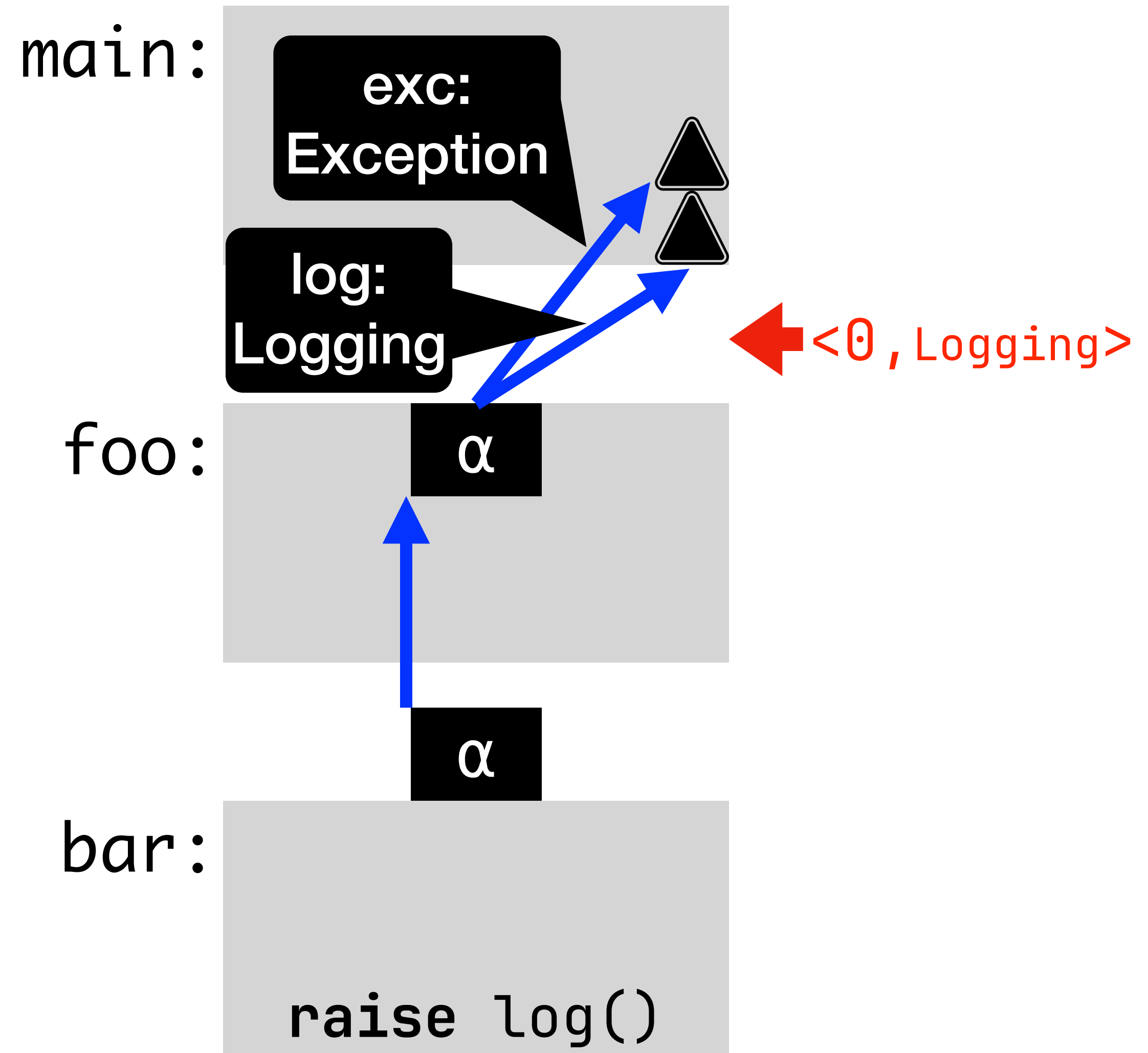


Zero Lexa, Example 2

```
def main():
    handle
    handle
    let
        bar =  $\lambda().\text{raise log}();$ 
              raise exc()

    in
        foo[log, exc](bar)
        with log: Logging = ...
        with exc: Exception = ...

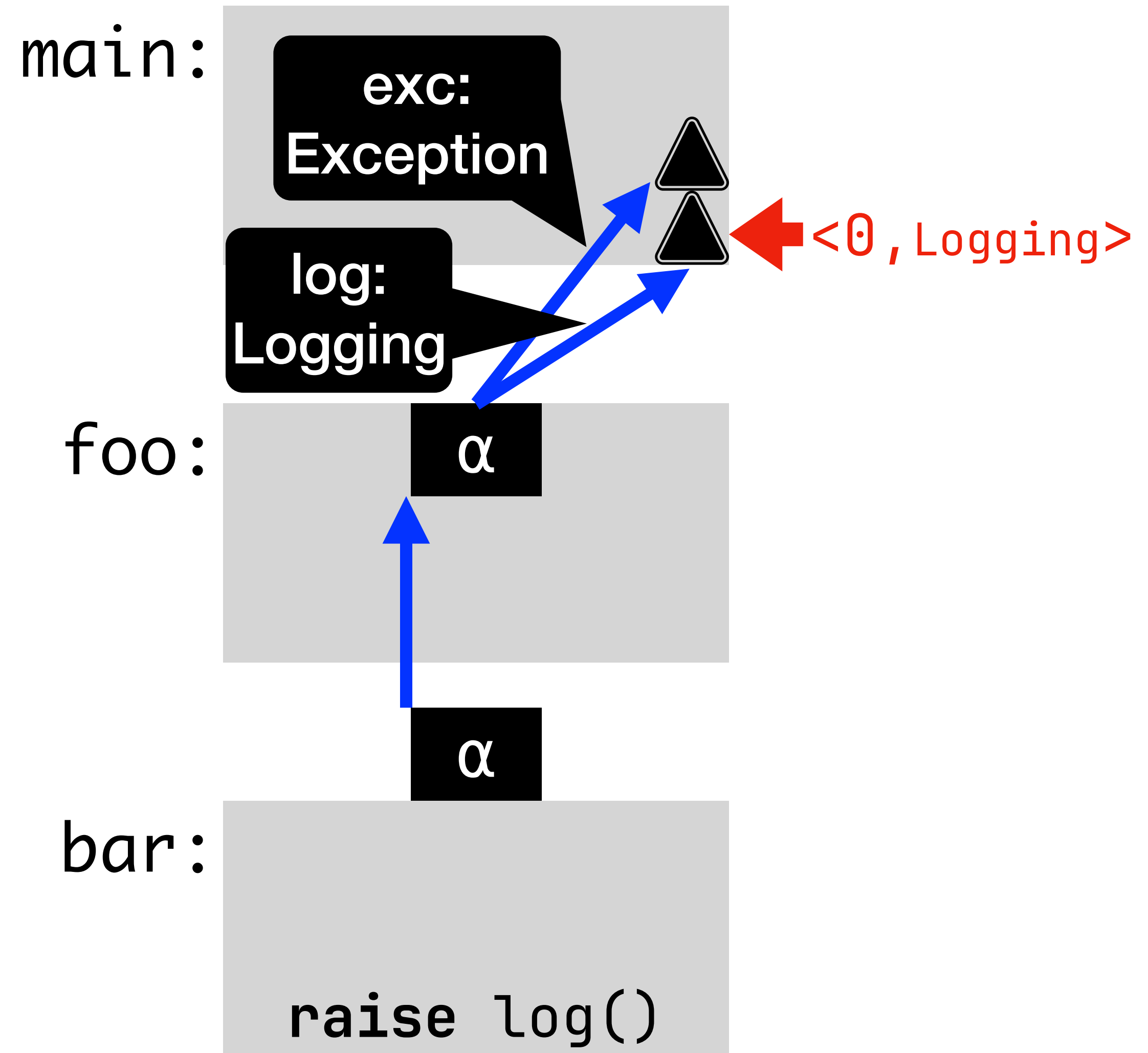
def foo[a](g):
    ...
    g()
    ...
```



Zero Lexa, Example 2

```
def main():
    handle
    handle
    let
        bar =  $\lambda().\text{raise log}();$ 
              raise exc()
    in
        foo[log, exc](bar)
    with log: Logging = ...
    with exc: Exception = ...

def foo[a](g):
    ...
    g()
    ...
```



Zero Lexa Implementation

Source Language



Target Language

```
handle  
  g(yield)  
with yield =  
  ...
```

```
handle  
  g()^{0→0}  
with _:  
  ...
```

callsite
metadata

Zero Lexa Implementation

Source Language



Target Language



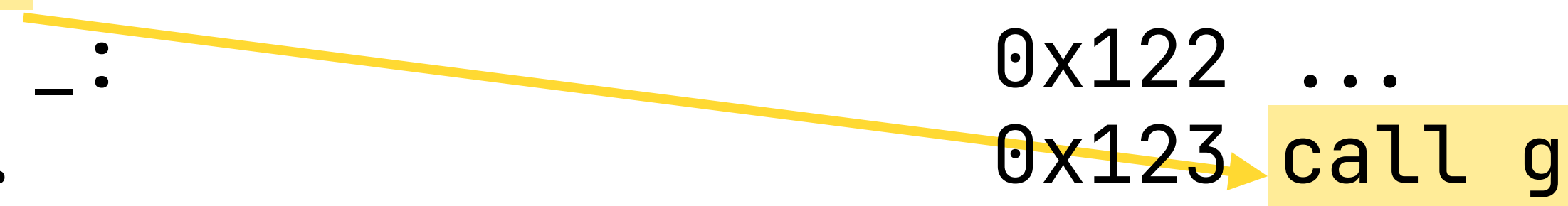
Binary

```
handle  
  g(yield)  
with yield =  
  ...
```

```
handle  
  g()^{0→0}  
with _:  
  ...
```

Code Segment:

```
0x122 ...  
0x123 call g  
0x124 ...
```



Zero Lexa Implementation

Source Language



Target Language



Binary

```
handle
  g(yield)
with yield =
  ...
```

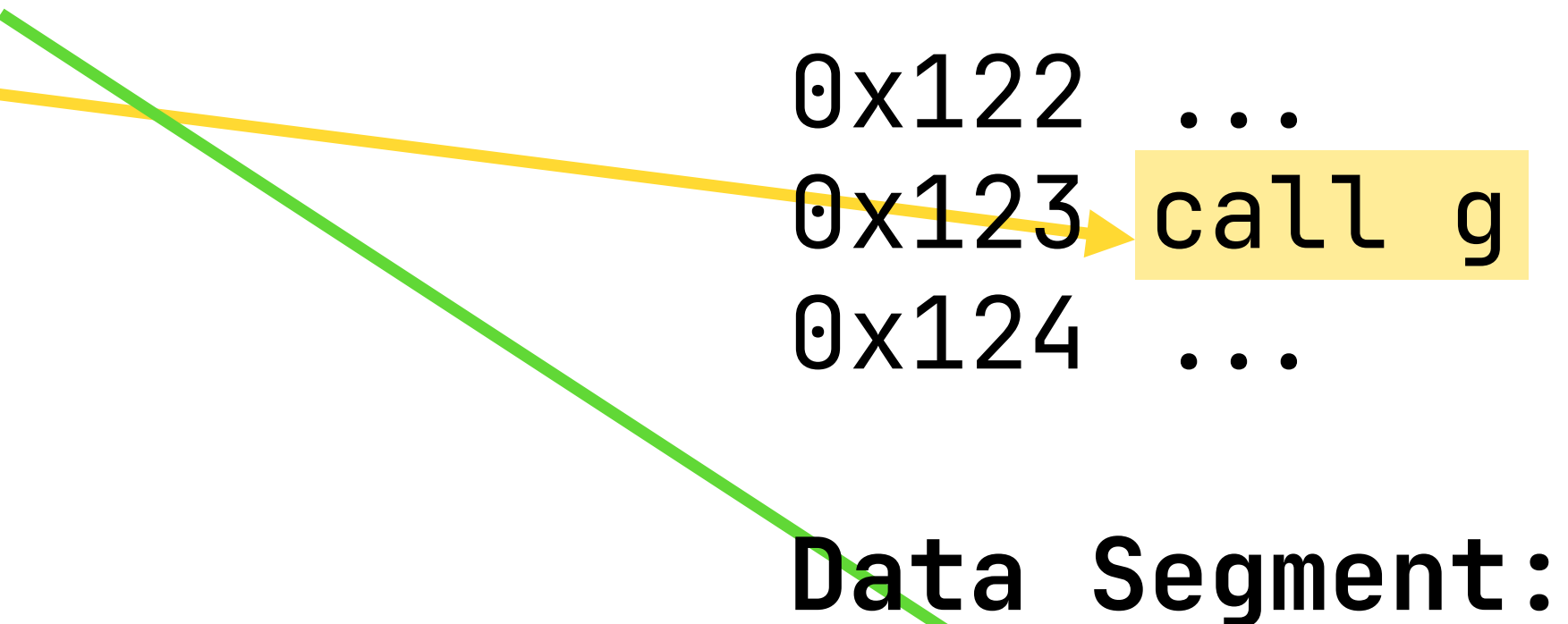
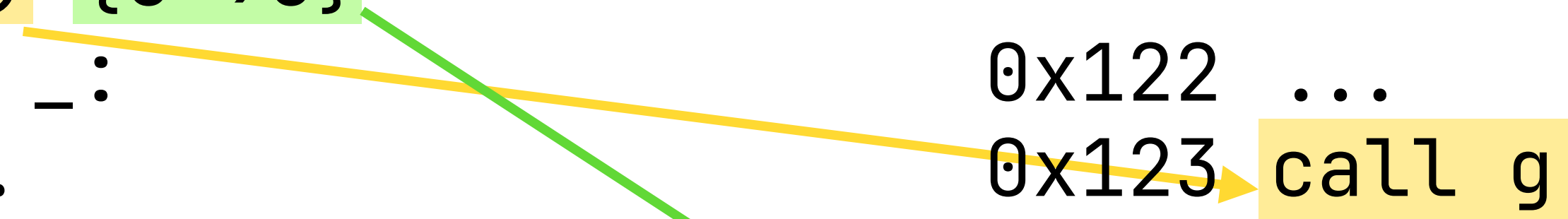
```
handle
  g()^{0→0}
with _:
  ...
```

Code Segment:

```
0x122 ...
0x123 → call g
0x124 ...
```

Data Segment:

```
0x83: { ... }
0x124: {0→0}
0x143: { ... }
```



Zero Lexa Implementation

Source Language



Target Language



Binary

```
handle
  g(yield)
with yield =
  ...
```

```
handle
  g()^{0→0}
with _:
  ...
```

Code Segment:

```
0x122 ...
0x123 call g
0x124 ...
```

Data Segment:

```
0x83: { ... }
0x124: {0→0}
0x143: { ... }
```

return
address of
call g

g()^{0→0}

with _:

0x123 call g

0x124

0x124: {0→0}

Zero Lexa

Lexa Compiler

```
effect State {  
  get: () → int  
  set: (int) → unit  
}
```

```
exceptional effect Exception {  
  exc: (int) → int  
}
```

...

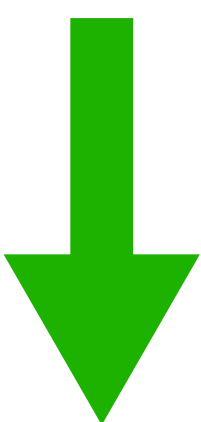
Evaluation

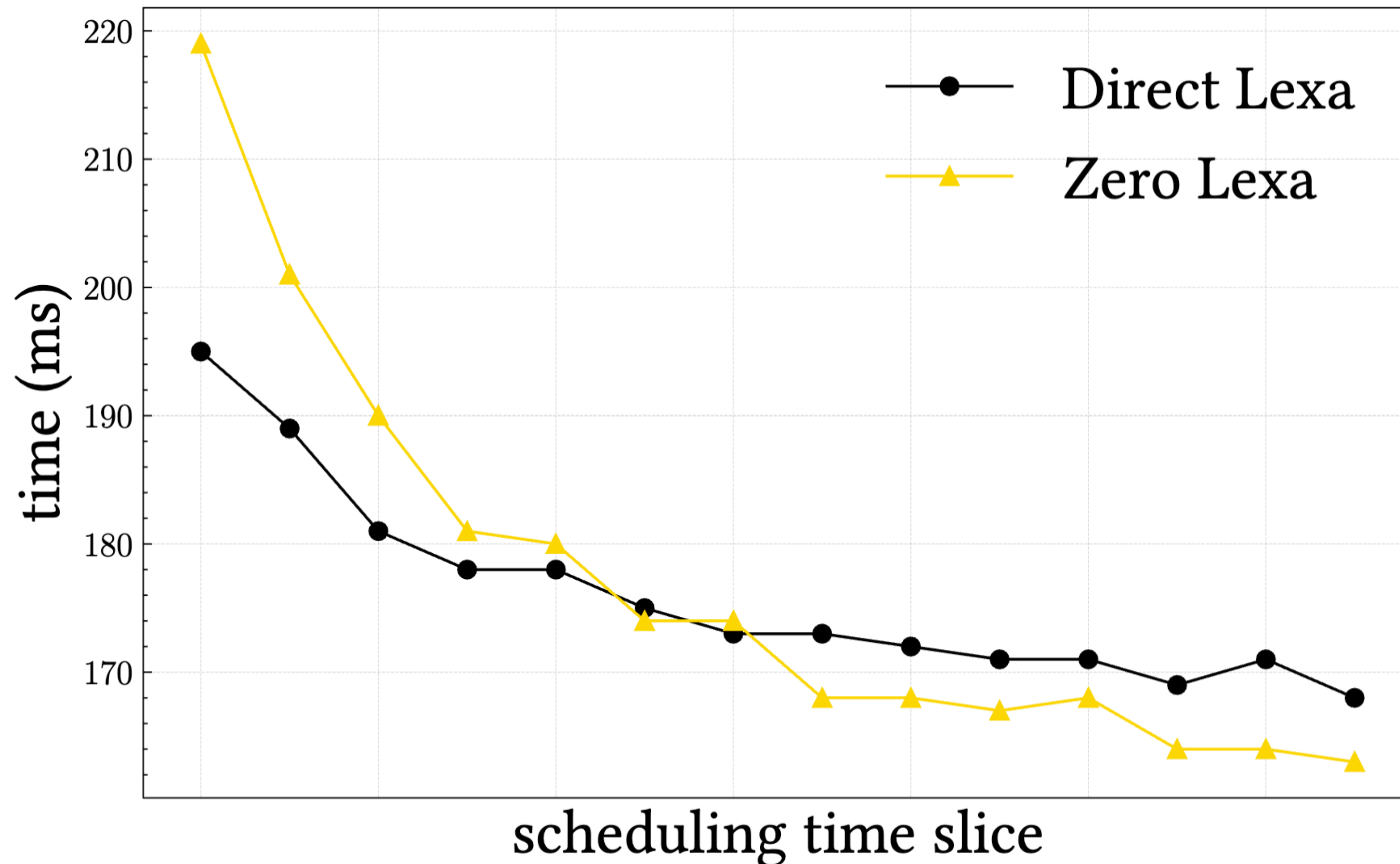
Hypothesis: zero-overhead implementation benefits exceptional handlers that are rarely used.

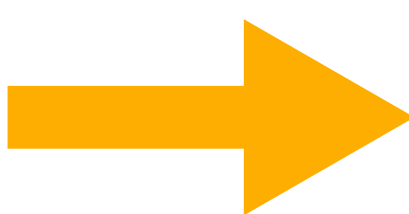
Evaluation

Case Study

A program with two cooperatively scheduled co-routine.


Better




**Infrequent
Yield**

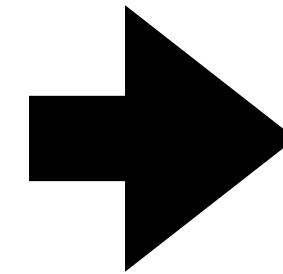
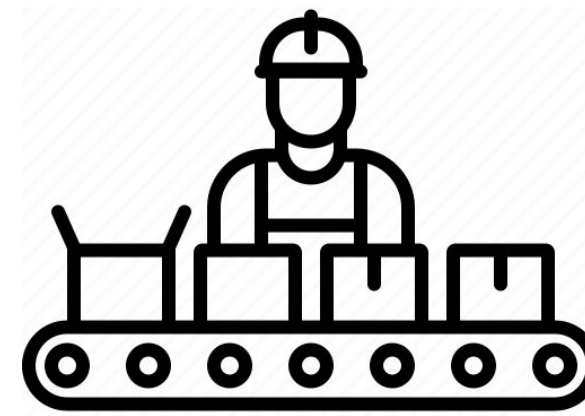
Evaluation

On new benchmarks

Benchmarks	Direct Lexa (ms)	Zero Lexa (ms)	Loss / Gain Zero vs. Direct
Catalan	344	301	12.5%
Bézout	680	671	1.3%
Golomb	651	563	13.5%
Hofstadter Q	576	559	3.0%
Karatsuba	804	702	12.7%
Ackermann	3975	3954	0.5%
Palindrome Partition	112	113	-0.9%
Lattice Path	1002	971	3.1%
Two Threads	153	140	8.5%

high-level effect handlers in Lexa

```
handle E with H  
  
raise ...  
  
resume ...
```



low-level stack switching in assembly

```
ENTER  
  
RAISE  
  
RESUME
```

Direct Lexa

A large, hollow black triangle pointing upwards. Inside the triangle, the text "tradeoffs between performance and expressivity" is centered.

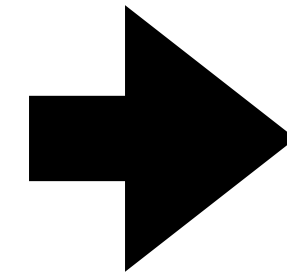
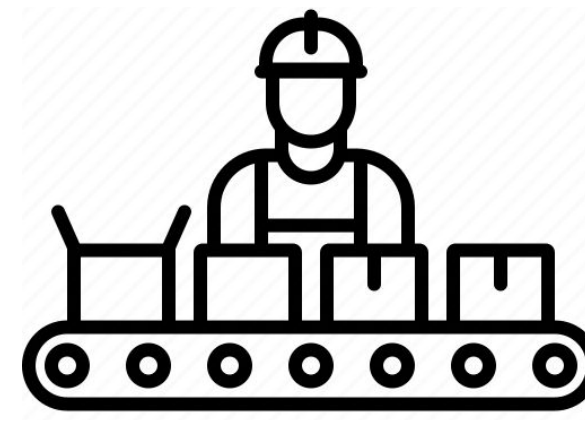
tradeoffs
between
performance
and expressivity

Multi Lexa

Zero Lexa

high-level effect handlers in Lexa

```
handle E with H  
  
raise ...  
  
resume ...
```



low-level stack switching in assembly

```
ENTER  
  
RAISE  
  
RESUME
```

Direct Lexa

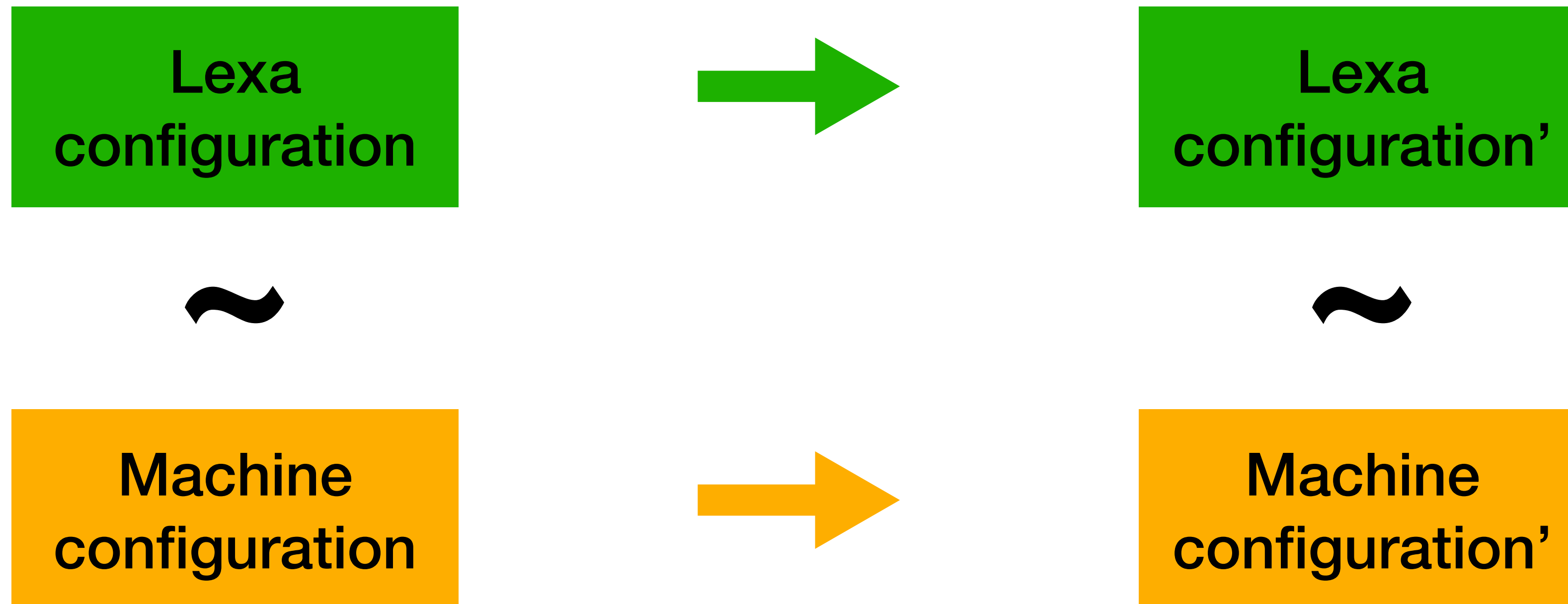
formally
proved
correct
[[*E*]] = [[*E*]]

Multi Lexa

Zero Lexa

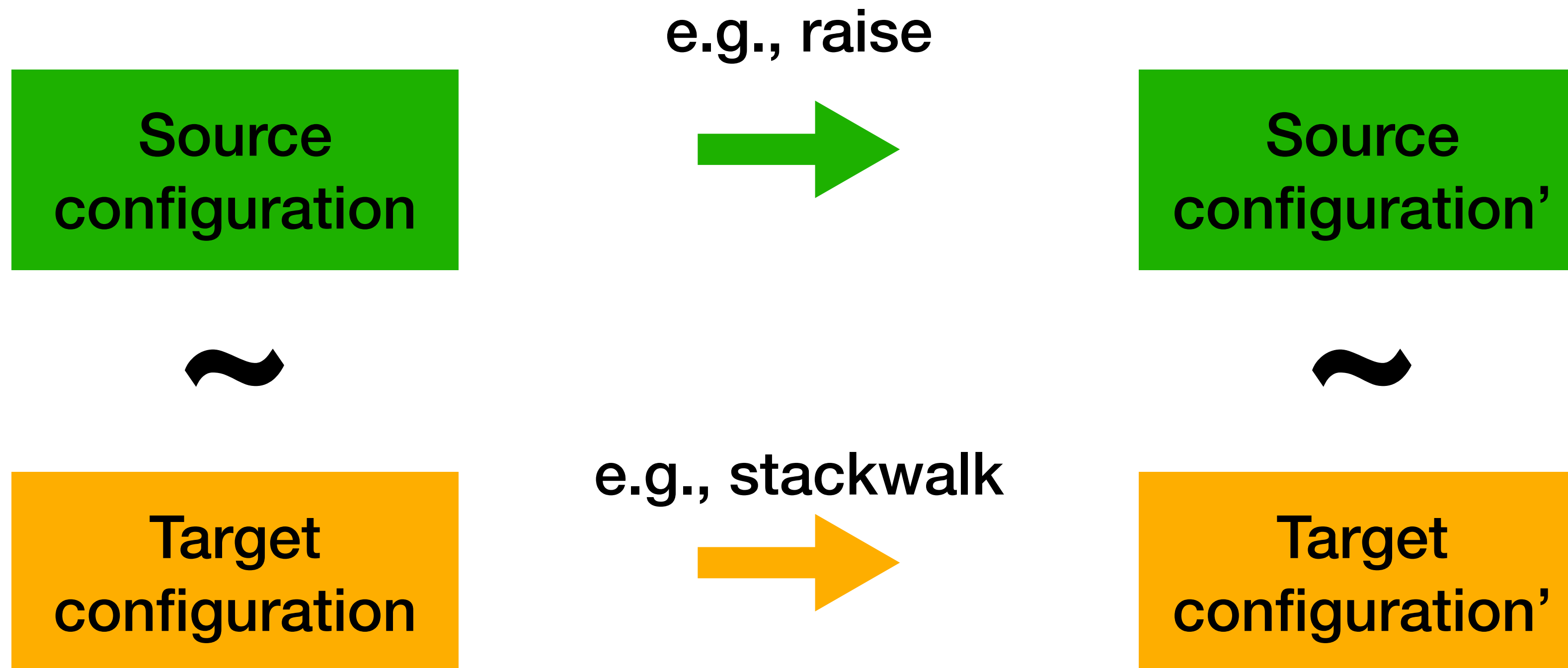
Compiler Correctness

CompCert-style Simulation Proofs for Direct, Multi & Zero Lexa



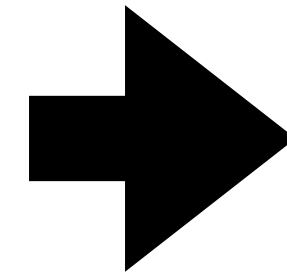
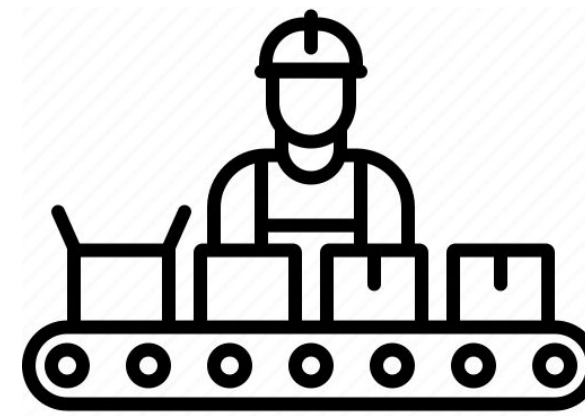
Compiler Correctness

e.g., Simulation Proofs for Zero Lexa



high-level effect handlers in Lexa

```
handle E with H  
  
raise ...  
  
resume ...
```



low-level stack switching in assembly

```
ENTER  
  
RAISE  
  
RESUME
```

Direct Lexa

A large, hollow black triangle pointing upwards. Inside the triangle, the text "A future of modular and efficient software" is centered and stacked vertically.

A
future
of
modular
and efficient
software

Multi Lexa

Zero Lexa